

# Practical 3D Geographic Routing for Wireless Sensor Networks

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# Geographic Routing Algorithms

Exploit geometric information  
(coordinates) of network  
topology to improve scalability  
of point-to-point routing

# Geographic Routing Algorithms

Greedy forwarding +  
Recovery mode when local  
minimum is encountered

1. Efficient

2. Storage

proportional to

density, not size

# Motivation

Previously proposed geographic routing algorithms assume “planar” network topology

⇒ Many modern sensor networks are three-dimensional

# Two Questions

1. How do we get geographic routing to work for 3D networks?

# Two Questions

2. How do existing point-to-point algorithms compare?

*(should we care?)*

# Outline

- Problem & Motivation
- Overview of related work & geographic routing
- Our Solution: GDSTR-3D
- Performance Evaluation
- Conclusion
- Future Work



# Related work

- 2D geographic routing
  - GPSR (Karp & Kung, Mobicom 2000)
  - GOAFR+ family (Kuhn et al., Mobihoc 2003)
  - CLDP (Kim et al., NSDI 2005)
  - GDSTR (Leong et al., NSDI 2006)
- 3D geographic routing
  - GRG (Flury & Wattenhofer, Infocom 2008)
  - GHG (Liu & Wu, Infocom 2009)

# Related work

## ➤ Point-to-point

- AODV (Perkins, Milcom 1997)
- VPCR (Newsome & Song, SenSys 2003)
- BVR (Fonseca et al., NSDI 2005)
- VRR (Caesar et al., SIGCOMM 2006)
- S4 (Mao et al., NSDI 2007)

## ➤ Virtual Coordinates

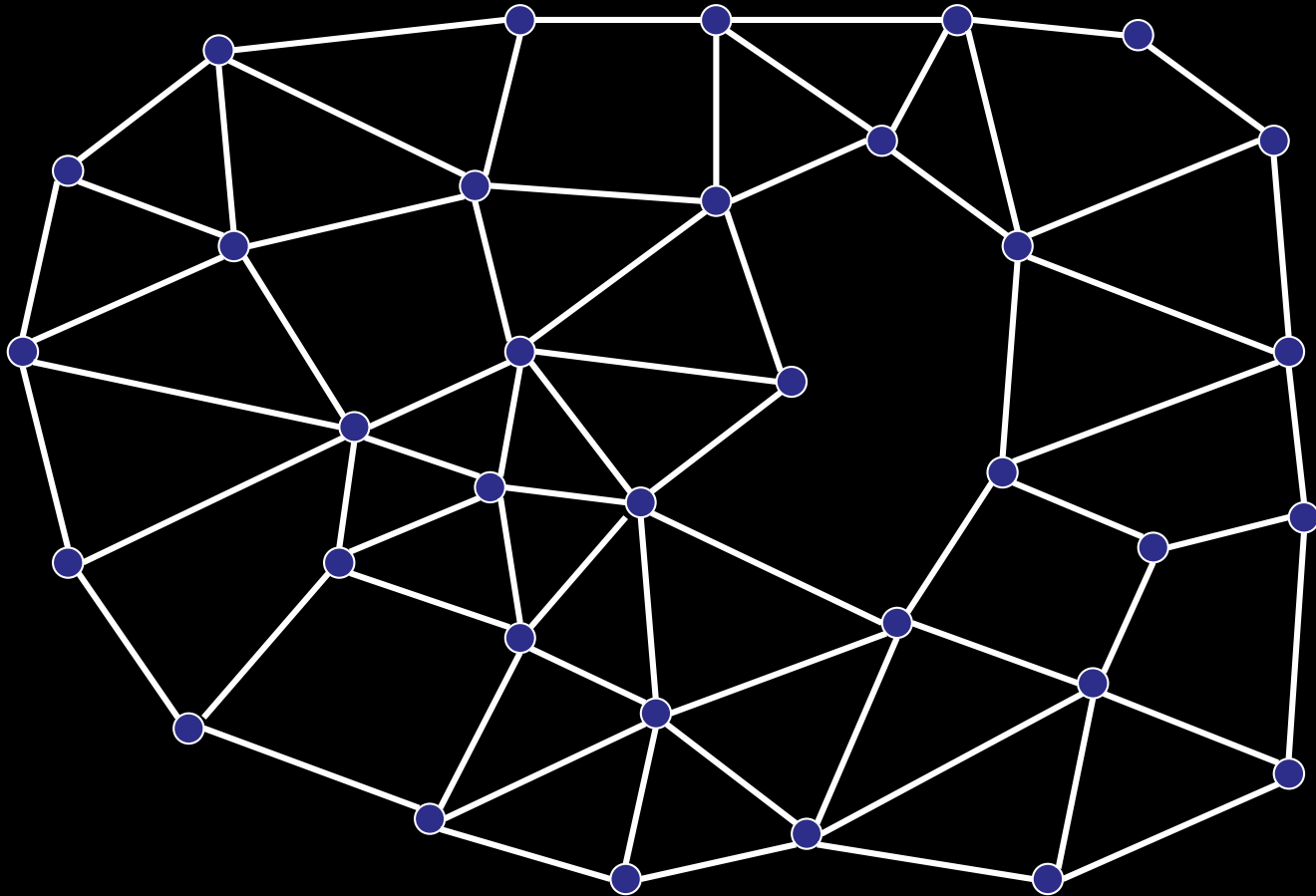
- NoGeo (Rao et al., Mobicom 2003)
- PSVC (Zhou et al., ICNP 2010)

Our Approach

Extend GDSTR  
to 3D

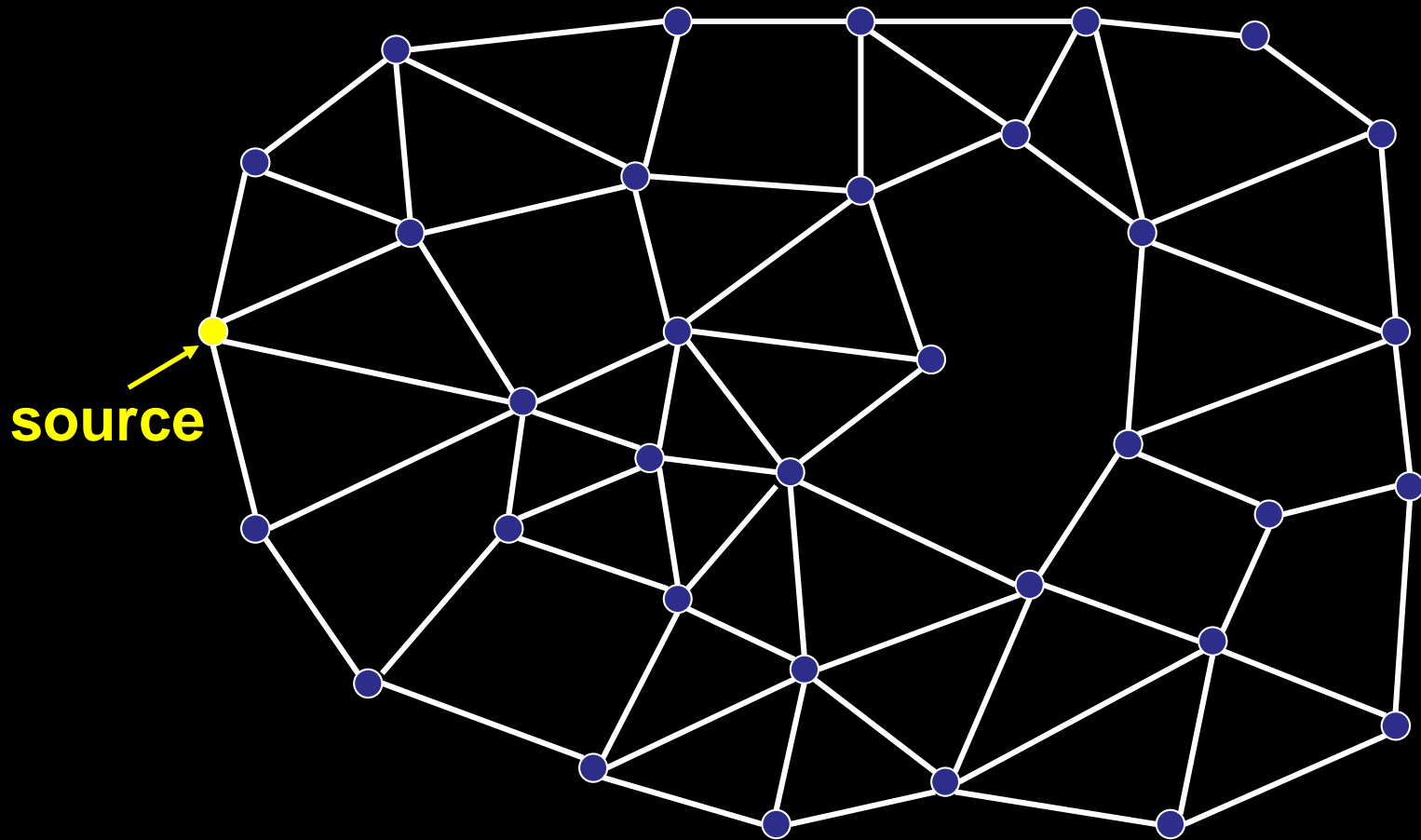
Complications!

# Overview: Geographic Routing



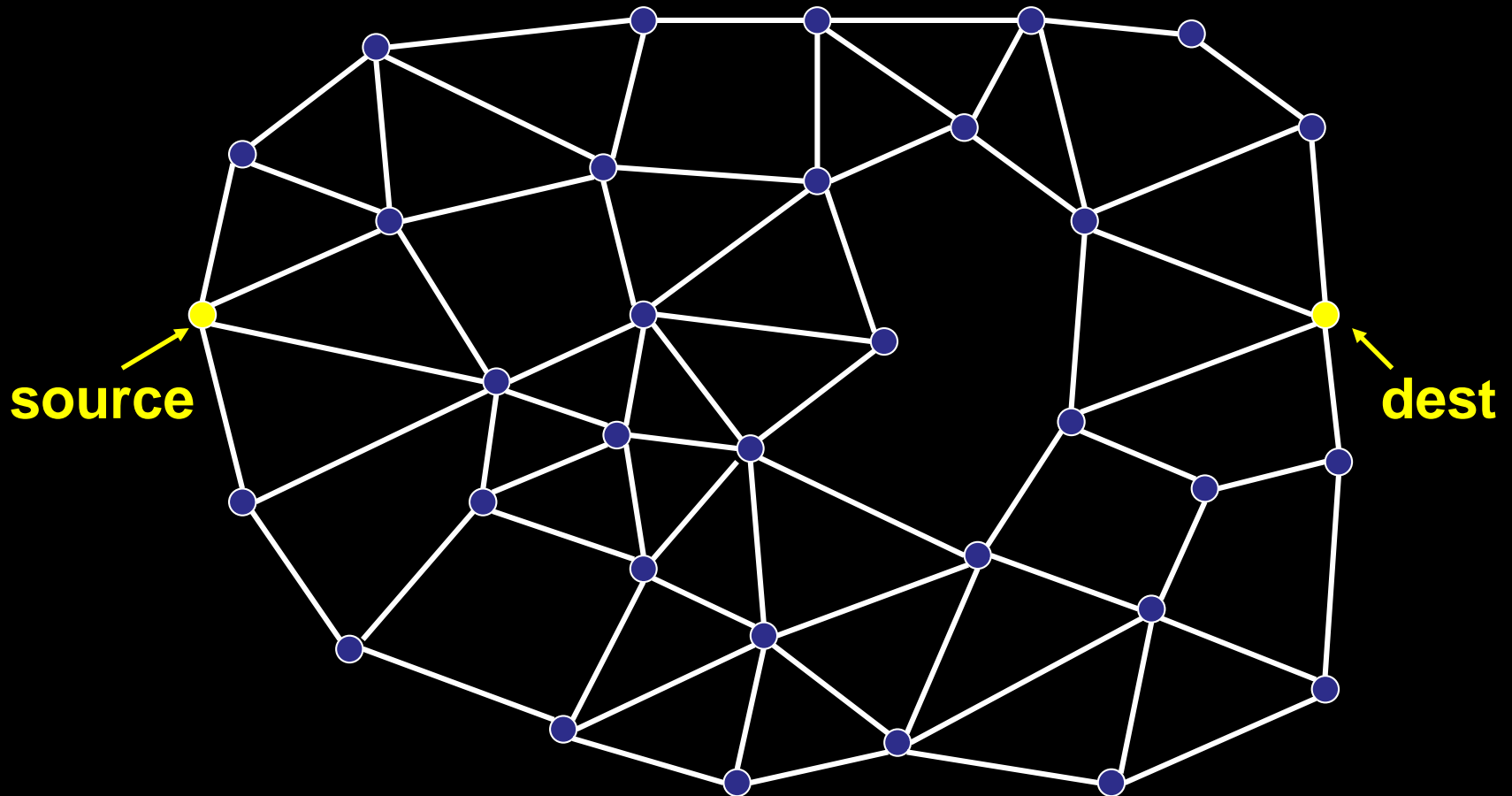
**Nodes have coordinates**

# Overview: Geographic Routing



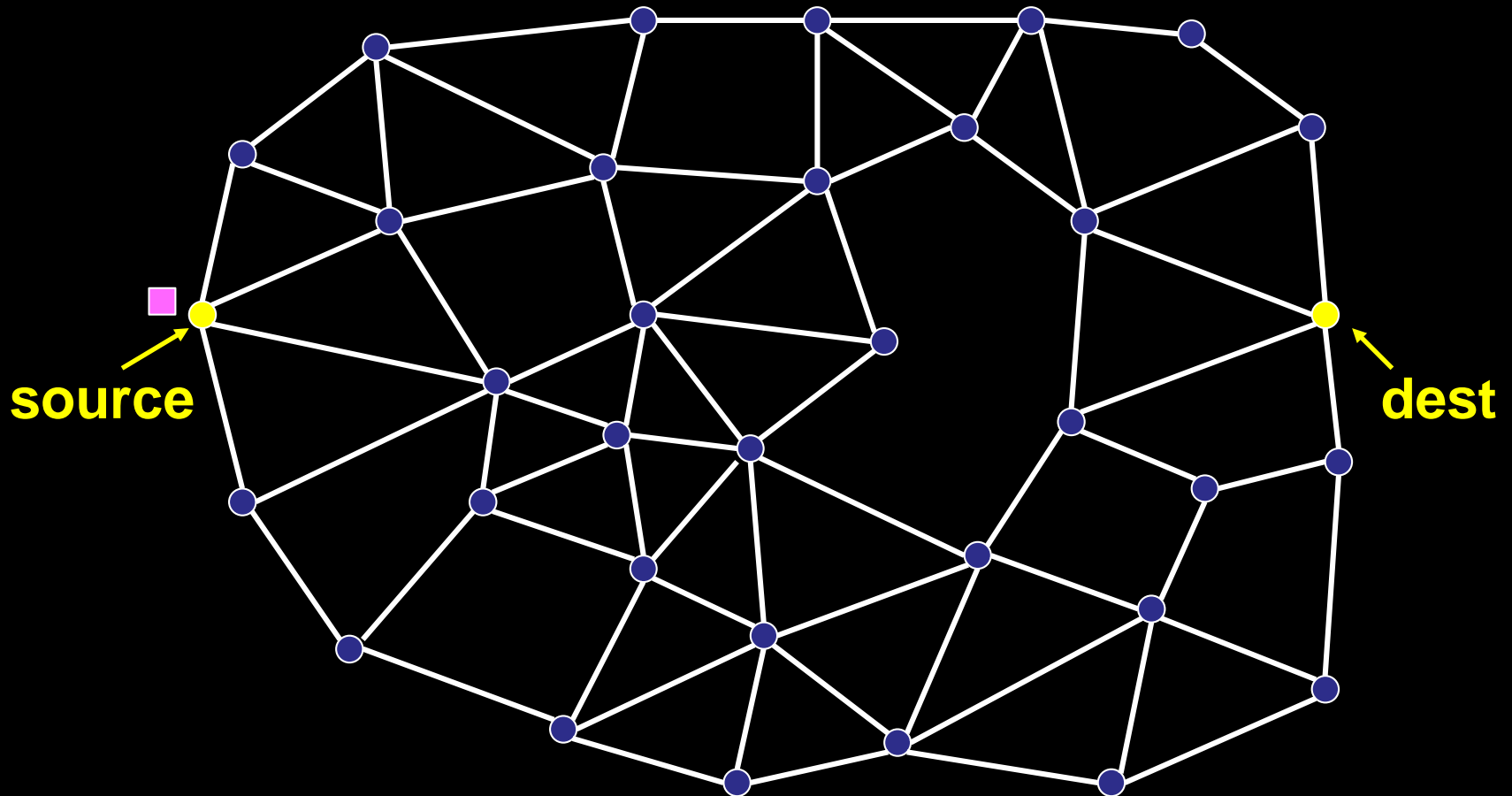
**Nodes have coordinates**

# Overview: Geographic Routing



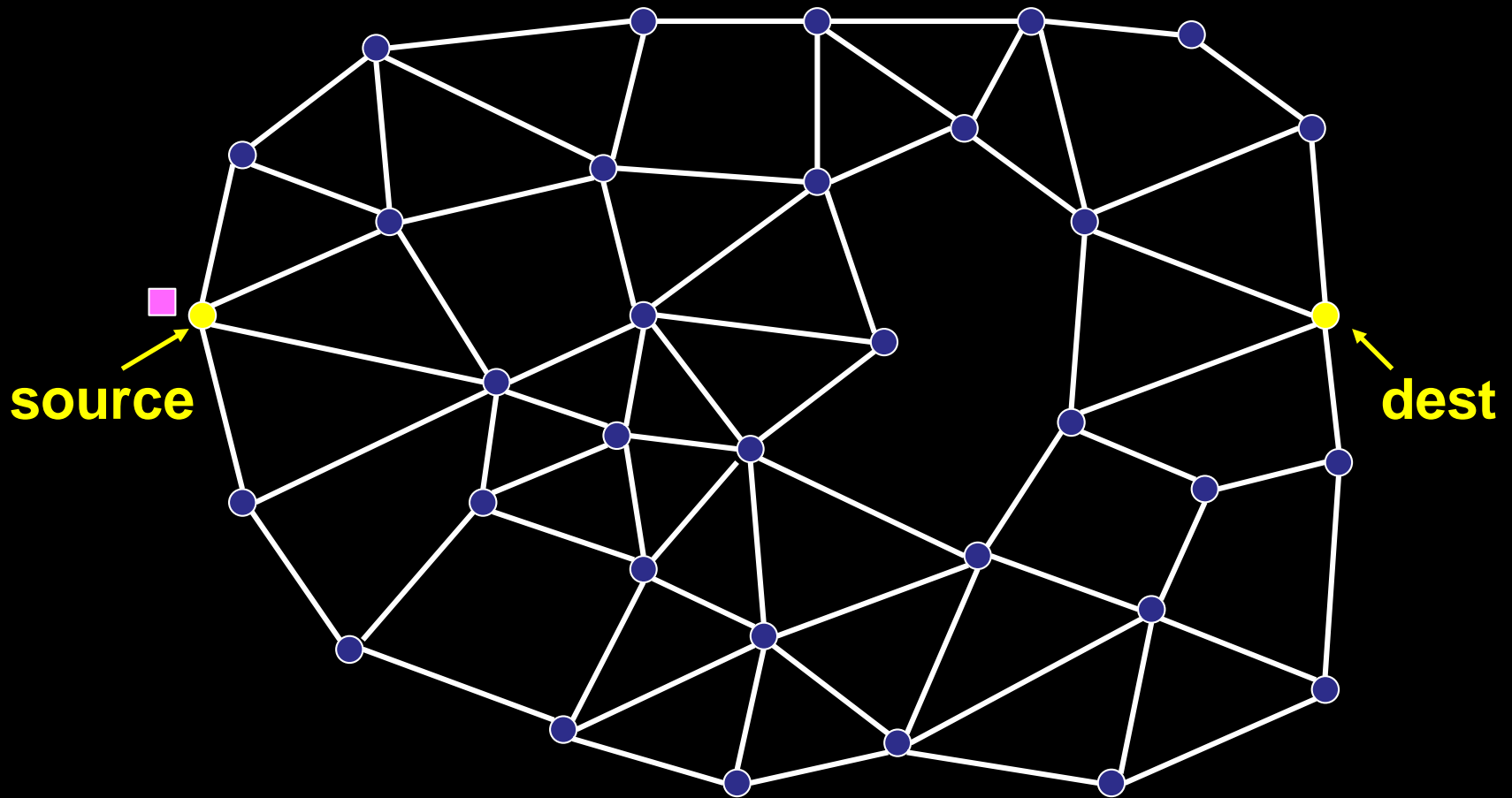
**Nodes have coordinates**

# Overview: Geographic Routing



**Packet contains coordinates of destination**

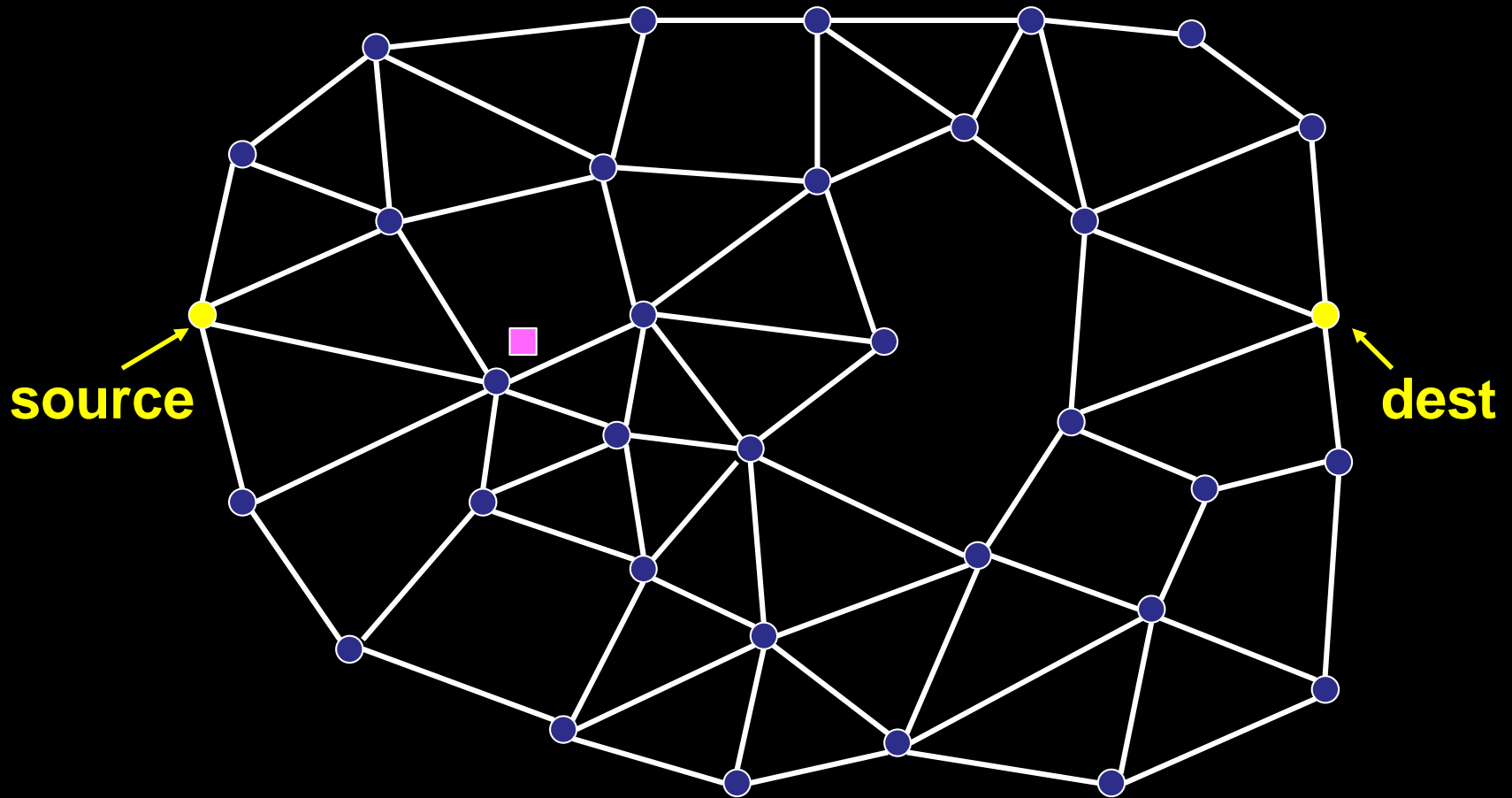
# Overview: Geographic Routing



**Greedy forwarding!**

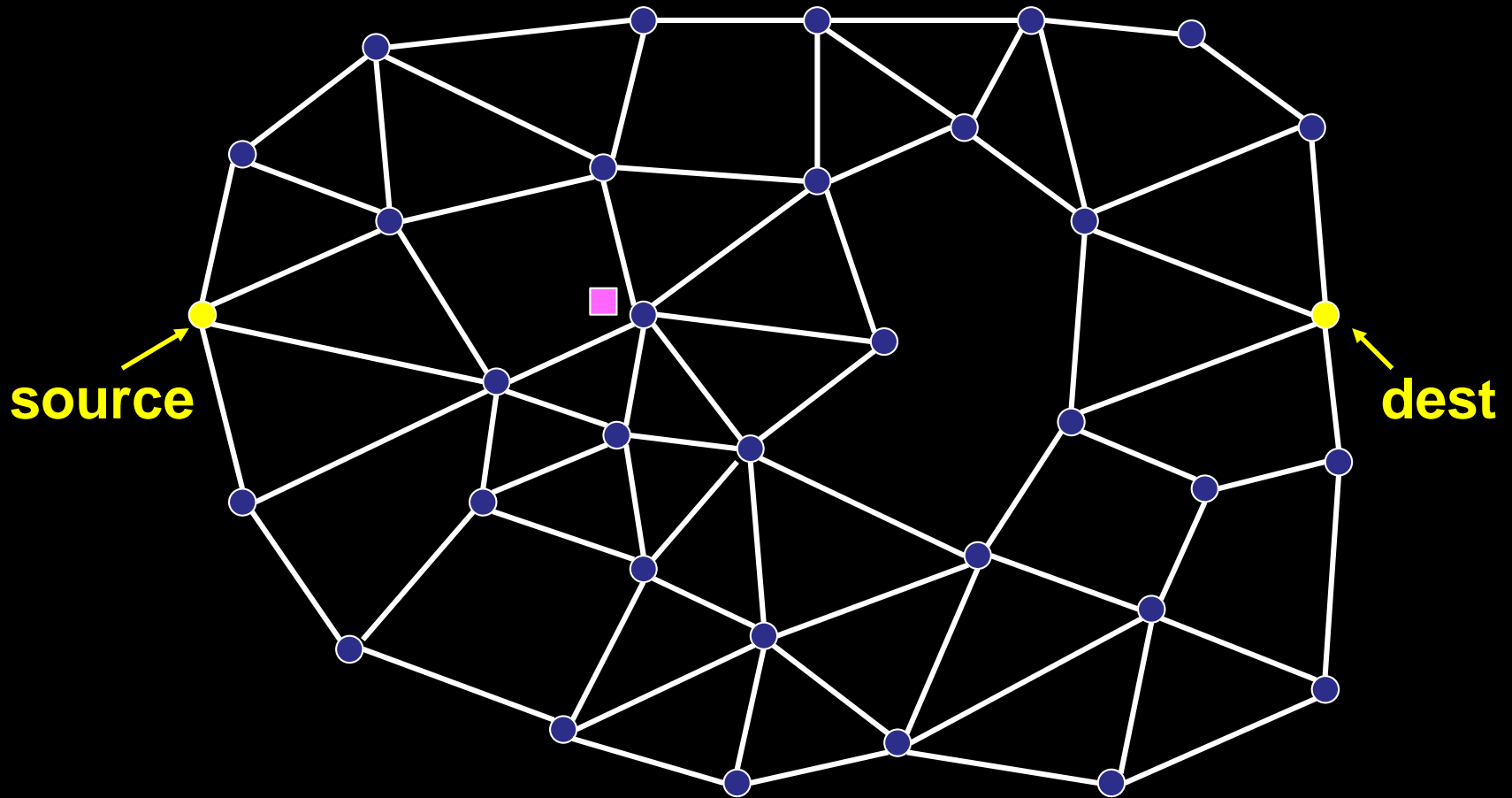


# Overview: Geographic Routing



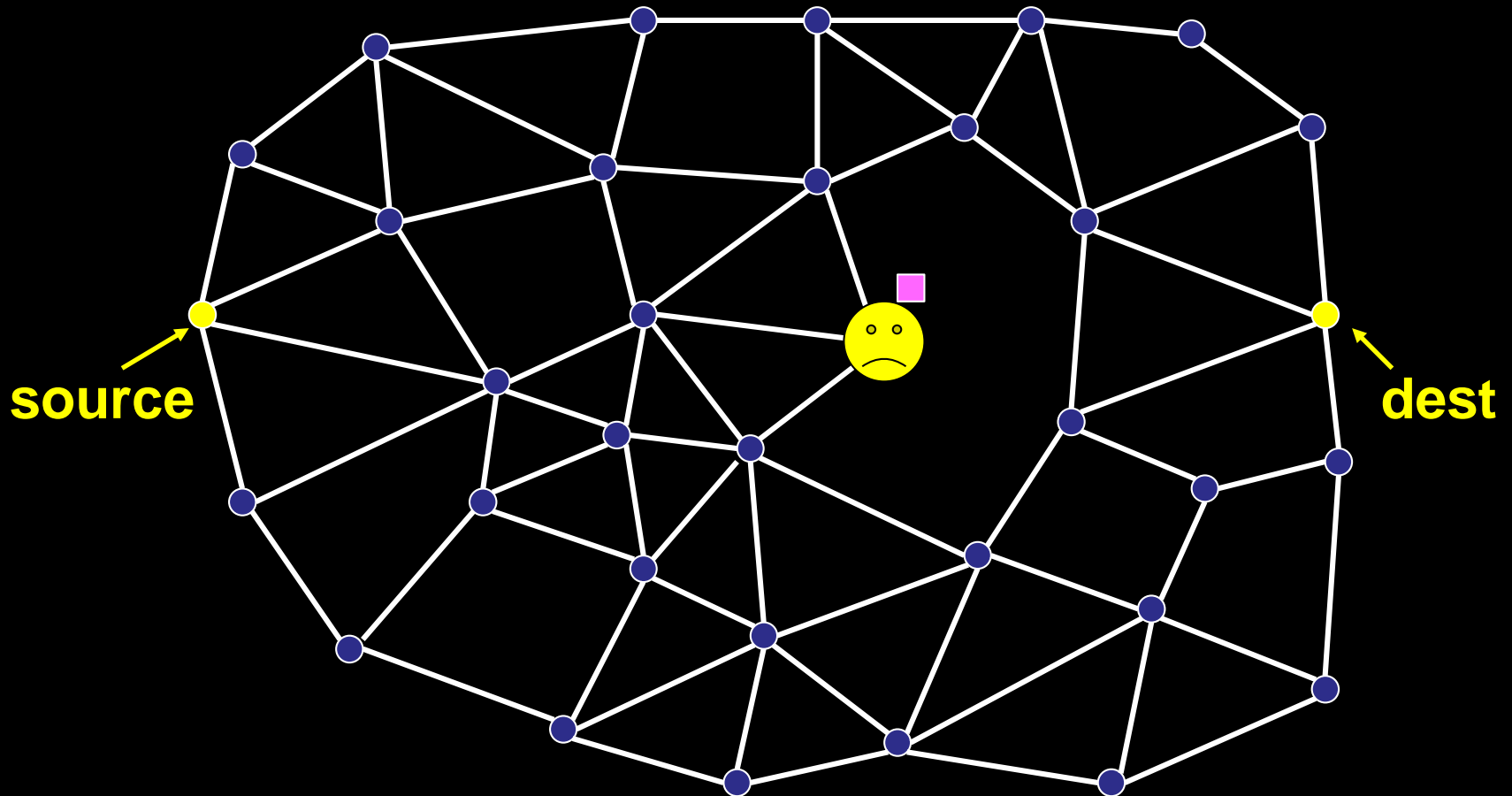
**Greedy forwarding!**

# Overview: Geographic Routing



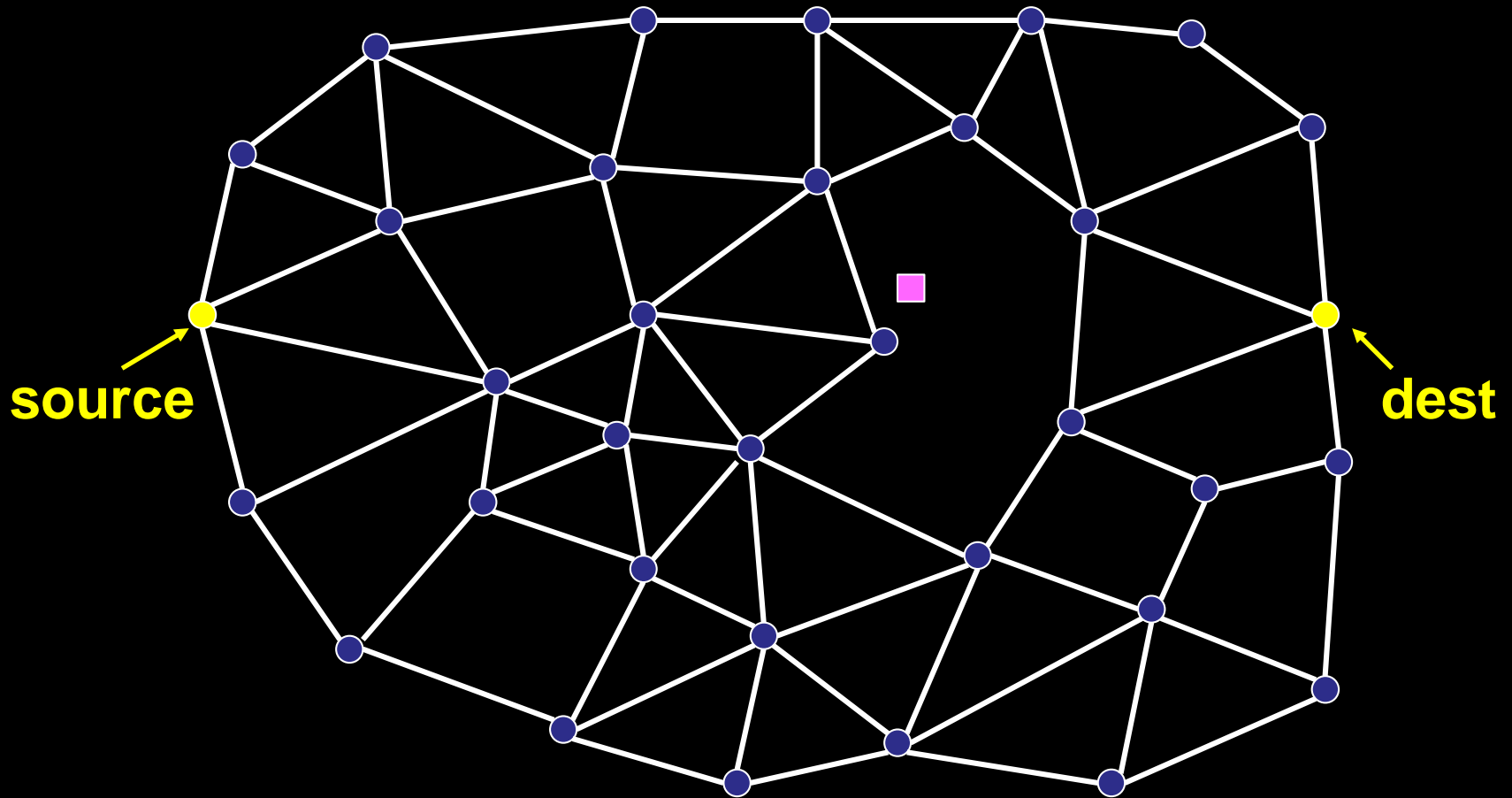
**Greedy forwarding!**

# Overview: Geographic Routing



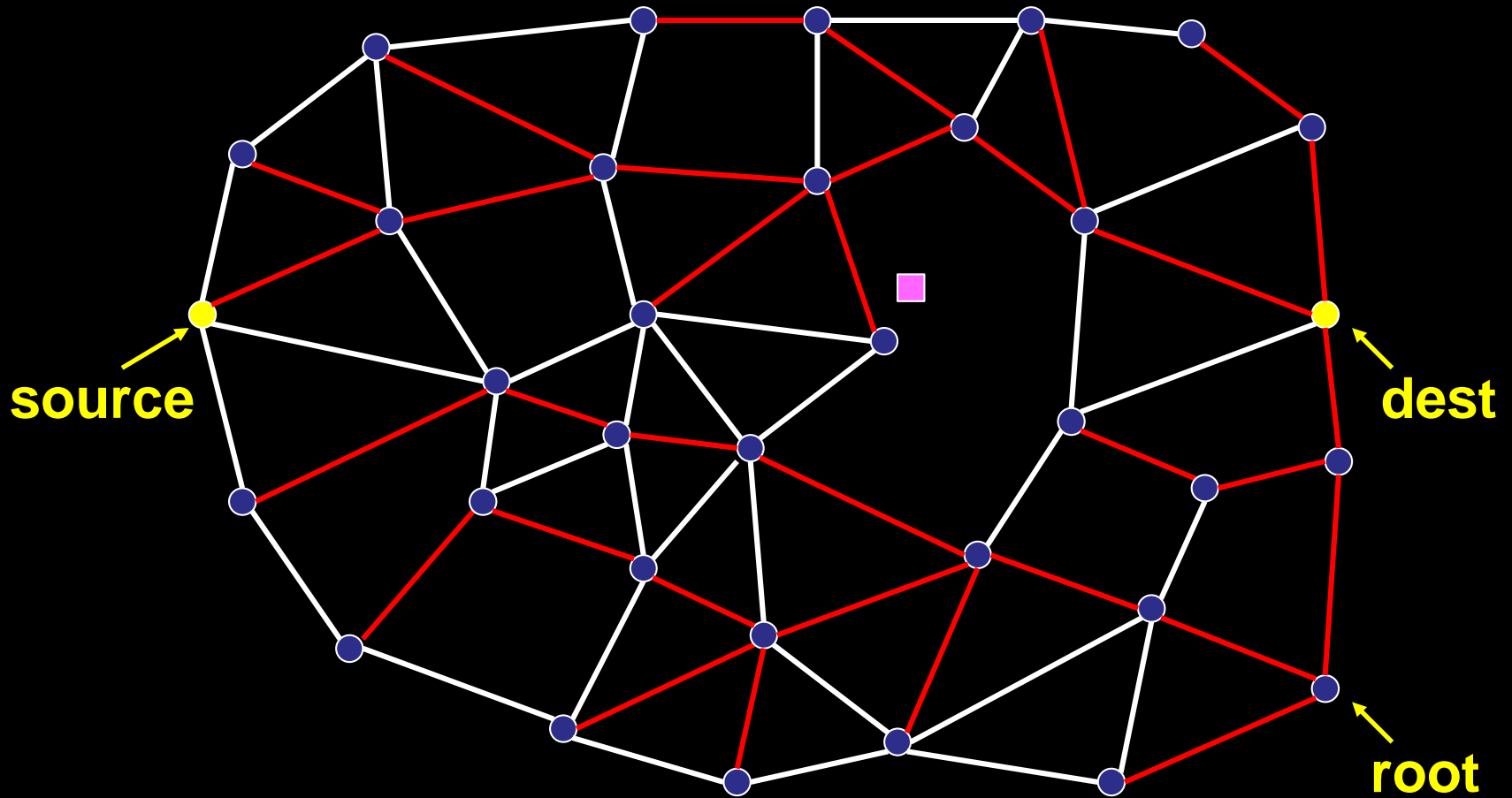
**Dead end! (local minima)**

# Overview: GDSTR



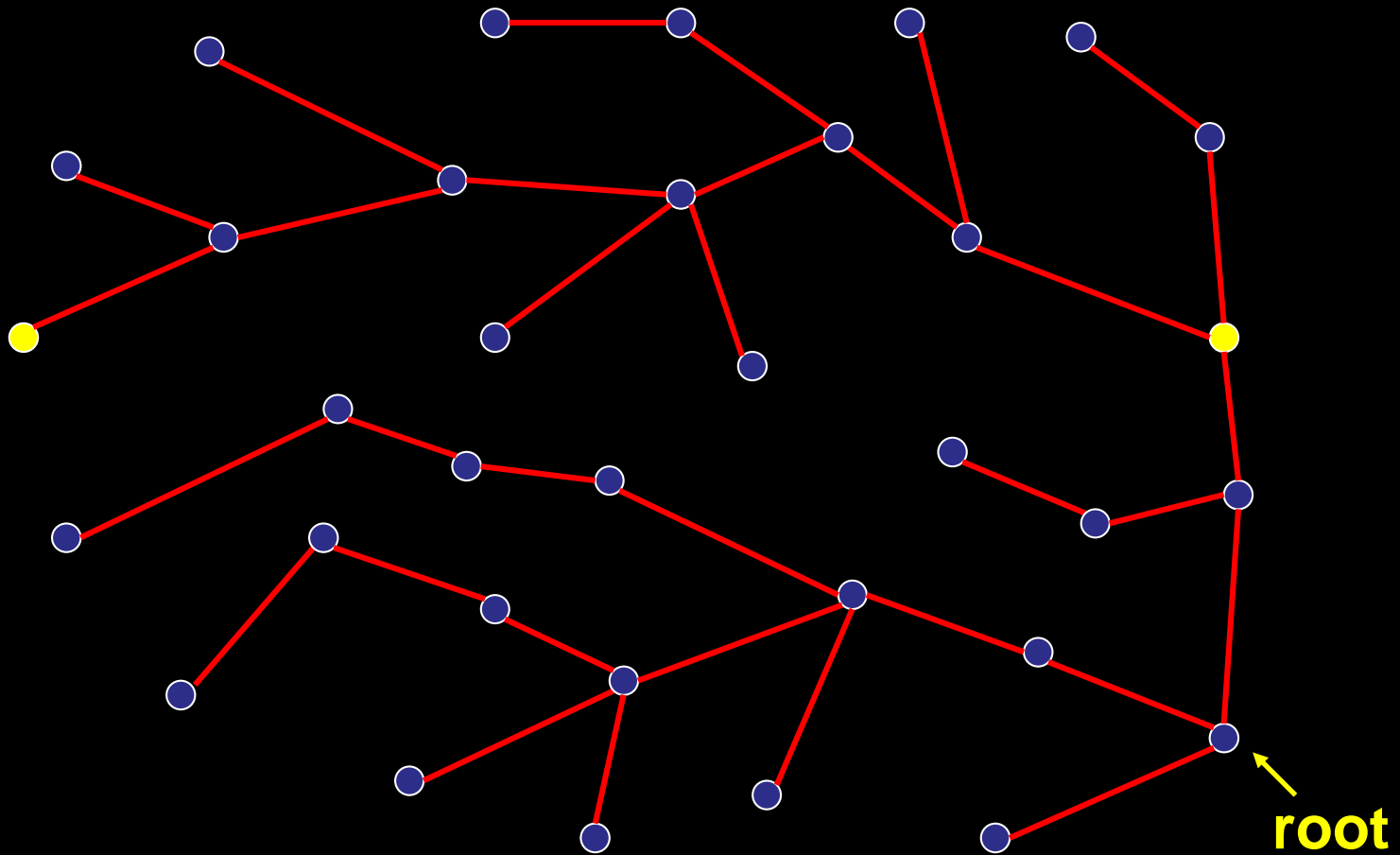
**Distributed Spanning Tree**

# Overview: GDSTR



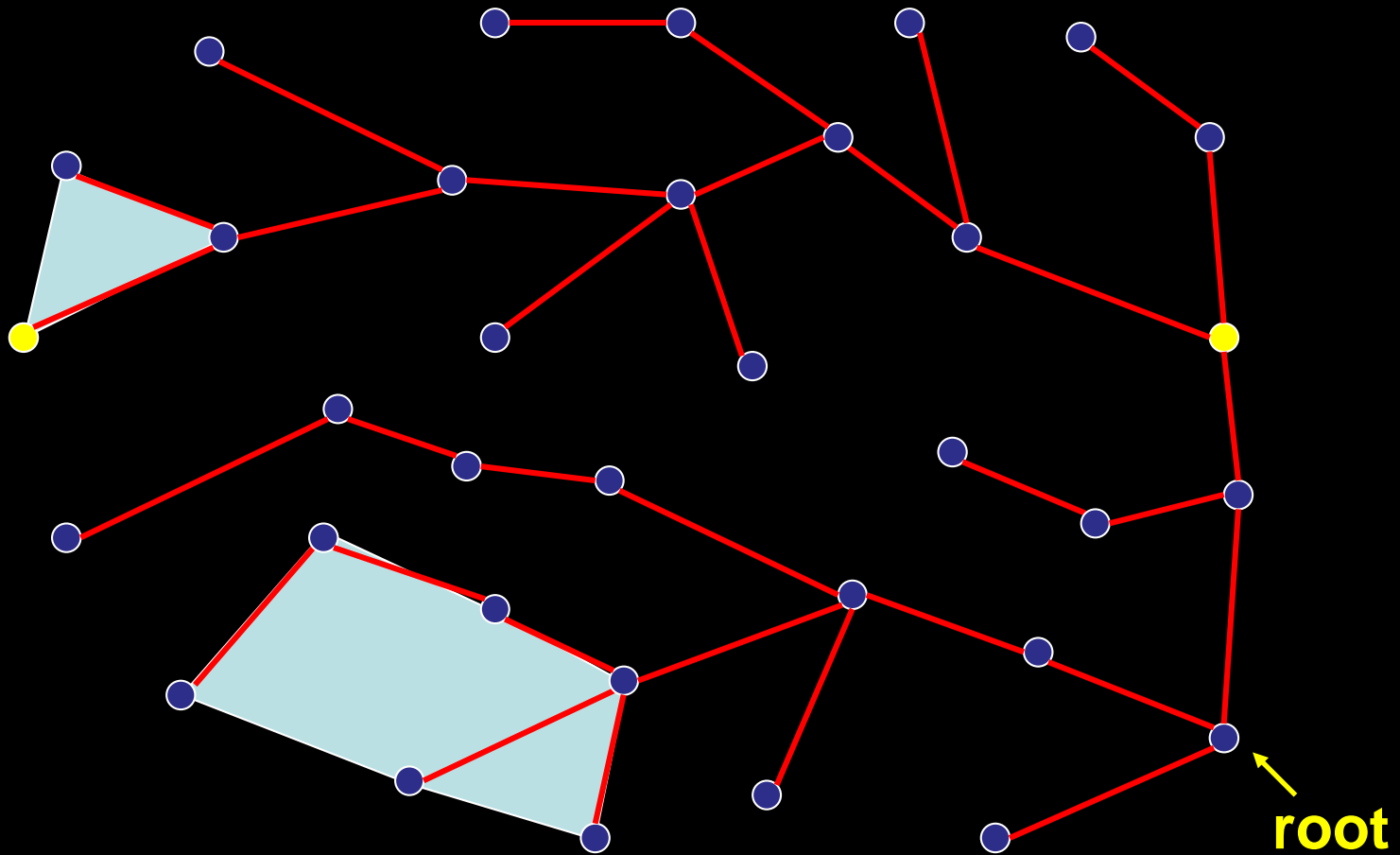
**Distributed Spanning Tree**

# Overview: GDSTR



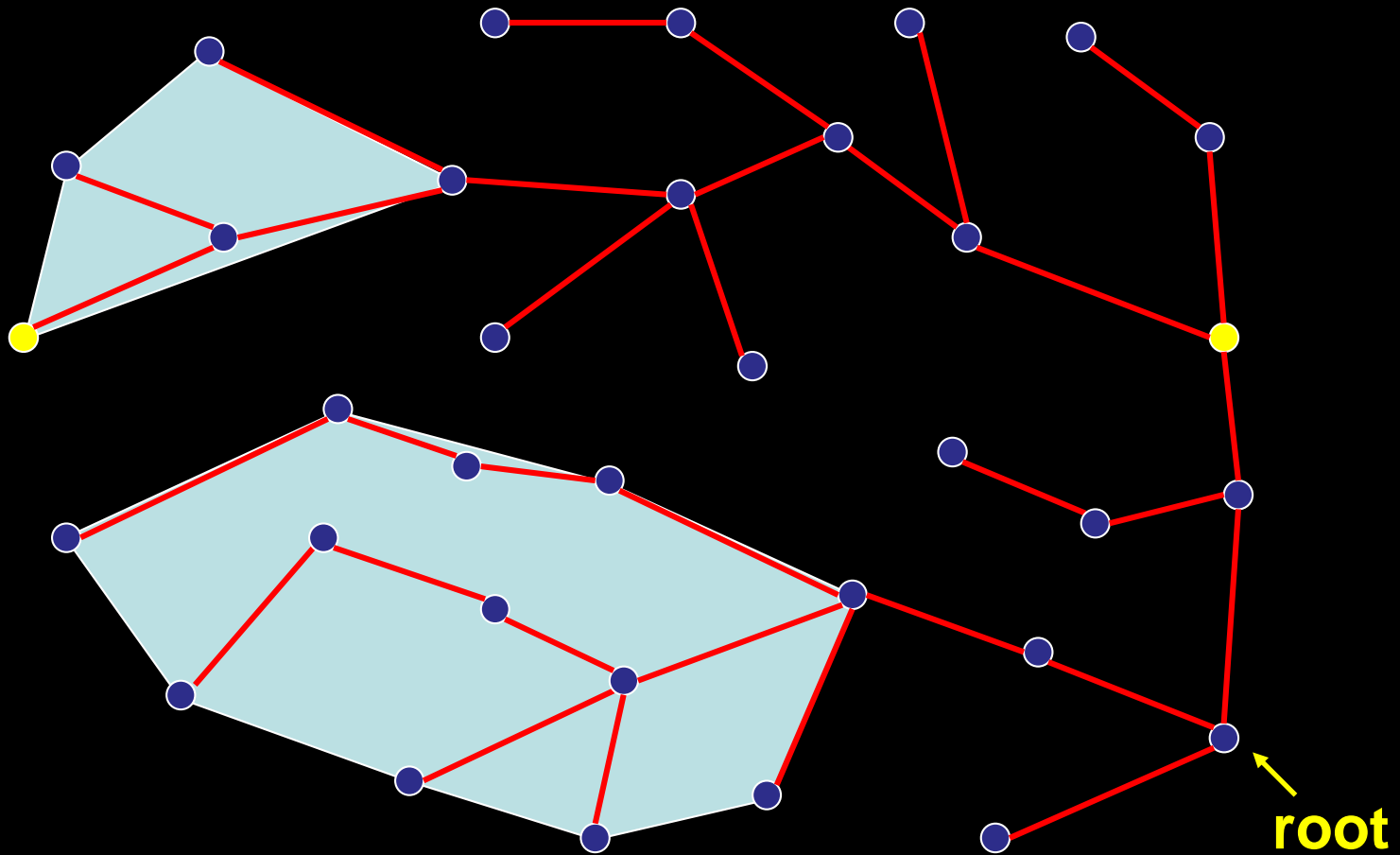
**Aggregate coordinates with convex hulls**

# Overview: GDSTR



Aggregate coordinates with convex hulls

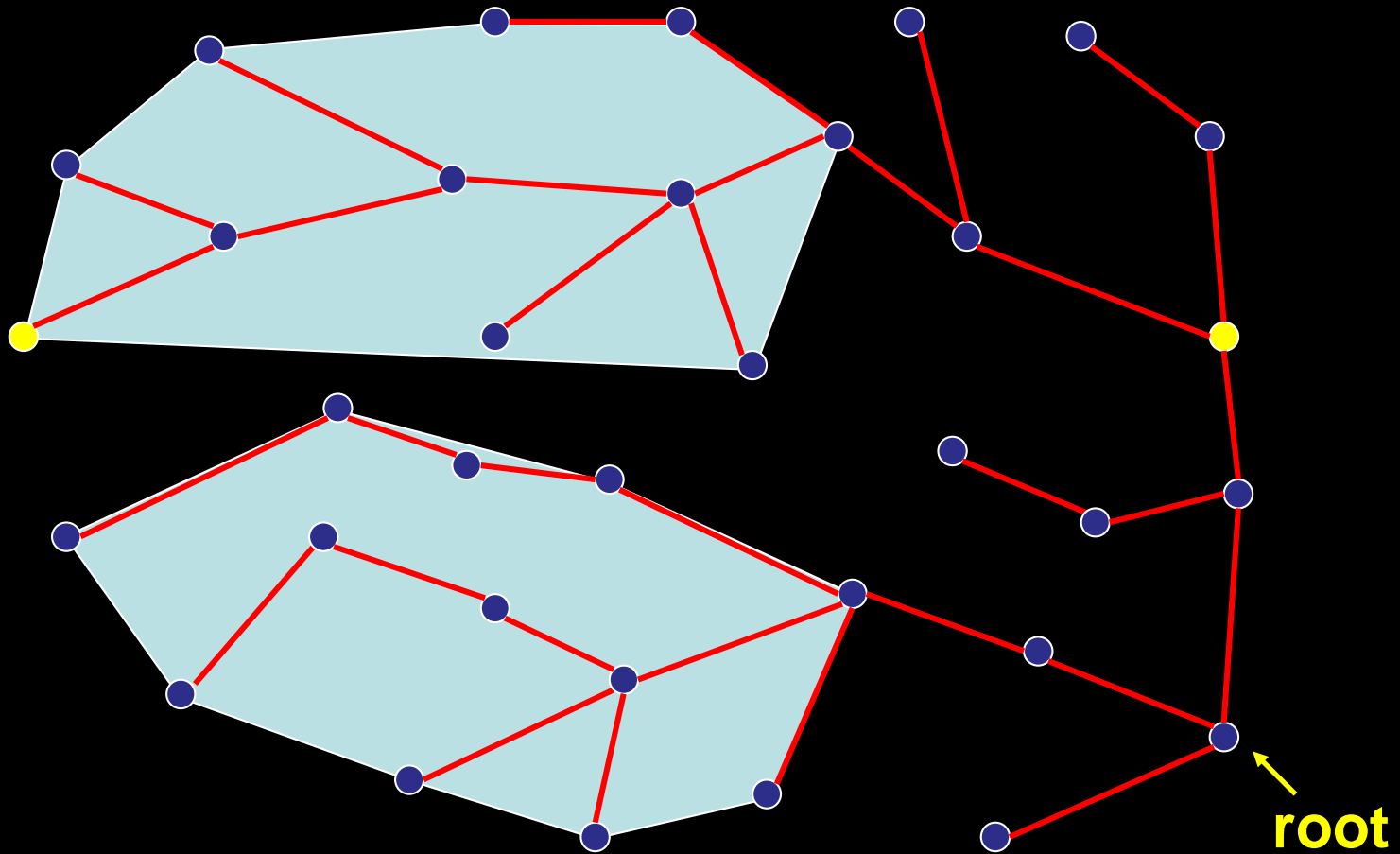
# Overview: GDSTR



Aggregate coordinates with convex hulls

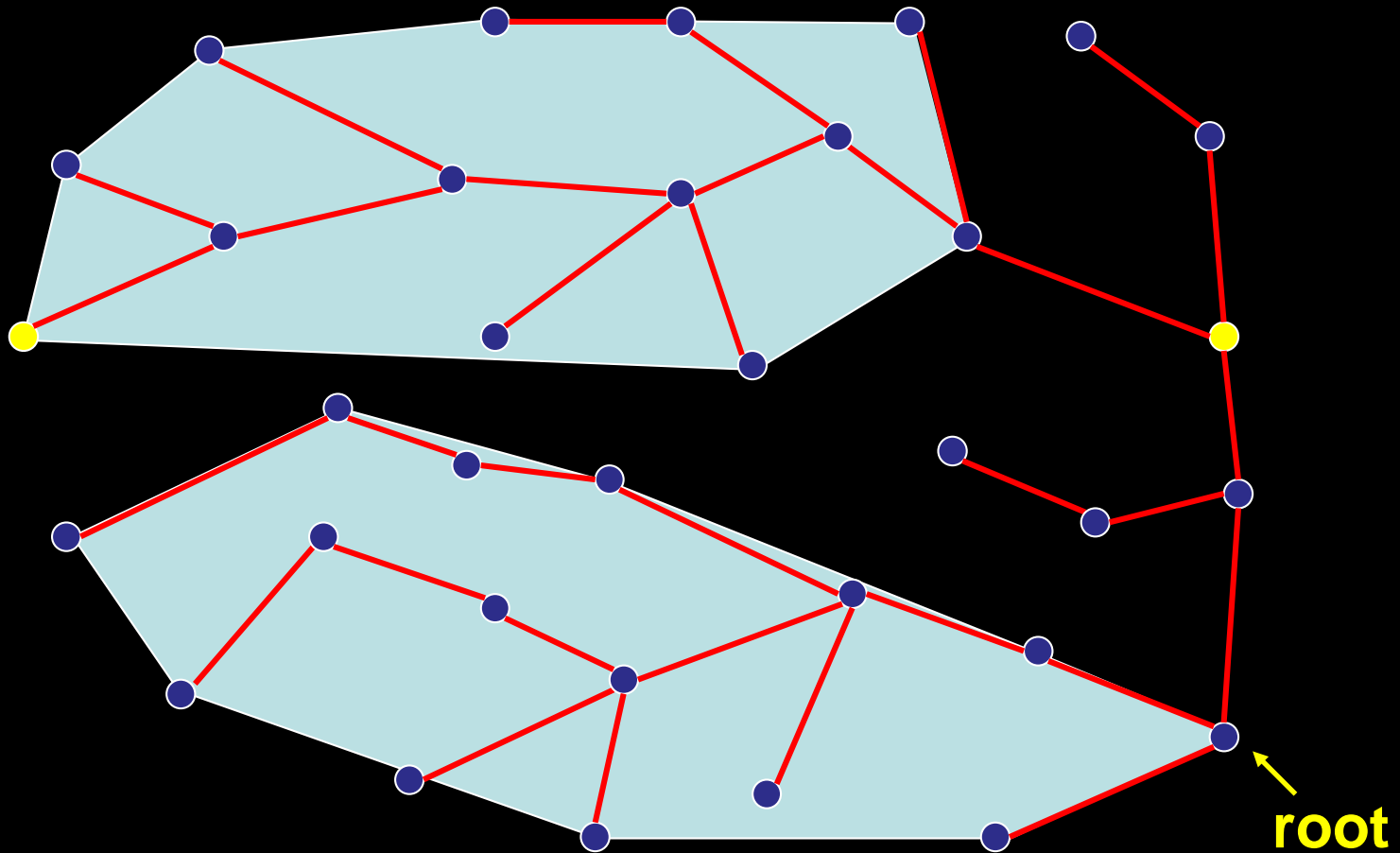


# Overview: GDSTR



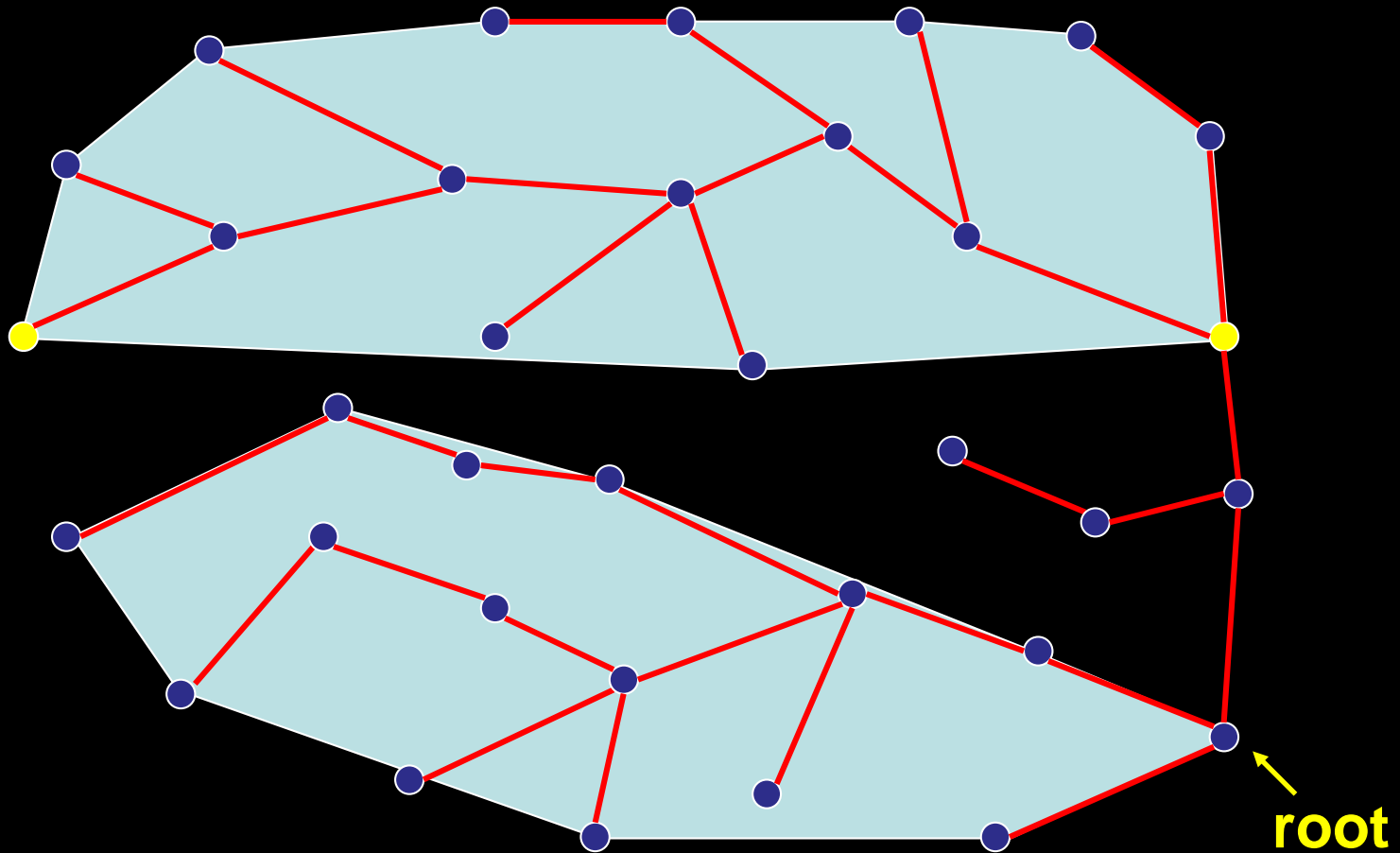
Aggregate coordinates with convex hulls

# Overview: GDSTR



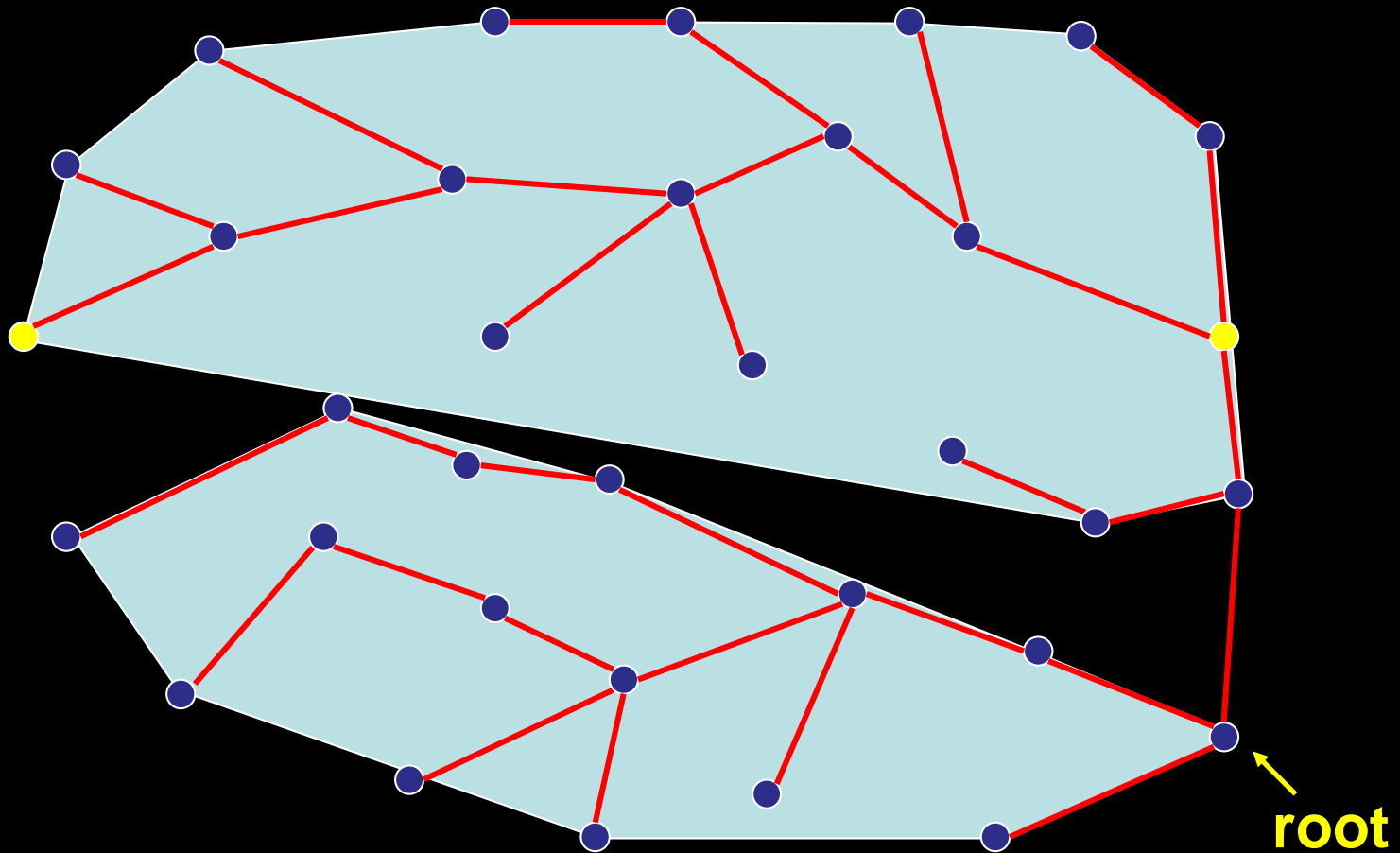
**Aggregate coordinates with convex hulls**

# Overview: GDSTR



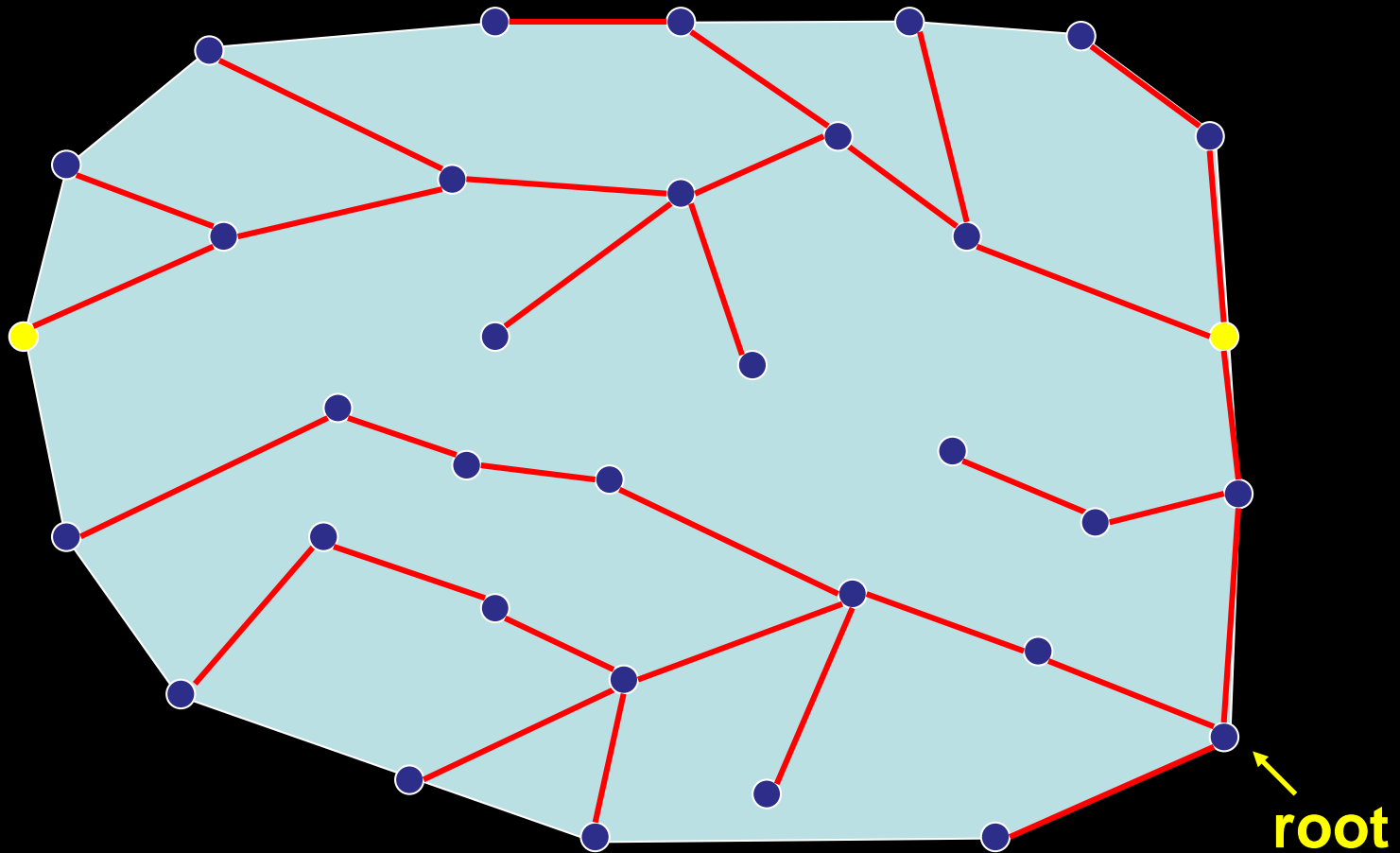
**Aggregate coordinates with convex hulls**

# Overview: GDSTR



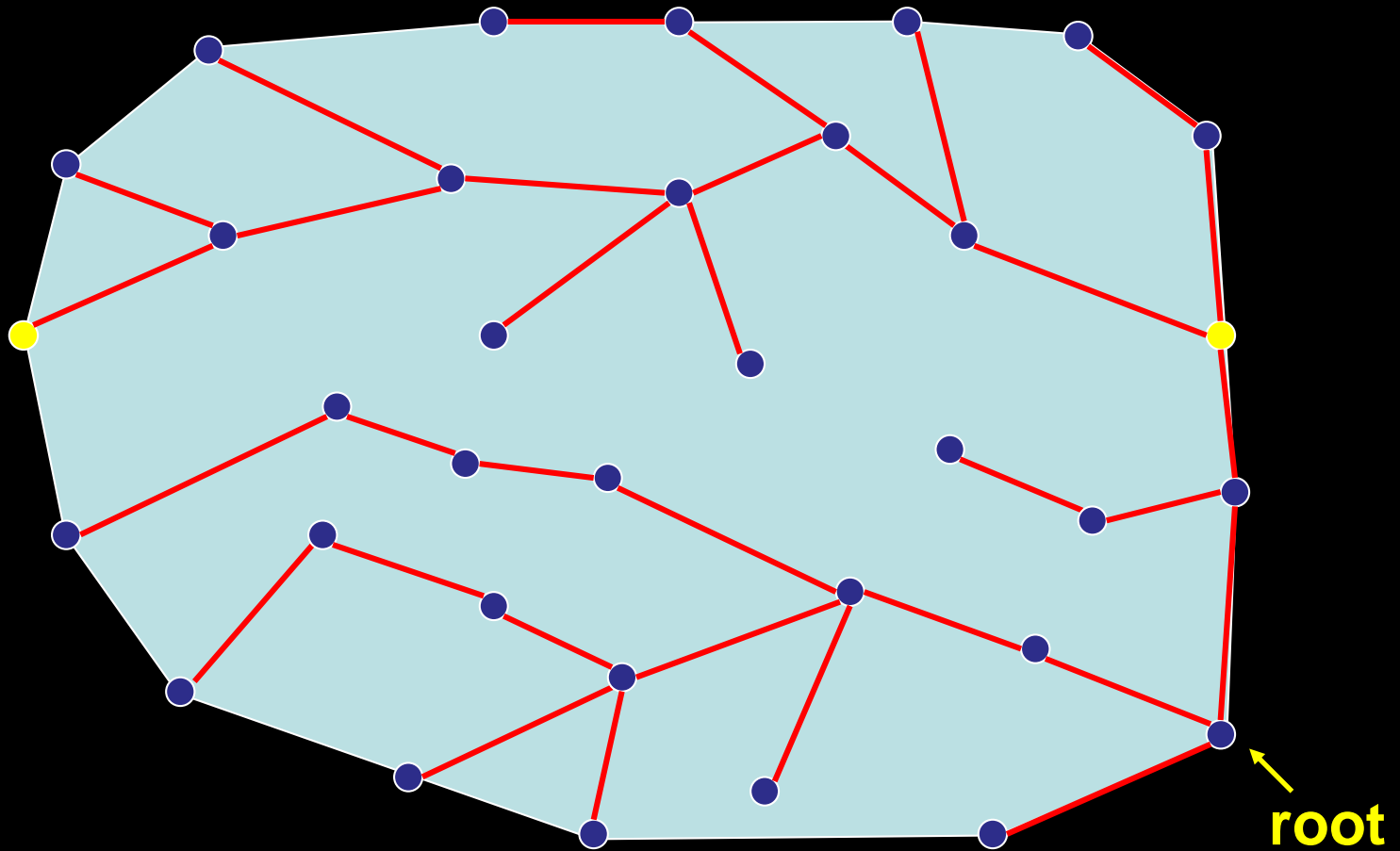
**Aggregate coordinates with convex hulls**

# Overview: GDSTR



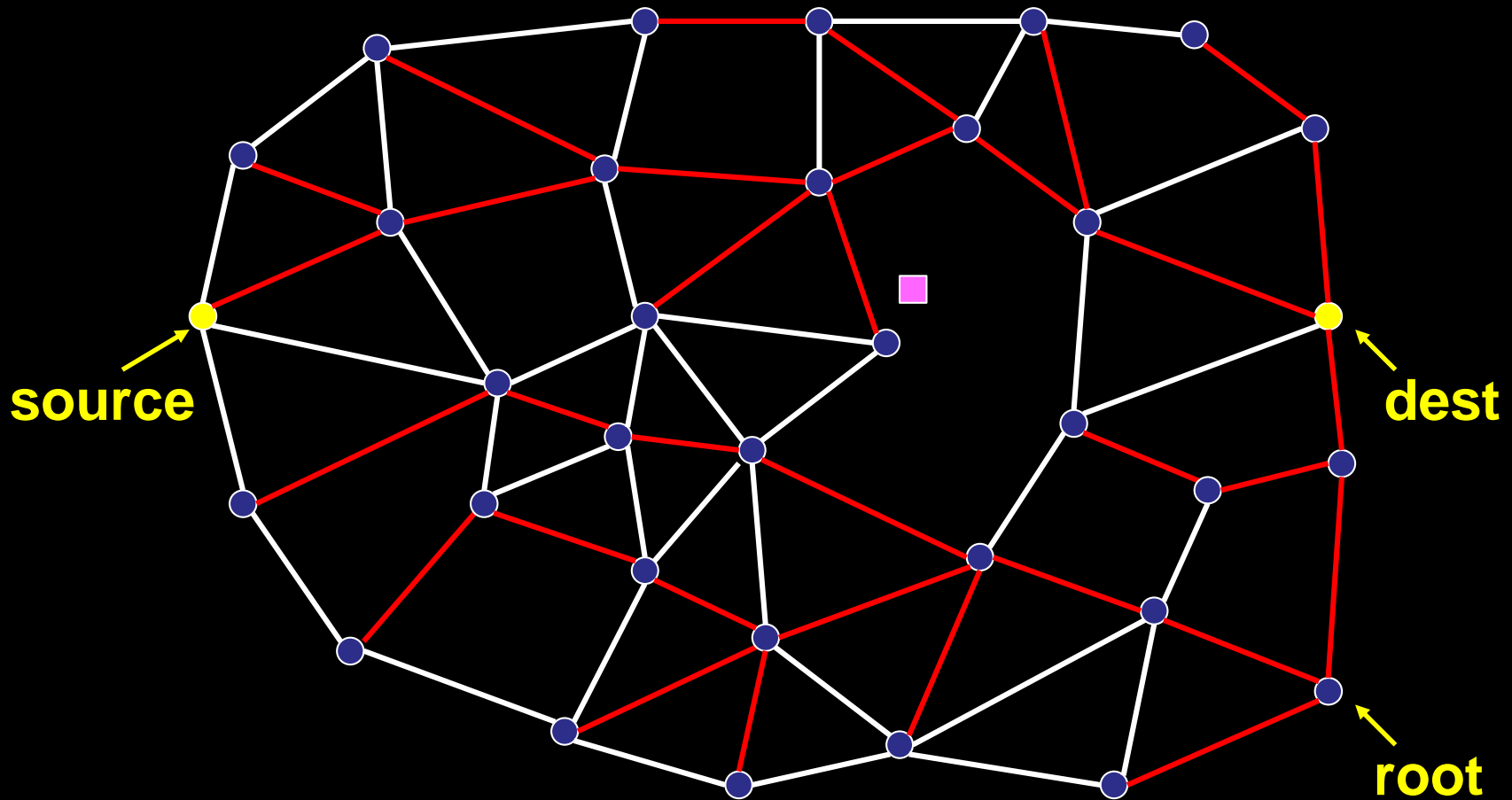
Aggregate coordinates with convex hulls

# Overview: GDSTR



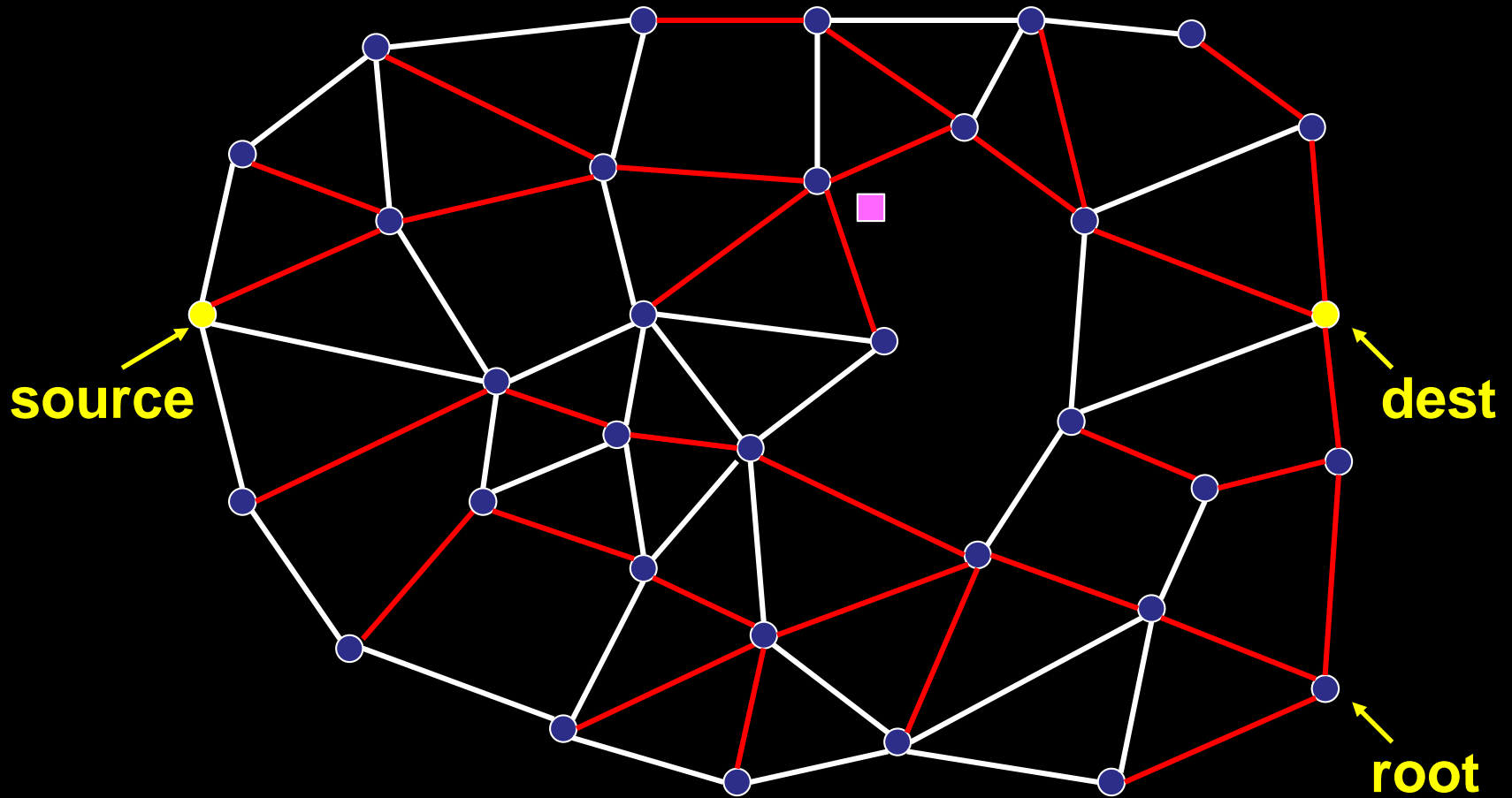
Hull Tree

# Overview: GDSTR



Remember minimum  $\Rightarrow$  tree traversal

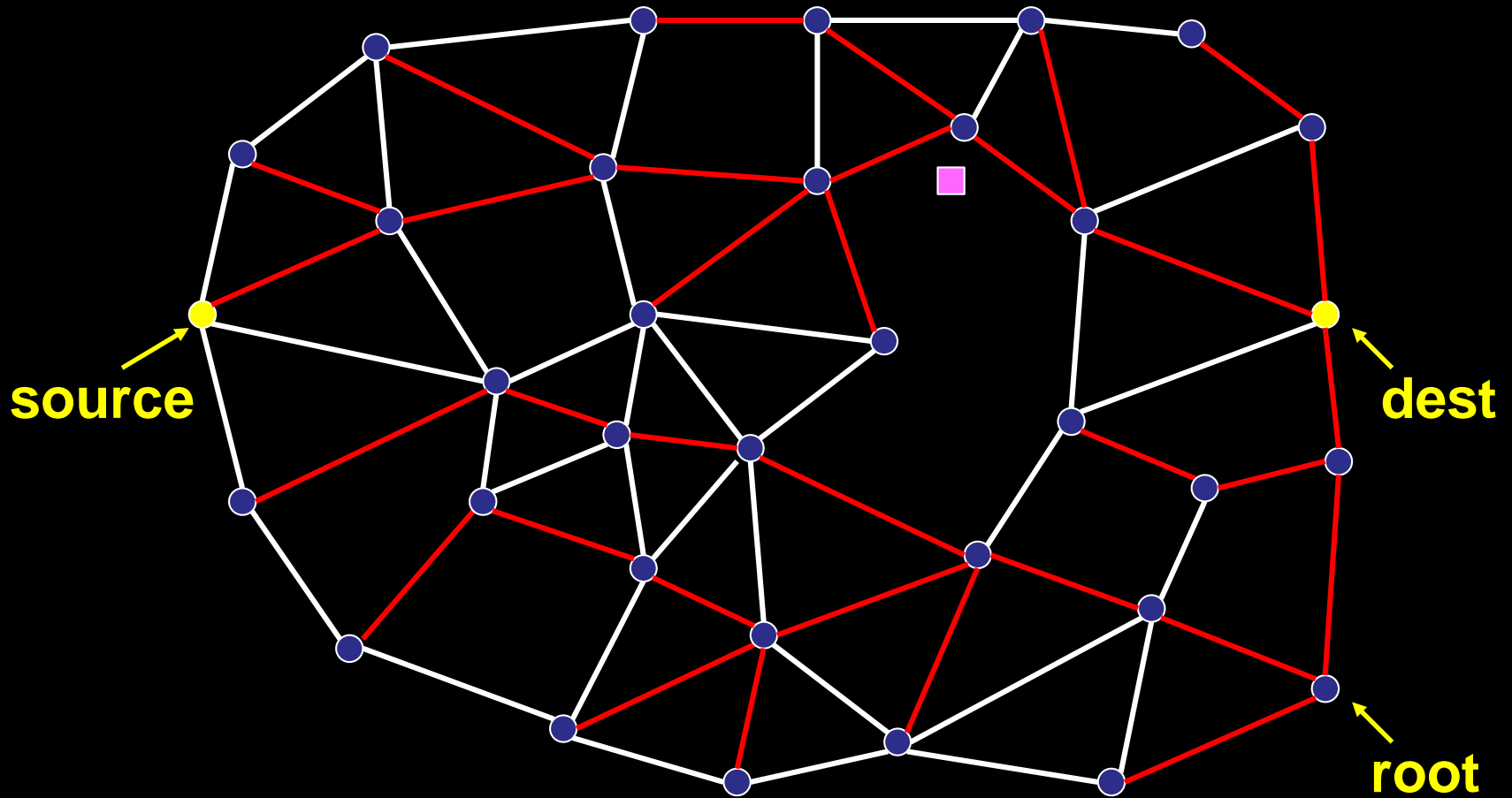
# Overview: GDSTR



**Tree Traversal**

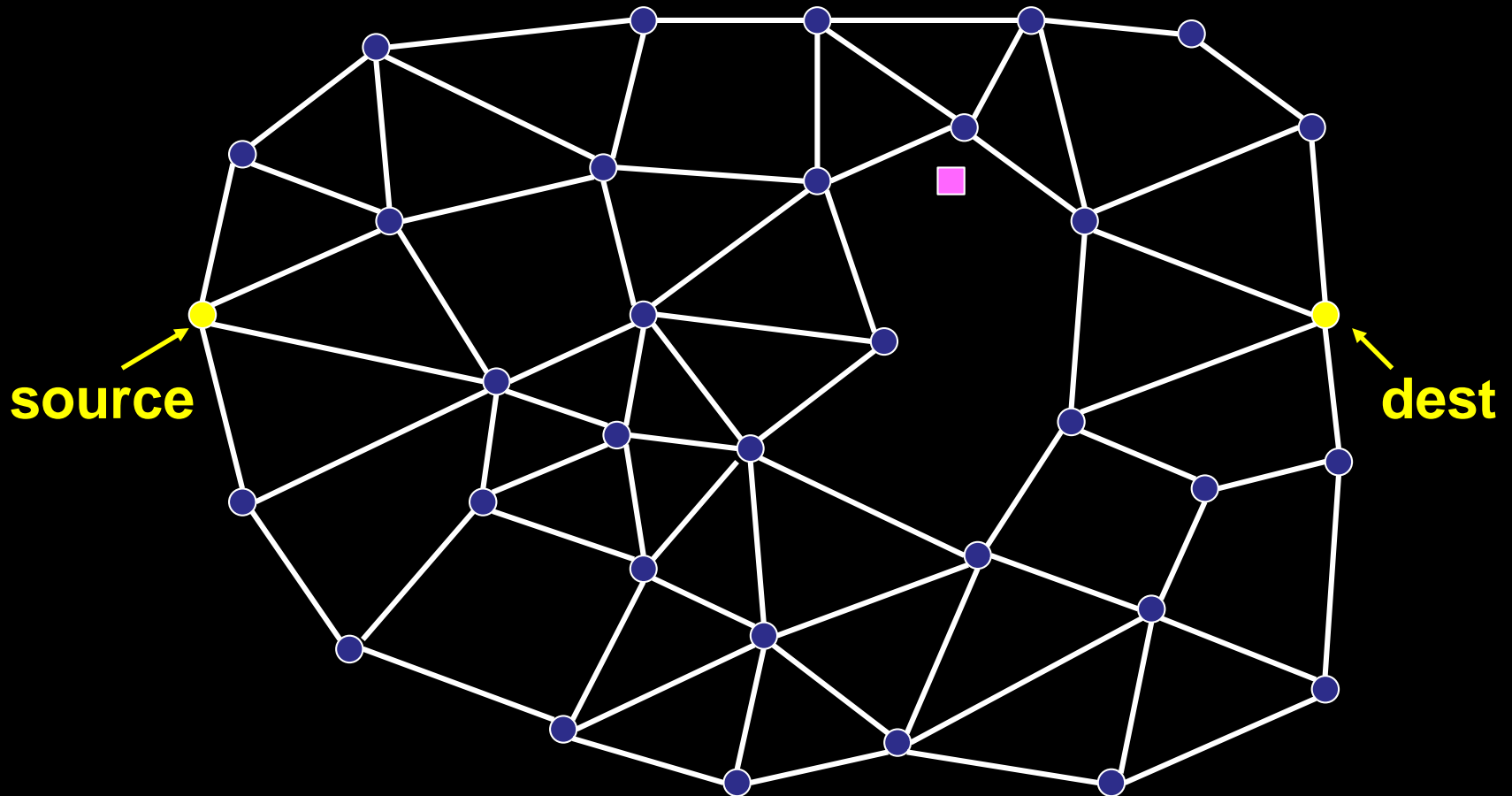


# Overview: GDSTR



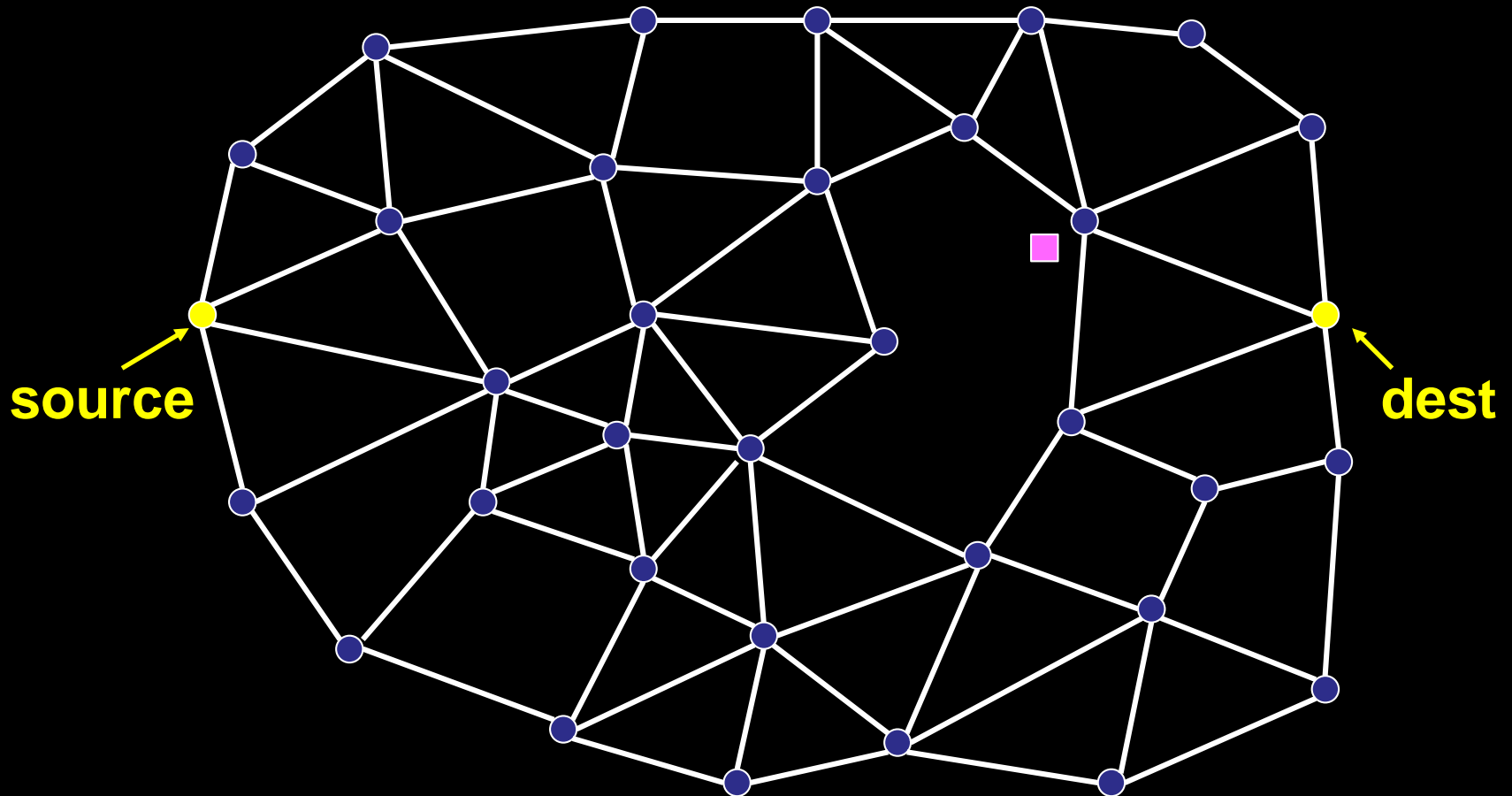
**Tree Traversal**

# Overview: GDSTR



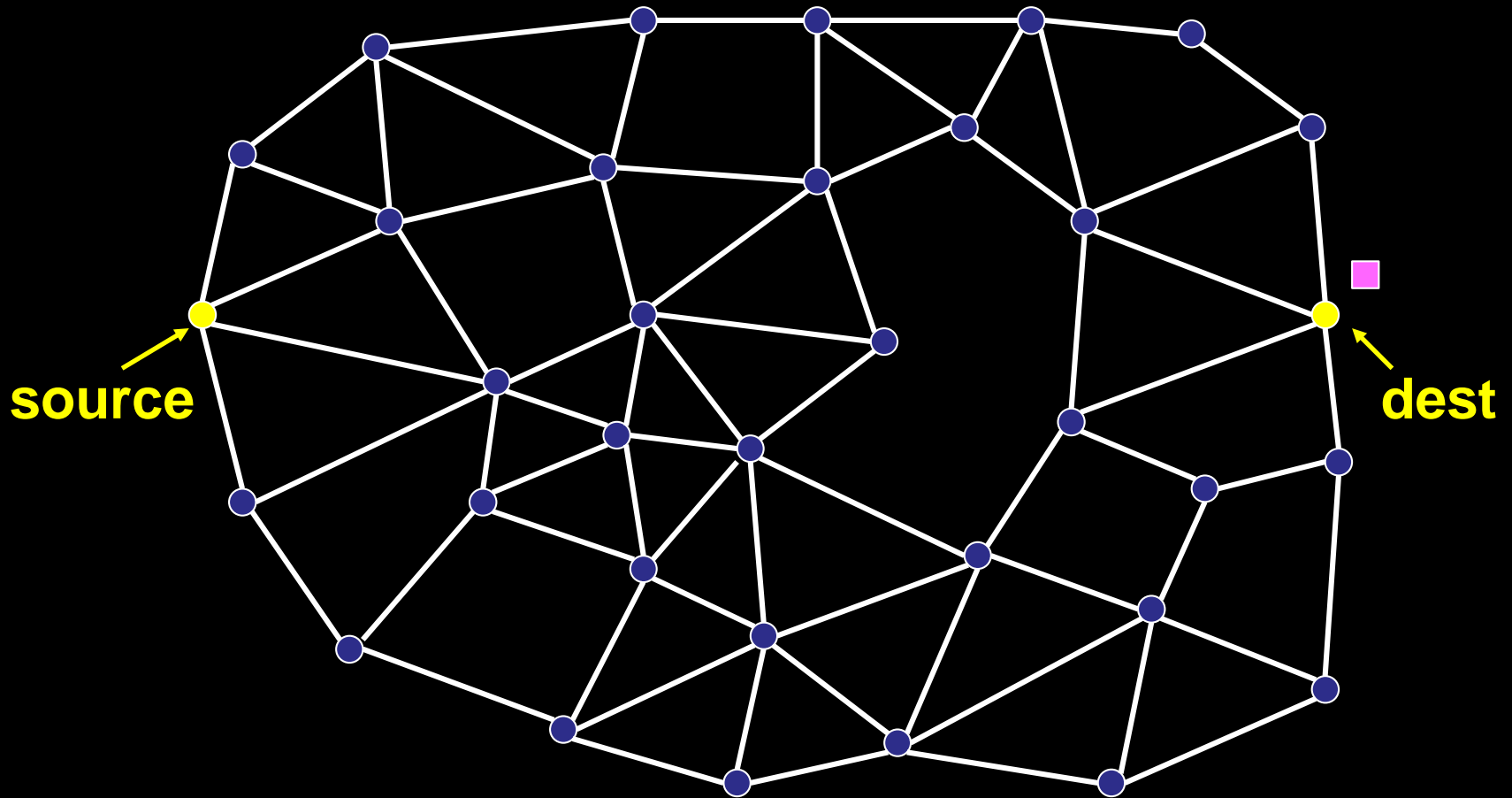
**Back to Greedy Forwarding!**

# Overview: GDSTR



**Back to Greedy Forwarding!**

# Overview: GDSTR

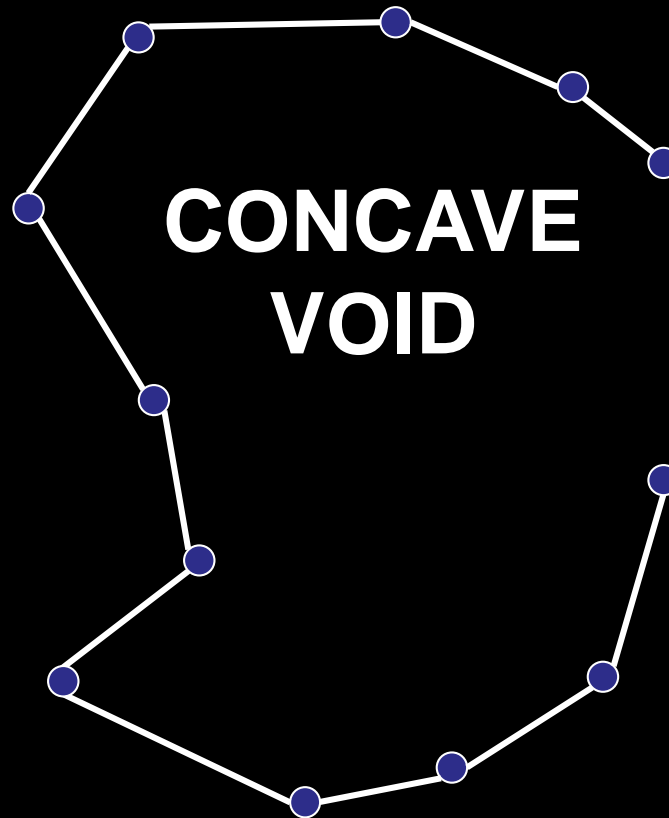


**Done!!**

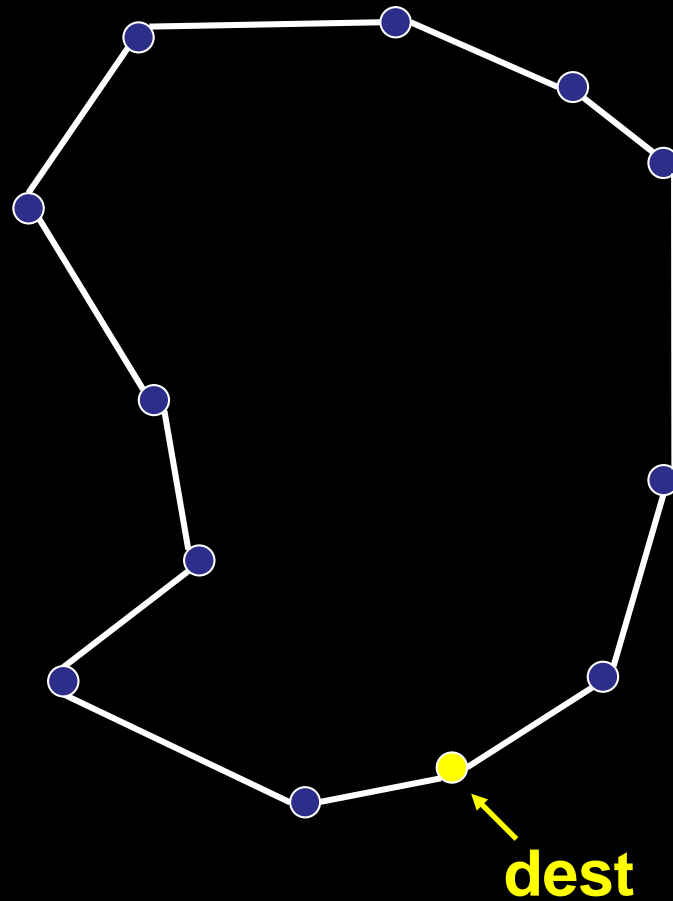
# Why do hull trees work?

- Used only to escape from local minimum
- Cheap to build –  $O(\log n)$

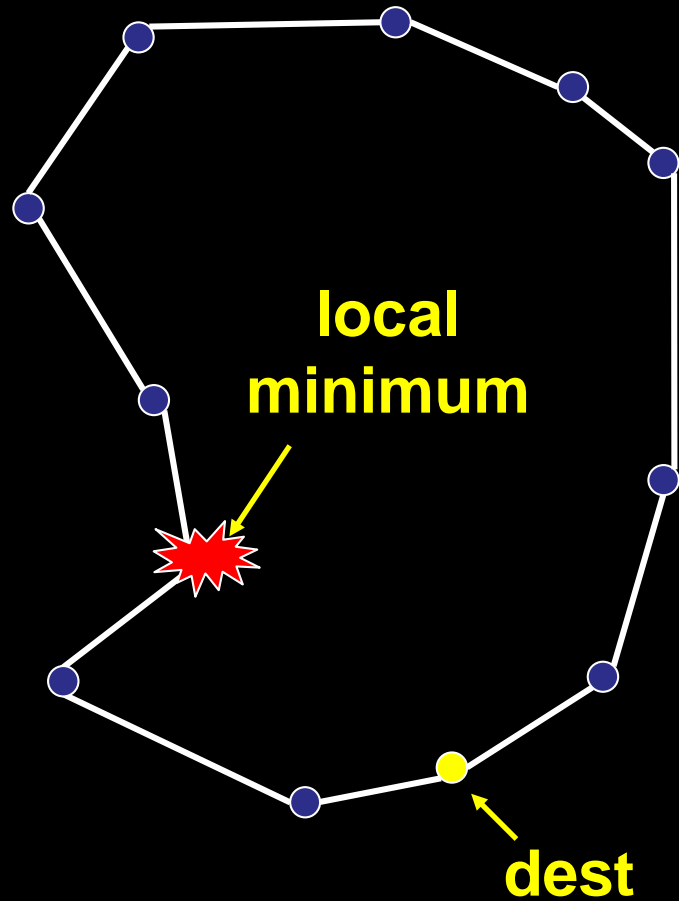
# Caveat



# Caveat

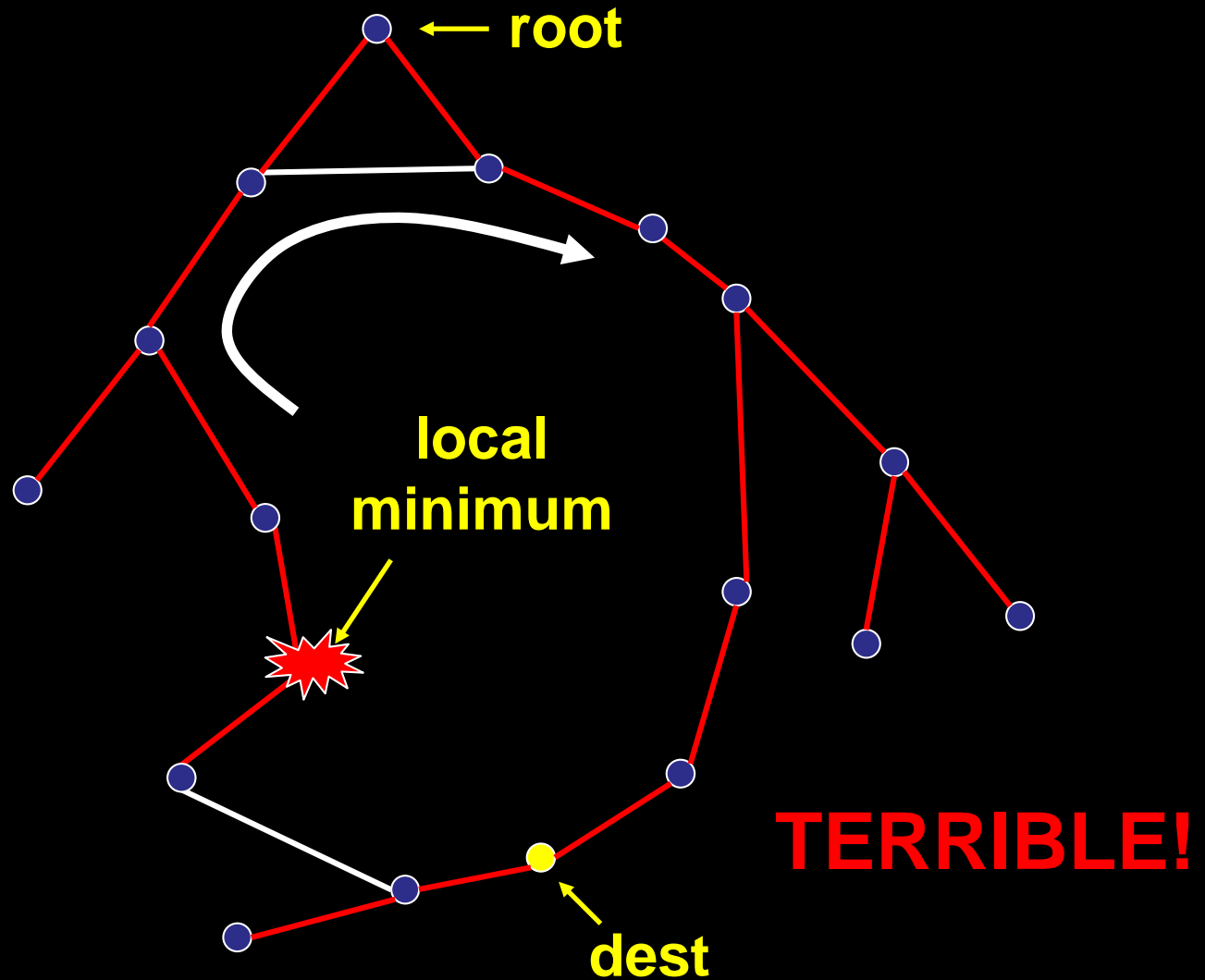


# Caveat



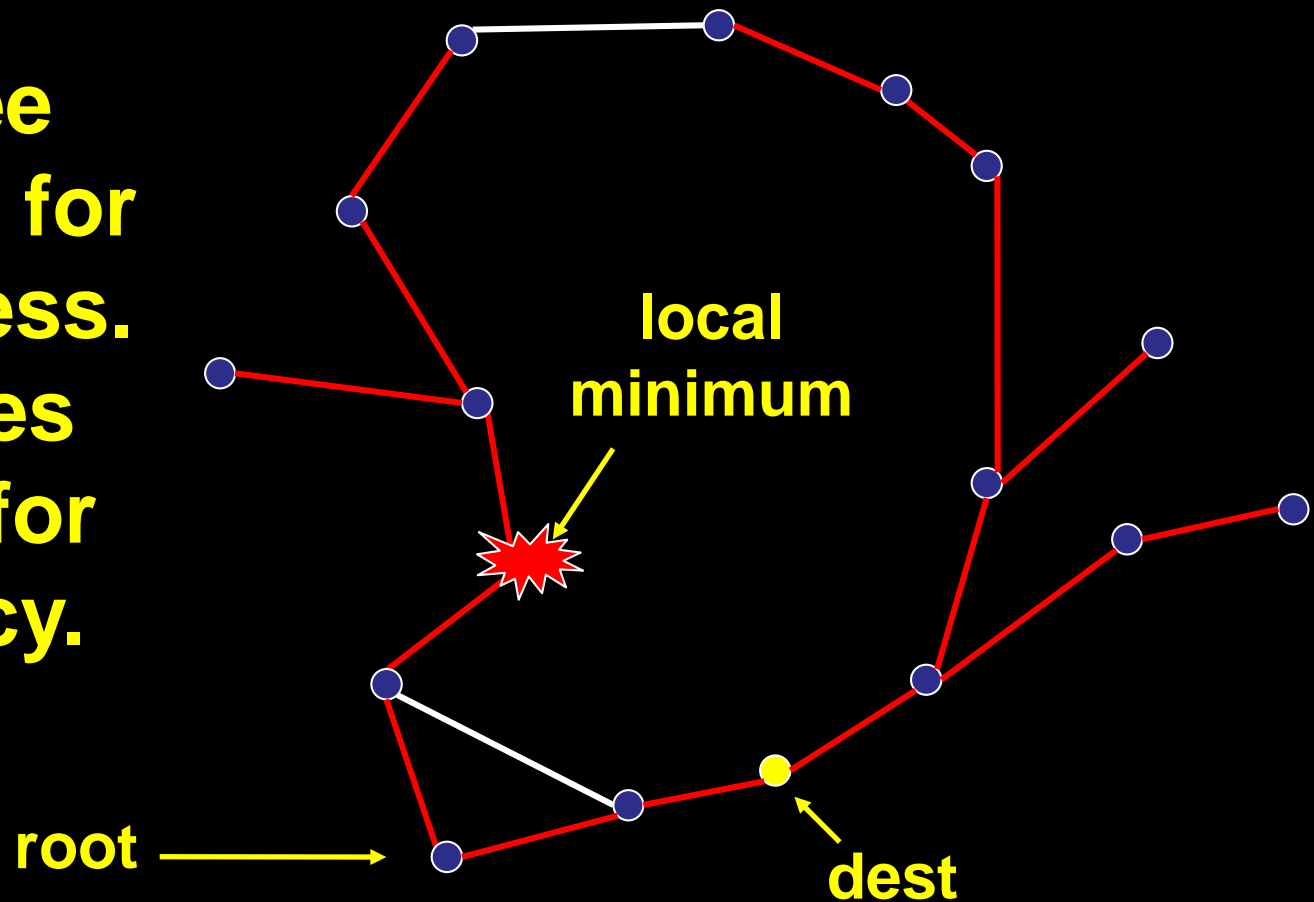


# Caveat



# Need TWO hull trees rooted at opposite ends

**One tree  
sufficient for  
correctness.  
Two trees  
needed for  
efficiency.**



Our Approach

Extend GDSTR  
to 3D

Complications!

# Challenges

(Why is it hard in TinyOS?)

- TinyOS does not support dynamic memory allocation
- CC2420 radio supports up to 128 bytes in size and has a limited data rate
- Limited DRAM and flash memory
- Precision of floating point operations is limited

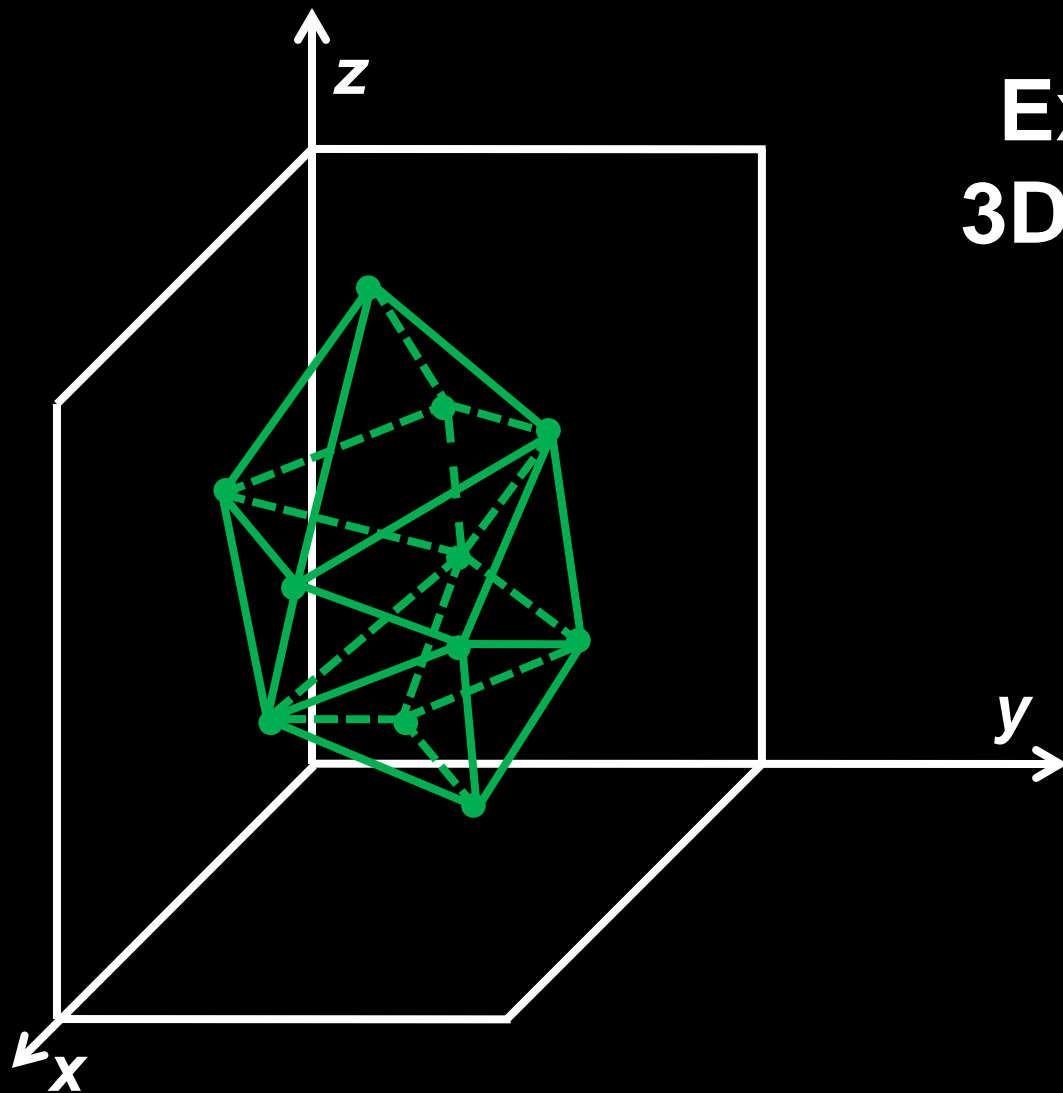
# Naïve Implementation of 3D Convex Hull

- Computations are costly
- Need to store auxiliary data structures for efficiency  $\Rightarrow$  storage costly
- Messages too big

# Key ideas

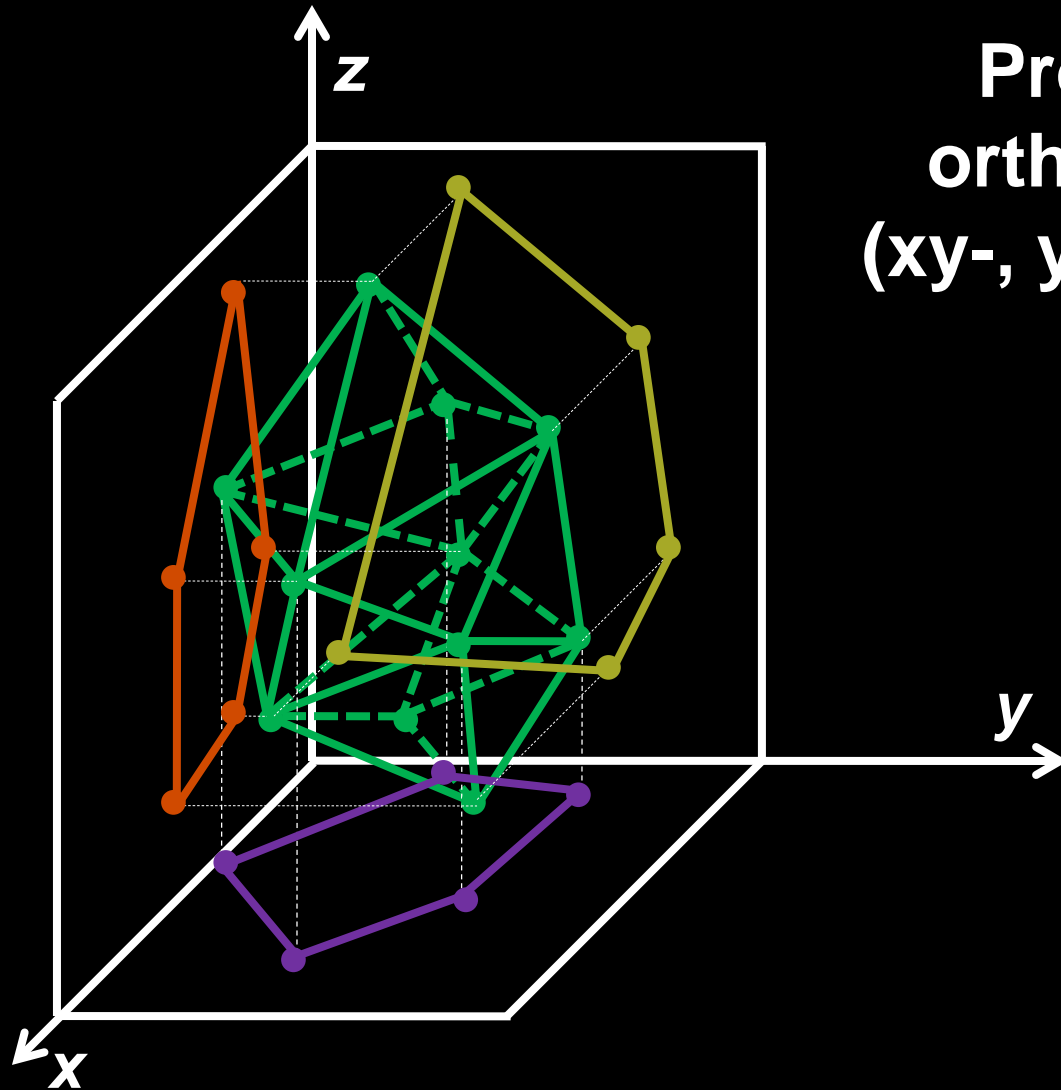
1. Approximate 3D Convex Hull with 2 x 2D Convex Hull
2. Use two-hop greedy forwarding
3. Simplify (details in paper)

# GDSTR-3D



**Example of a  
3D convex hull**

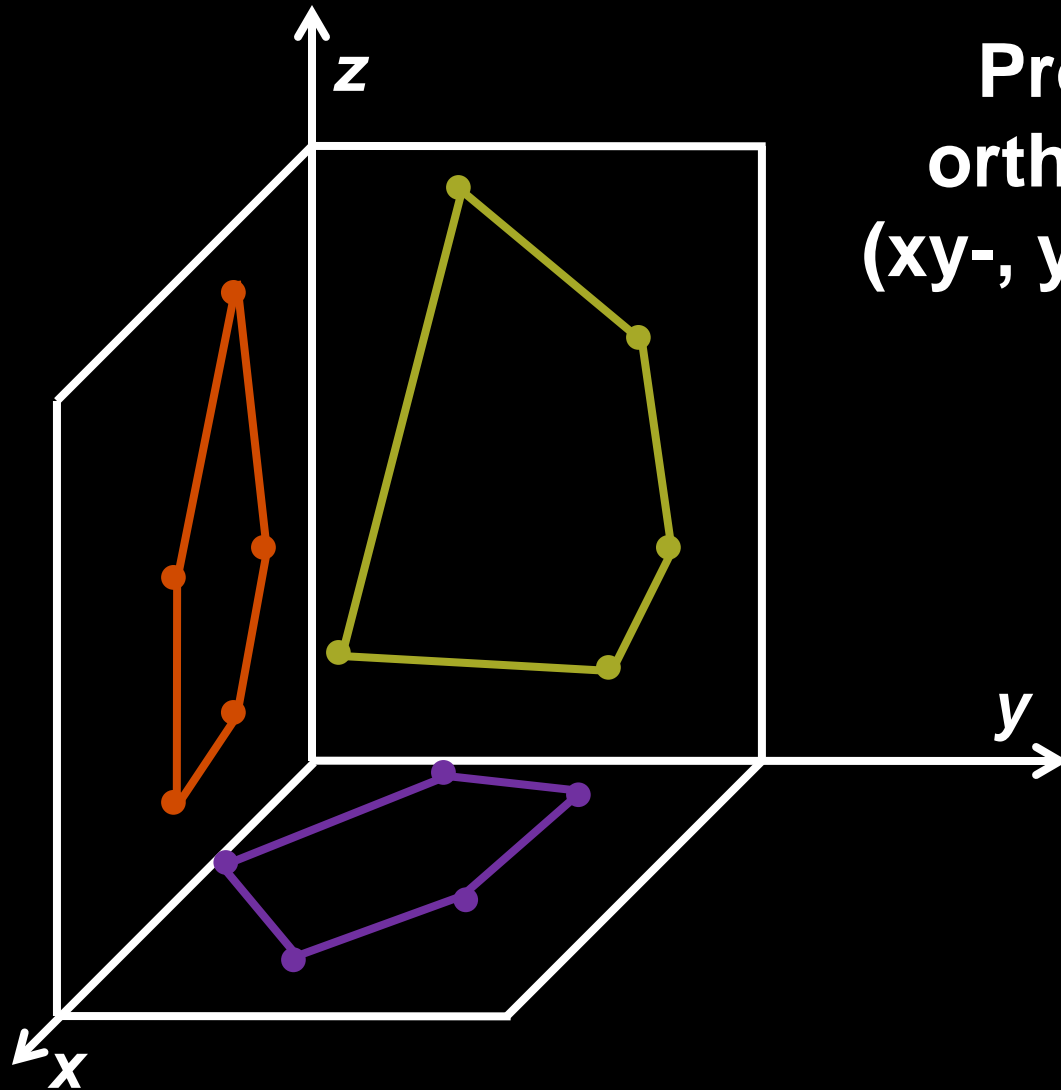
# GDSTR-3D



Projection onto  
orthogonal planes  
( $xy$ -,  $yz$ -, and  $zx$ -plane)

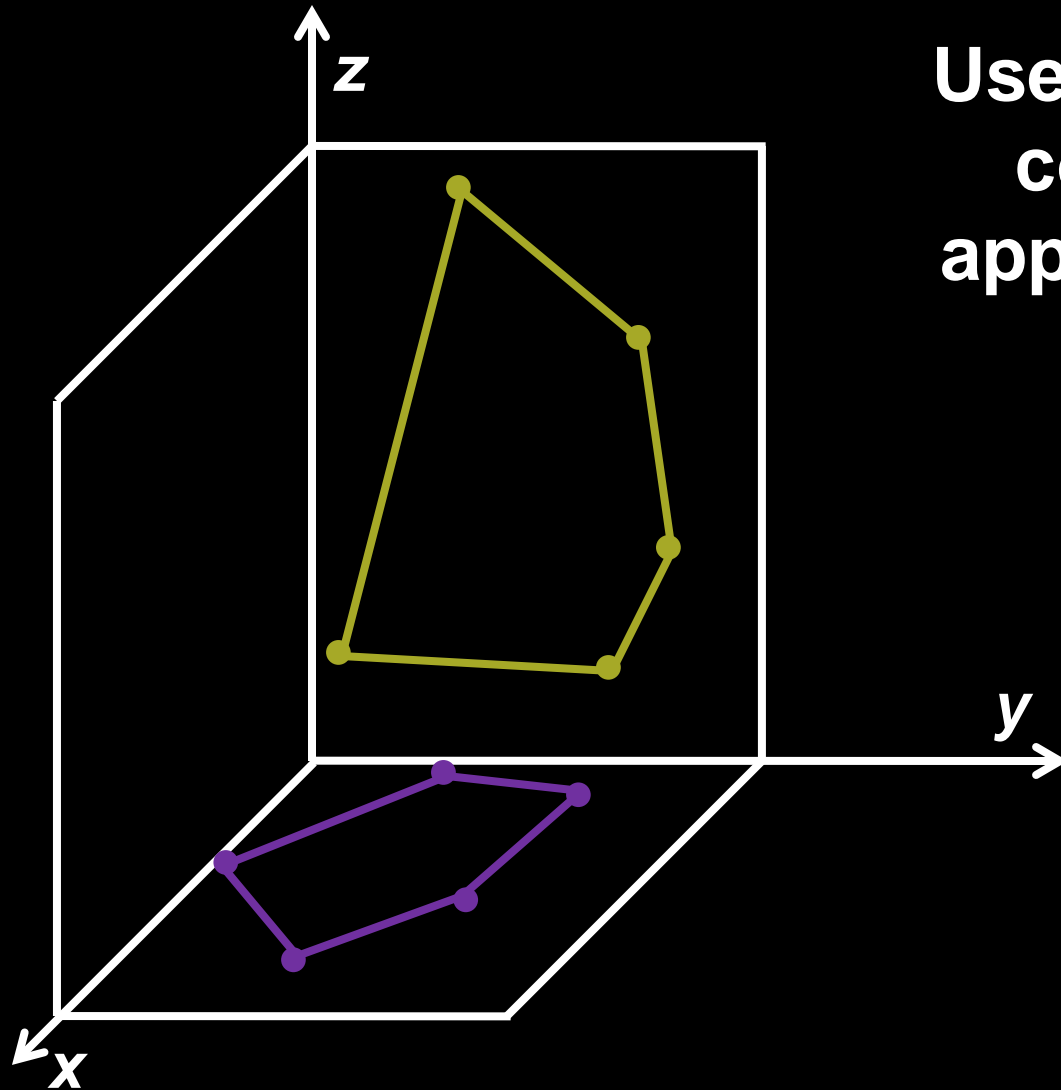


# GDSTR-3D



Projection onto  
orthogonal planes  
( $xy$ -,  $yz$ -, and  $zx$ -plane)

# GDSTR-3D



Use two of these 2D  
convex hulls to  
approximate the 3D  
convex hull

# PERFORMANCE EVALUATION

# Metrics

1. Success rate
2. Hop stretch
3. Maximum Storage
4. Message Overhead

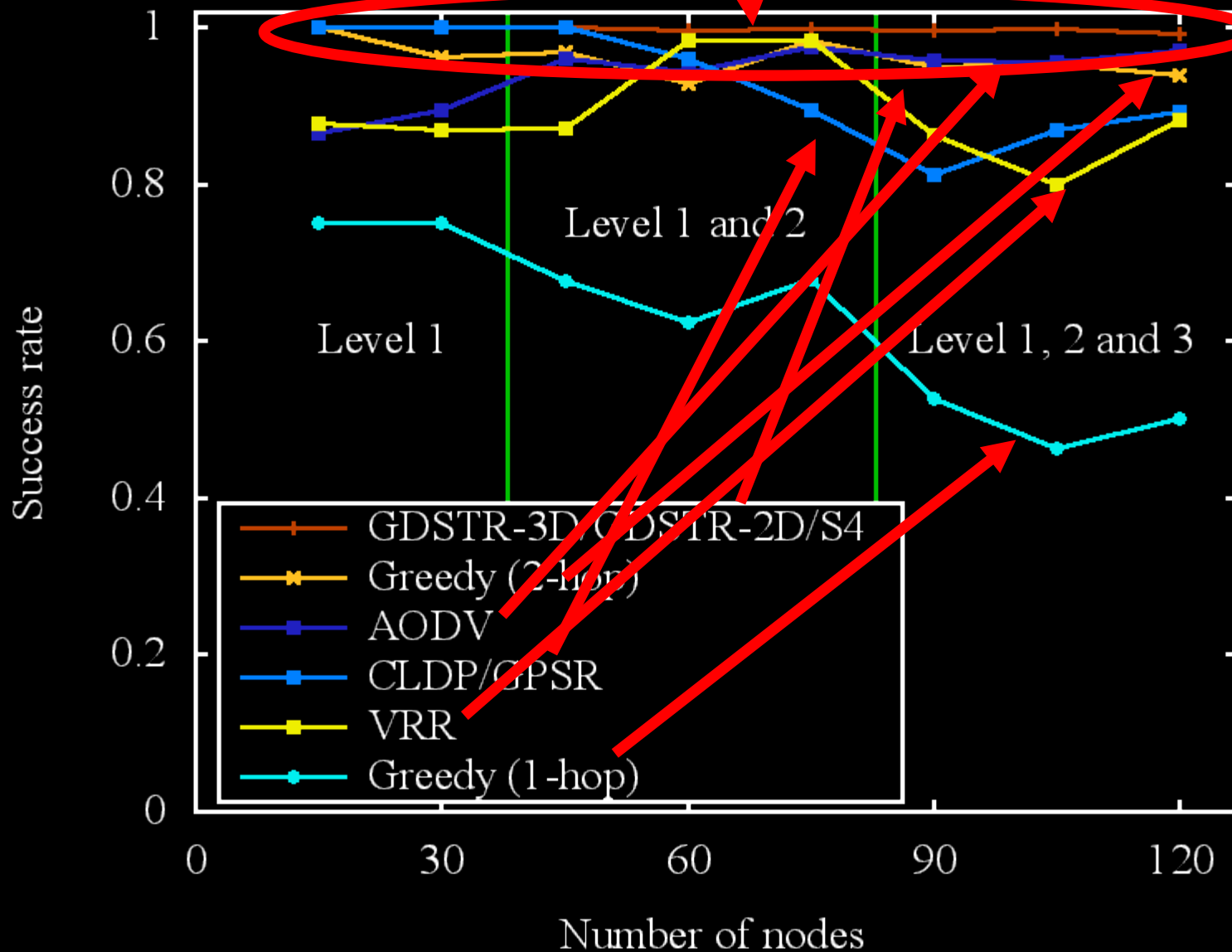
# Indriya Testbed (NUS)

- 127 TelosB motes  
distributed over 3 floors
- Picked random subsets  
of nodes on 1, 2 and 3  
floors

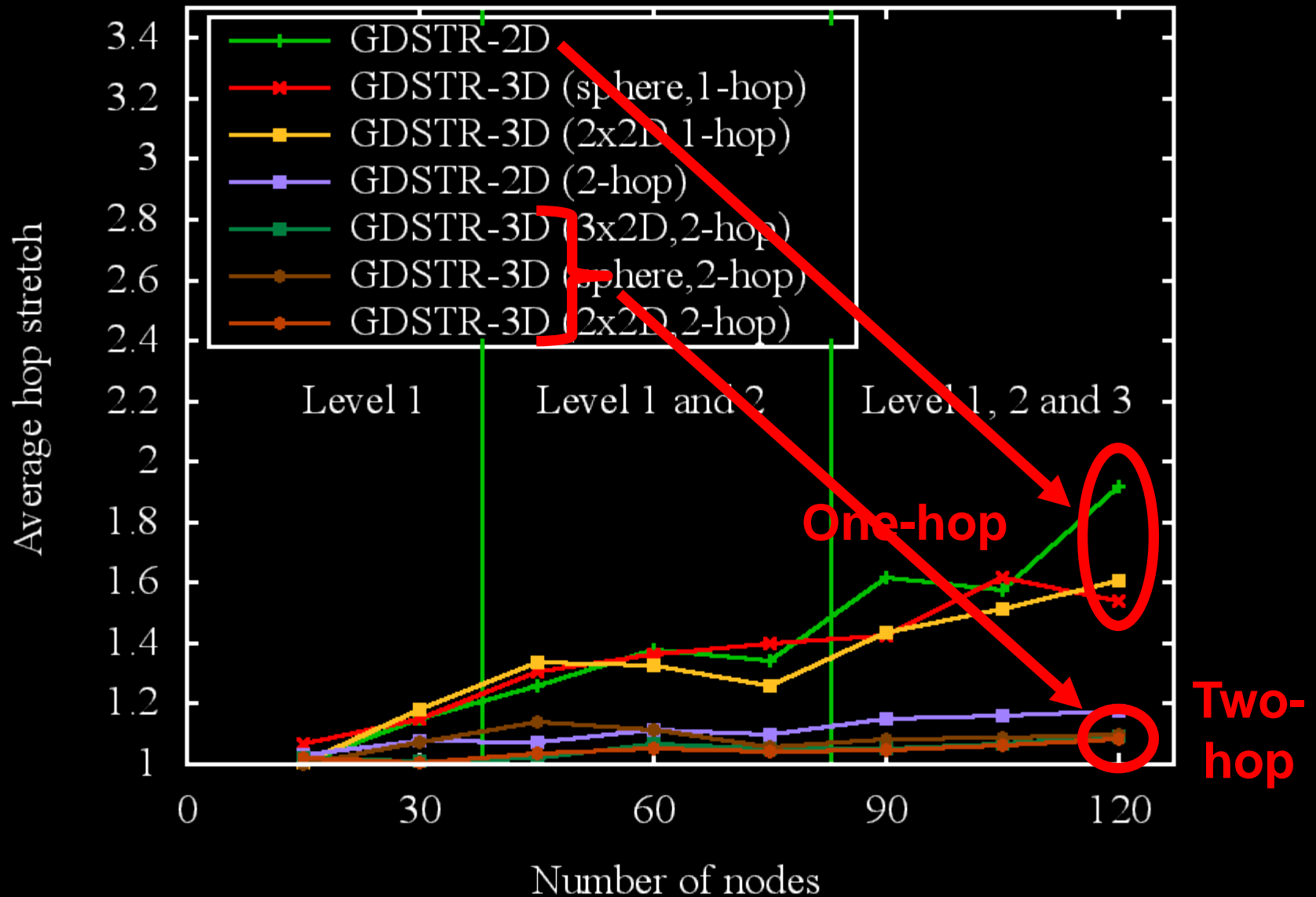
# Algorithms

1. GDSTR-3D
2. GDSTR(-2D)
3. CLDP/GPSR (2D Face Routing)
4. AODV
5. VRR
6. S4

# Success rate vs. network size

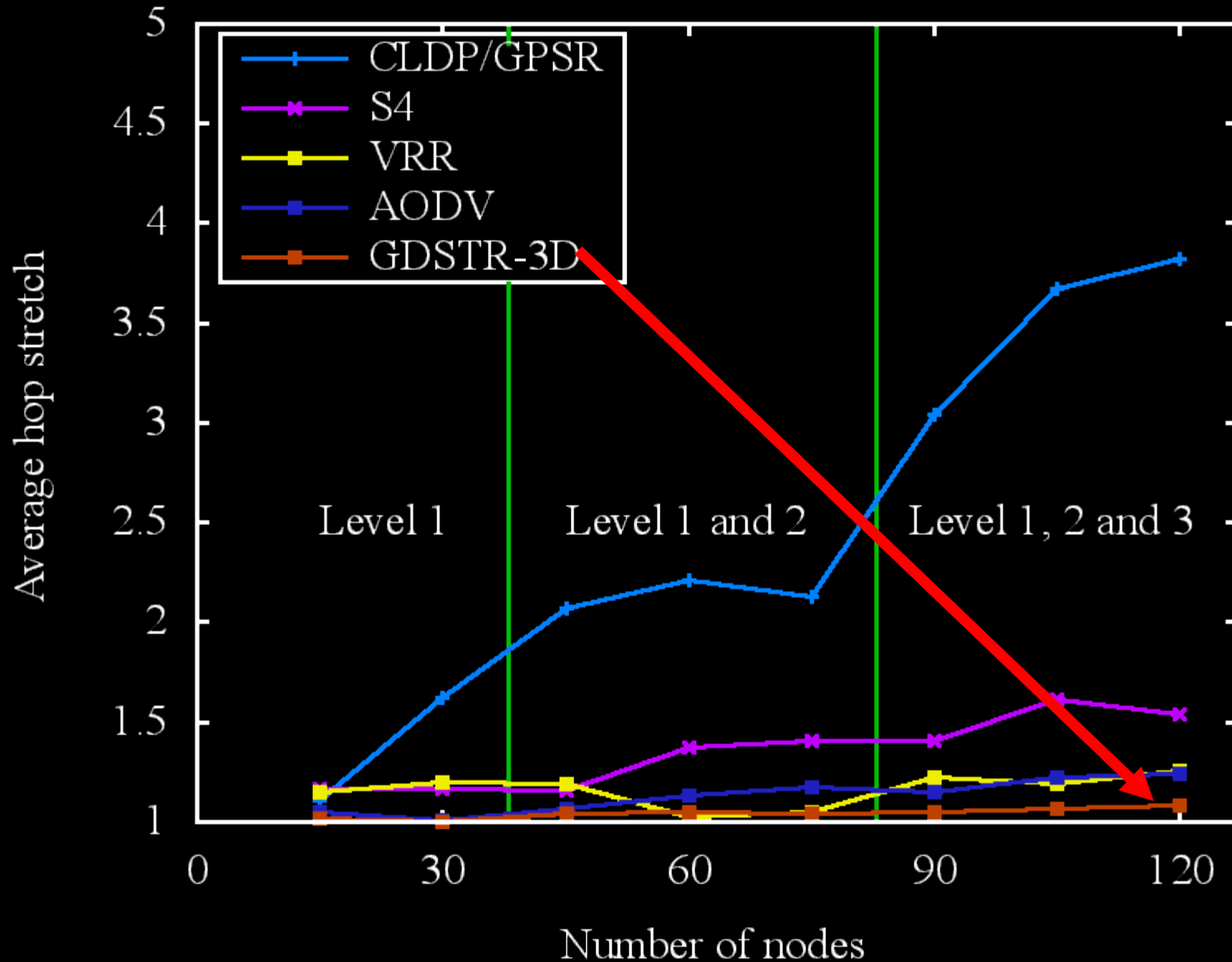


# Hop stretch (GDSTR+) vs. network size





# Hop stretch vs. network size

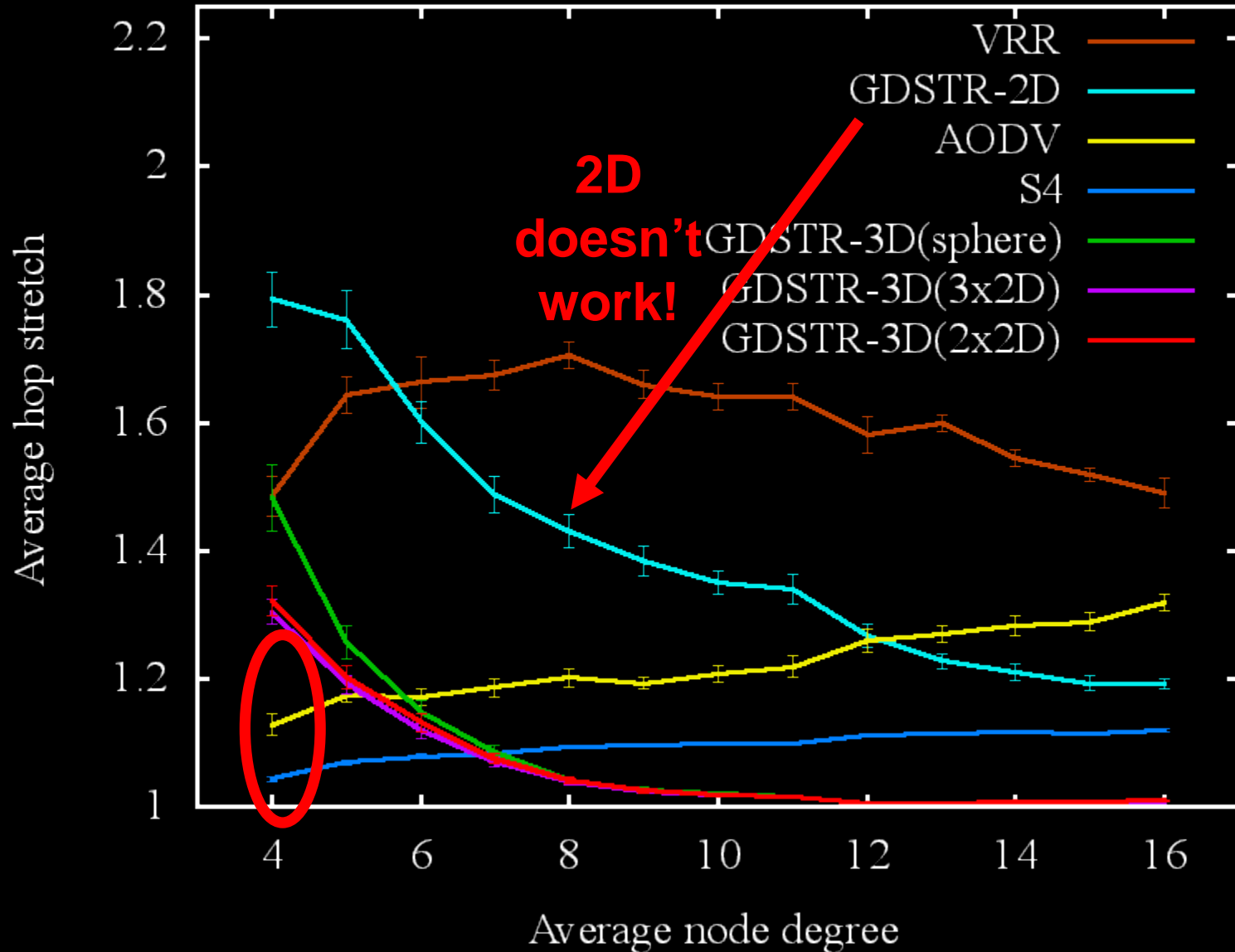


# Size of compiled binaries & source code

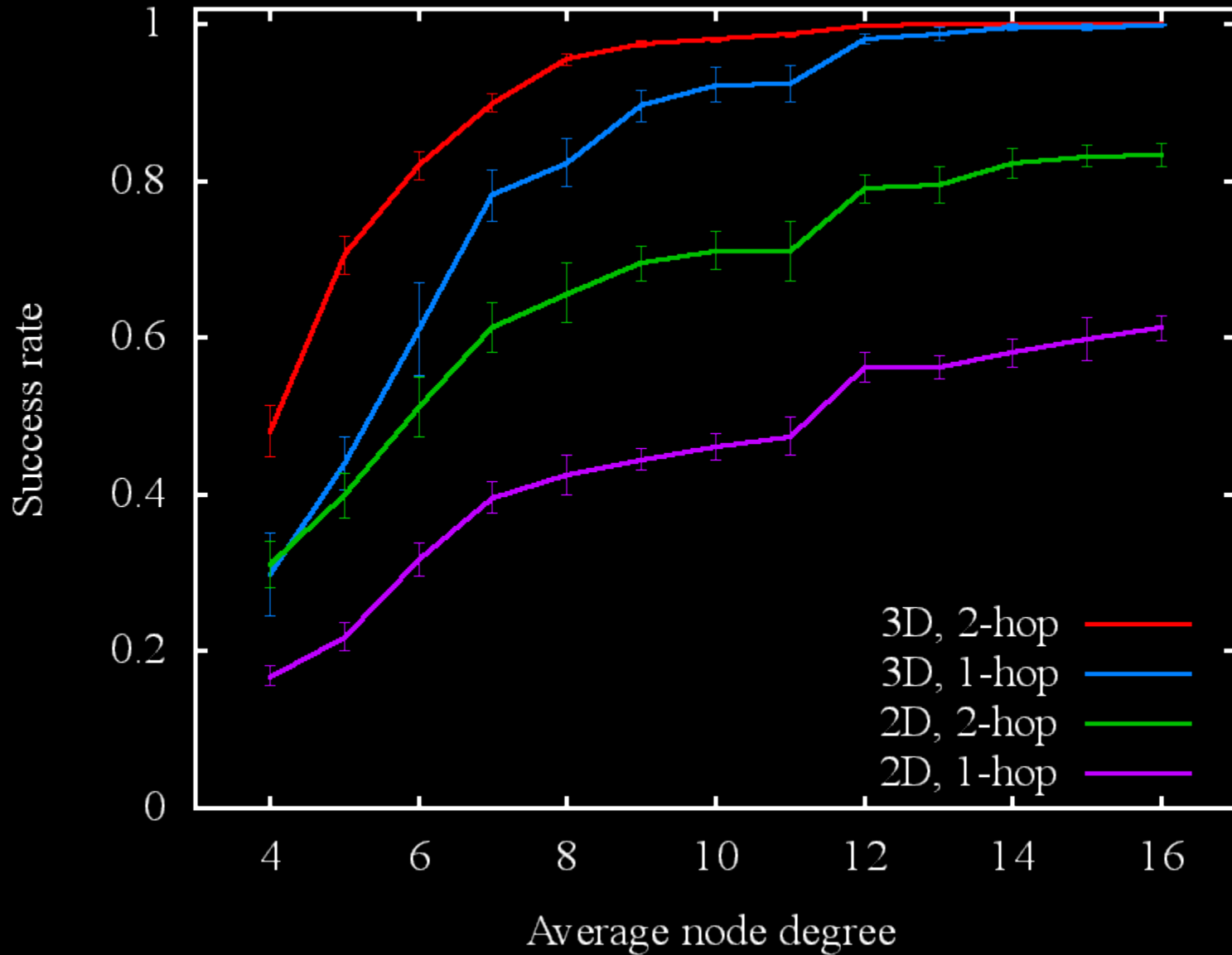
Algorithm	Compiled binary Size (KB)	Lines of code
<b>GDSTR-3D</b>	<b>39.5</b>	<b>2,757</b>
<b>GDSTR</b>	<b>33.8</b>	<b>2,641</b>
<b>CLDP/GPSR</b>	<b>47.5</b>	<b>2,500</b>
<b>S4</b>	<b>43.2</b>	<b>3,997</b>
<b>VRR</b>	<b>45.1</b>	<b>4,135</b>
<b>AODV</b>	<b>21.1</b>	<b>1,294</b>

# TOSSIM Experiments

# Hop stretch vs. network density

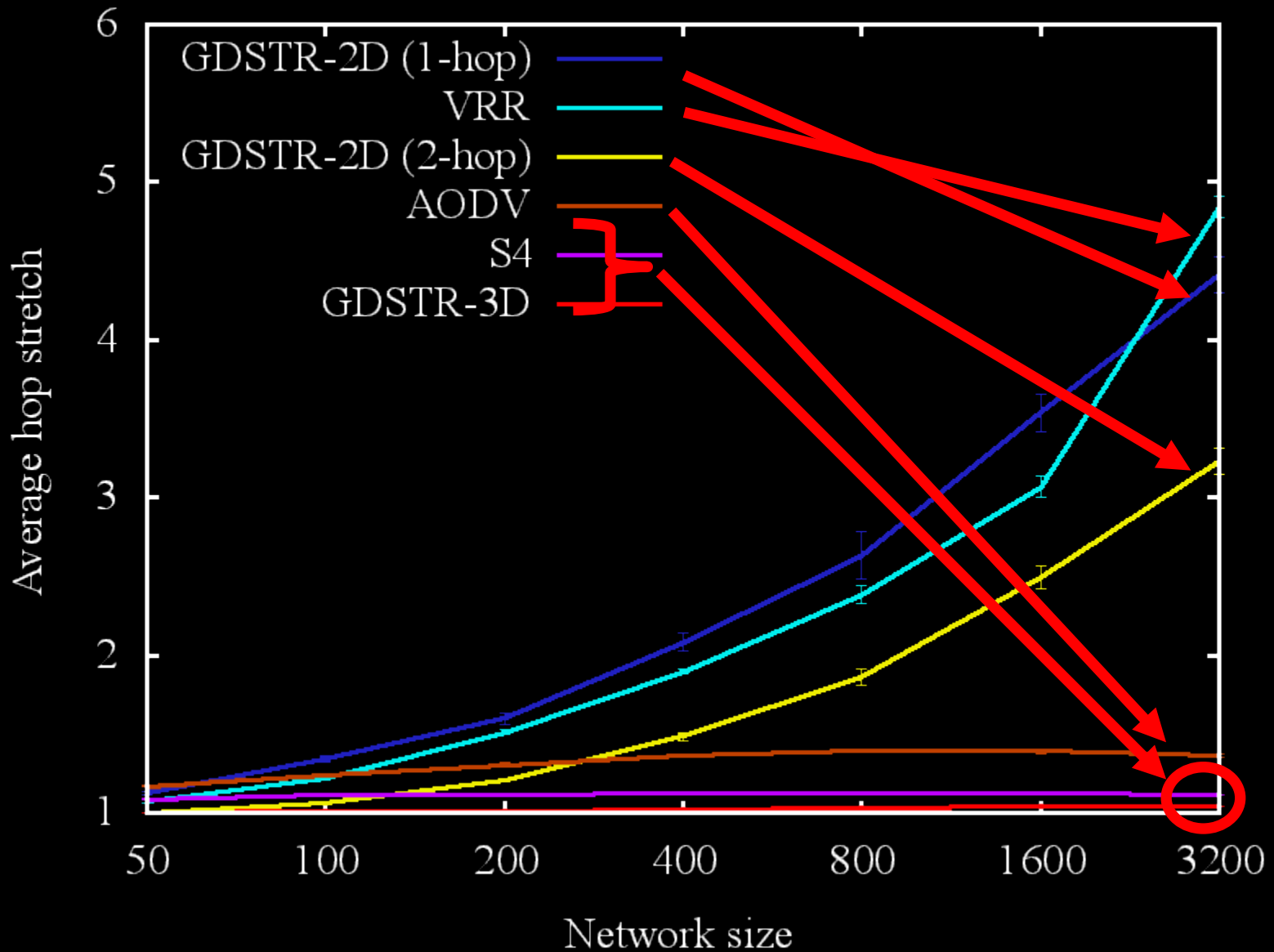


# Greedy forwarding success rate vs. network density

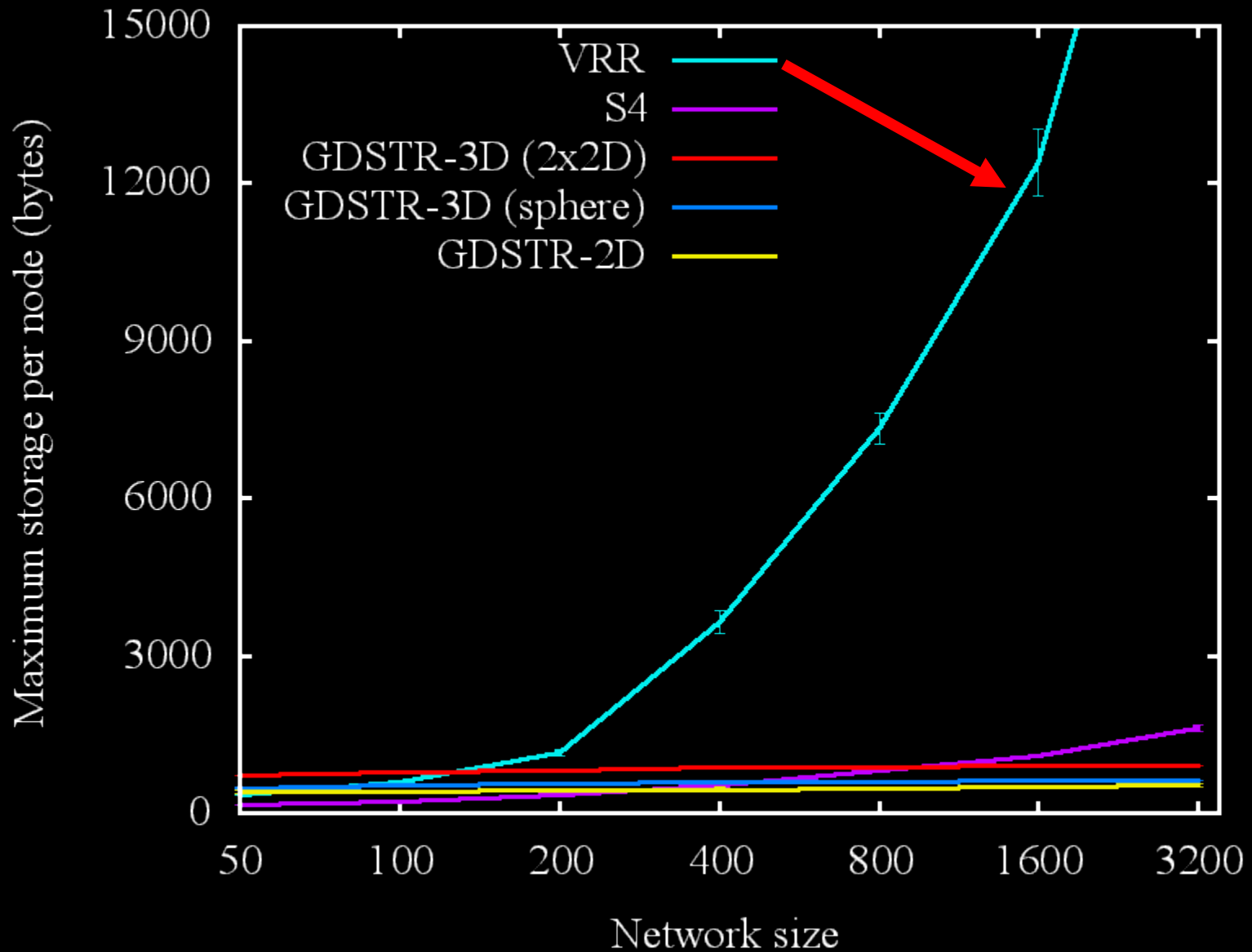


Scaling Up

# Hop stretch vs. network size

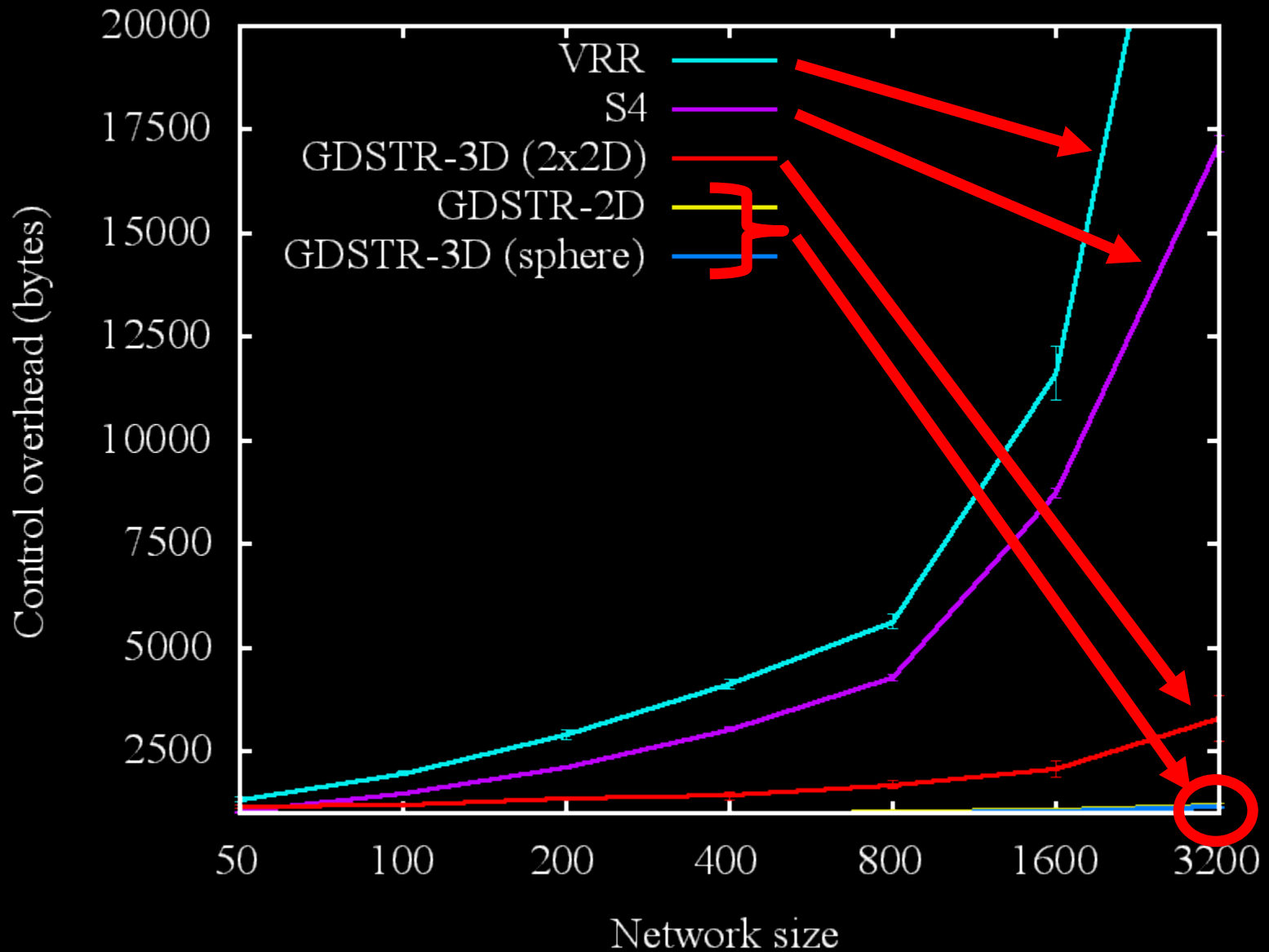


# Maximum storage vs. network size





# Message overhead (bytes) vs. network size



# Summary: Scaling Up (3,200 nodes)

Algorithm	Stretch	Storage	Overhead
<b>GDSTR-3D</b>	✓	✓	-
<b>GDSTR-2D</b>	✗	✓	✓
<b>S4</b>	✓	✓	✗
<b>VRR</b>	✗	✗	✗
<b>AODV</b>	-	✓	?

Comprehensive comparison  
of GDSTR-3D to

1. AODV

2. VRR

3. S4

Details in the paper.

# Key Contributions

1. Practical 3D geographic routing
  - 2x2D hulls for aggregation
  - Two-hop greedy
2. Comprehensive comparison of state-of-art point-to-point algorithms for TinyOS

# Summary

For small sensor  
networks (<200 nodes):  
pick your favorite  
algorithm. 😊

For large sensor networks (~3,200 nodes), geographic routing algorithms are most scalable:

- relatively low overheads
- storage matters, but is not overriding consideration

Life's  
complicated

# Tradeoffs at a glance

Algorithm	Needs coordinates?	Needs location service	Reactive?
<b>GDSTR-3D</b>	✓	✓	✗
<b>S4</b>	✗	✓	✗
<b>VRR</b>	✗	✗	✗
<b>AODV</b>	✗	✗	✓



# Future Work

- More Thorough Comparison
  - link losses
  - quantify cost of location service/  
coordinate assignment
  - resilience
  - incremental costs
  - traffic pattern/load
- Sleep-wake duty cycle
- Reduce memory footprint

# TinyOS Source Code

Available here:

<https://sites.google.com/site/geographicrouting>

Or email me: [benleong@gmail.com](mailto:benleong@gmail.com)

Questions?

Thank  
You

**For large sensor networks ,  
geographic routing algorithms are  
most scalable:**

- guarantee packet delivery
- storage cost is proportional to network density but size
- motes have small RAM

# Choice:

- Extend existing 2D geographic routing algorithms to implement a 3D routing algorithm
- GDSTR is a natural candidate for extension

# Routing in 3D:

- Geographic routing in 3D topologies is intrinsically harder than routing in 2D topologies since greedy forwarding tends to encounter more local minima in general 3D topologies
- It is not entirely straightforward to extend GDSTR to 3D because that 3D convex hulls require significantly more storage and are much more computationally costly

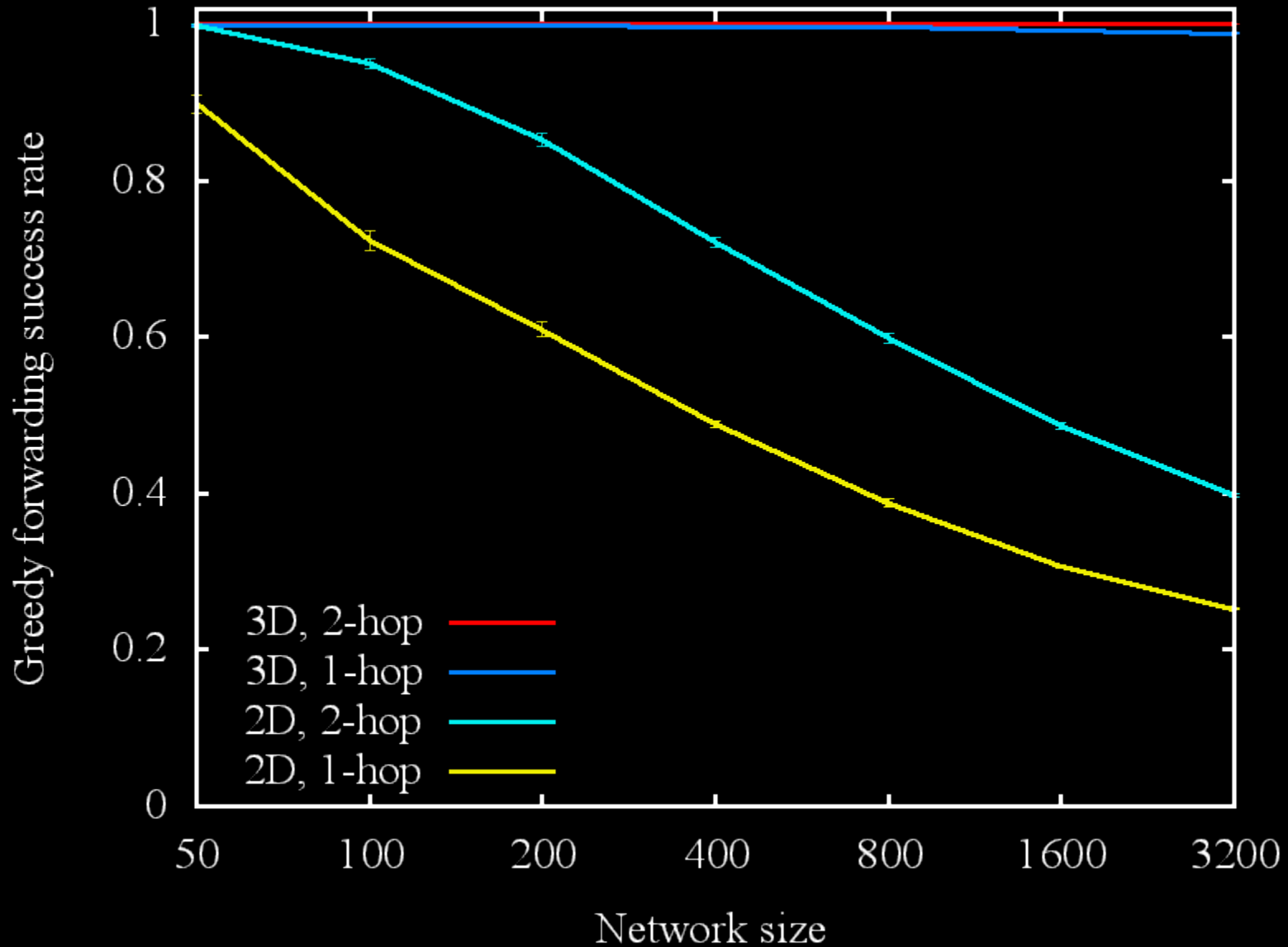
# Solution:

- Extend greedy forwarding by using 2-hop neighbor information to improve the greedy forwarding success rate in 3D networks
- Approximate 3D convex hulls with two 2D convex hulls

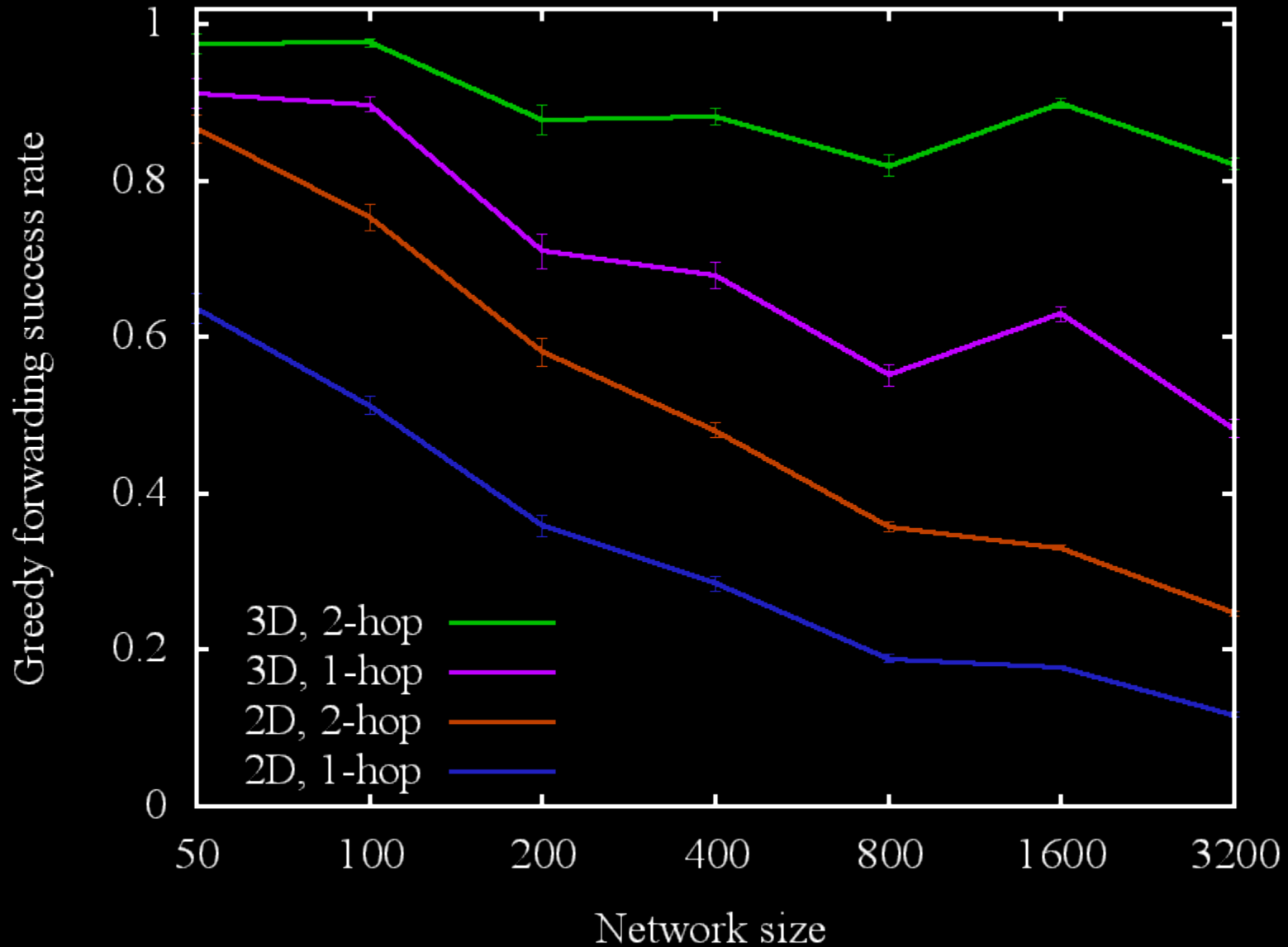


All graphs

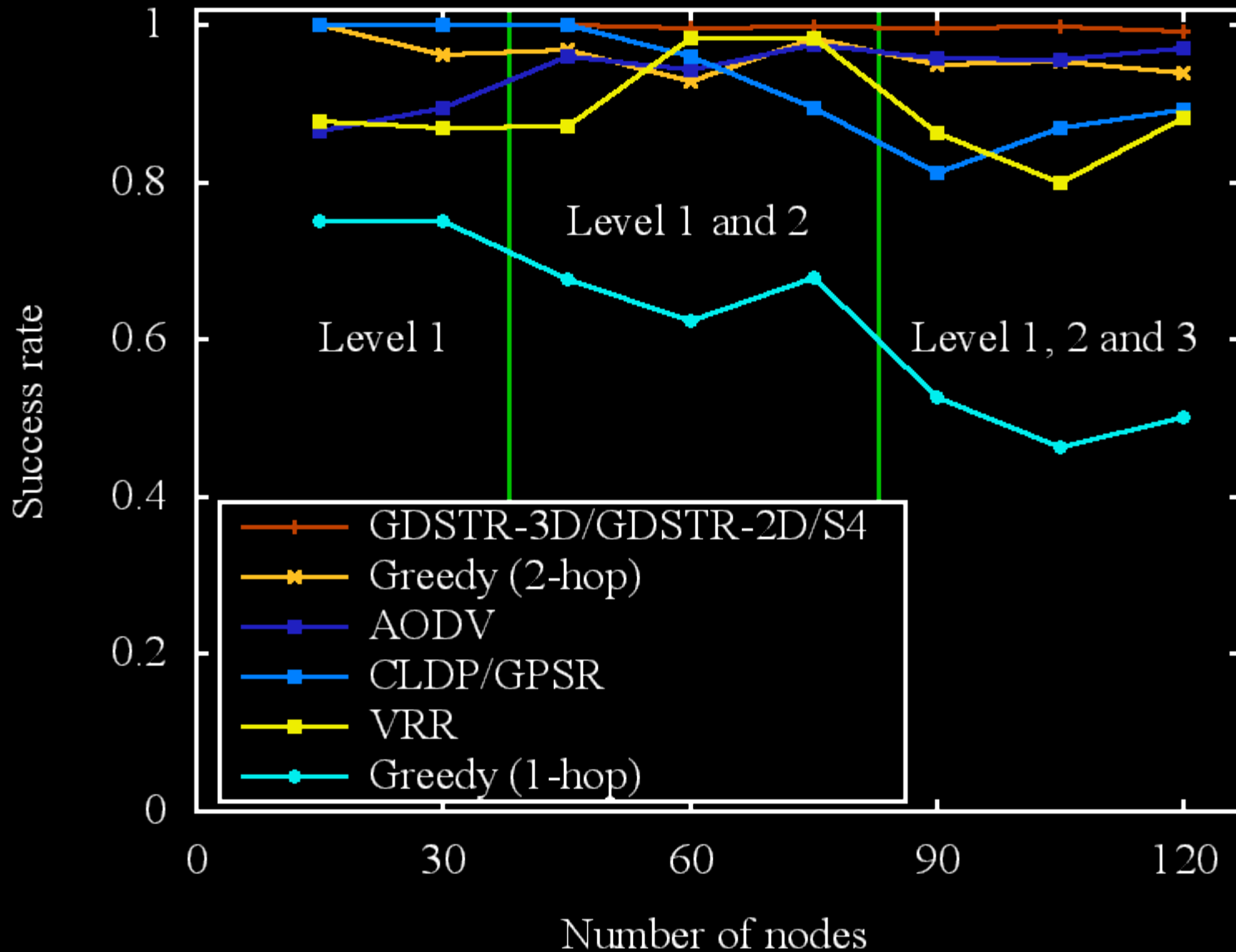
# Greedy forwarding success rate vs. network size (high density)



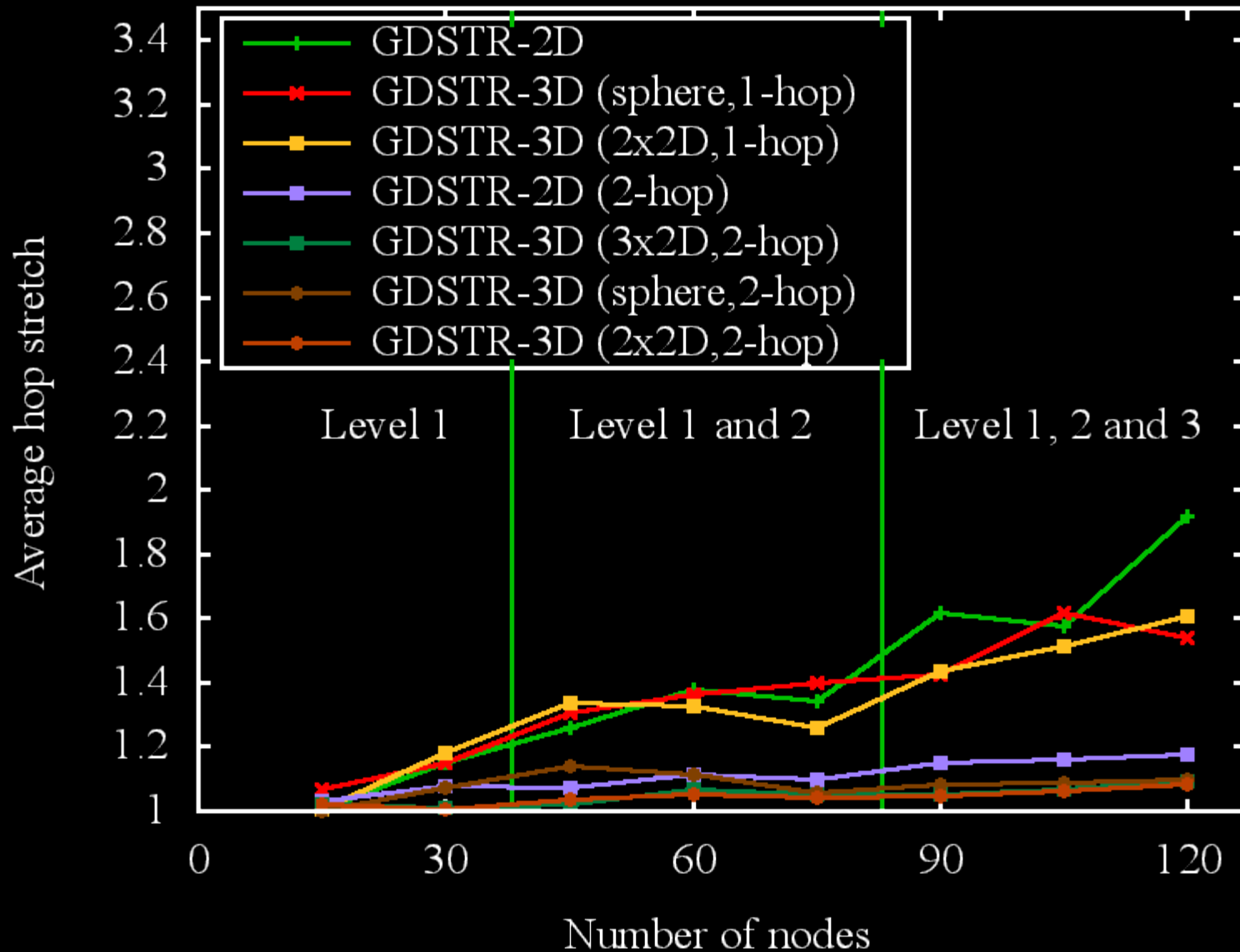
# Greedy forwarding success rate vs. network size (low density)



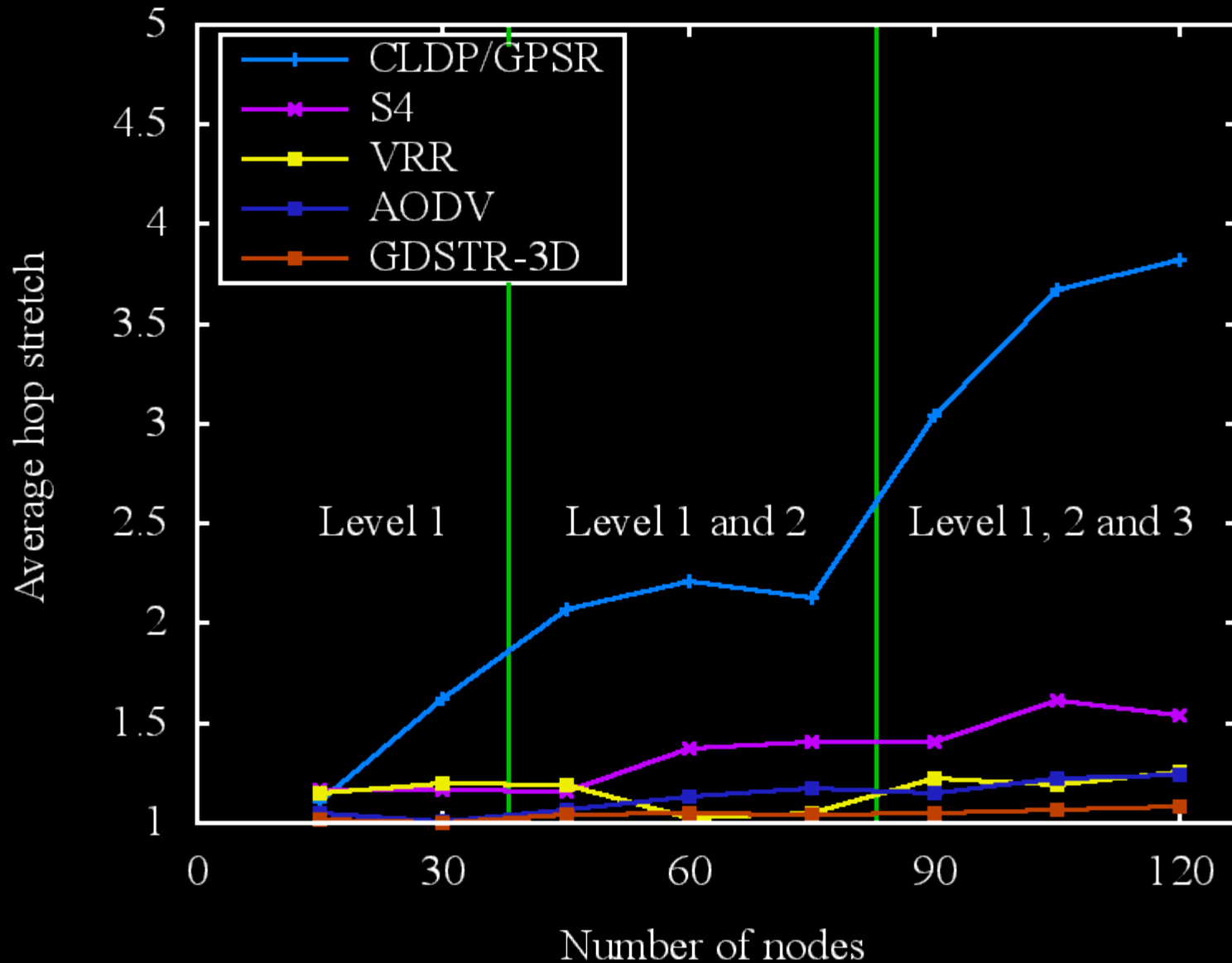
# Success rate vs. network size



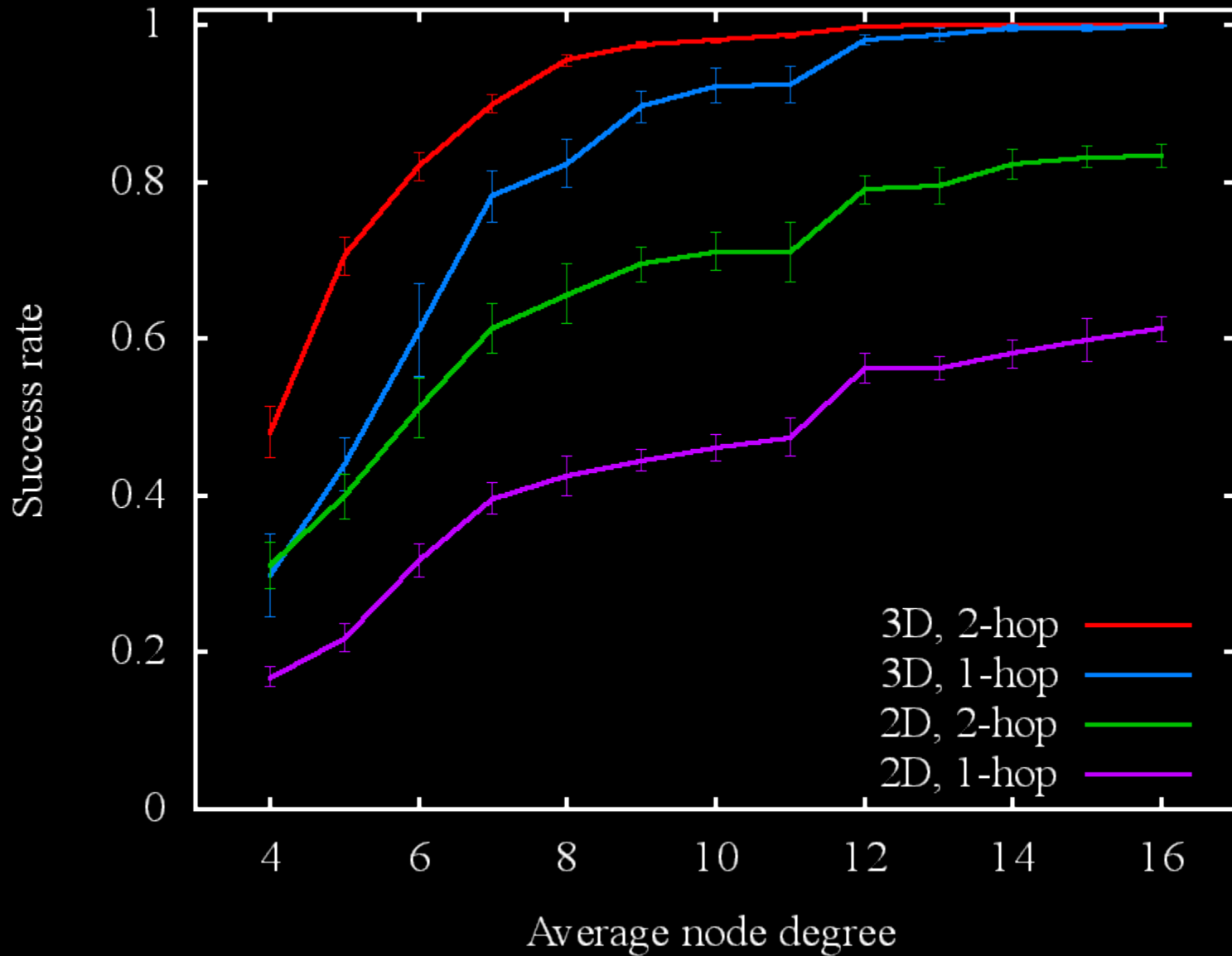
# Hop stretch(GDSTR+) vs. network size



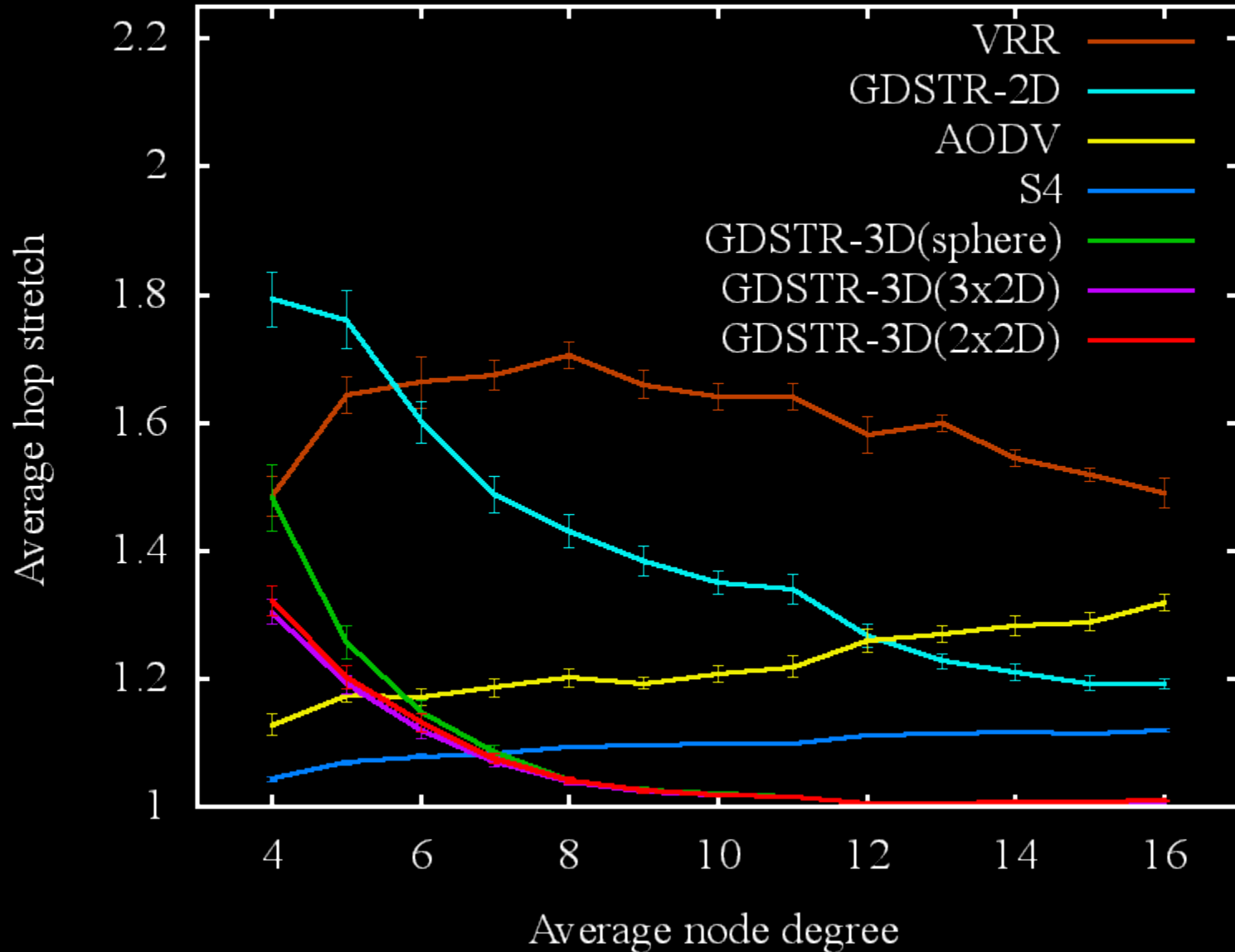
# Hop stretch vs. network size



# Greedy forwarding success rate vs. network density

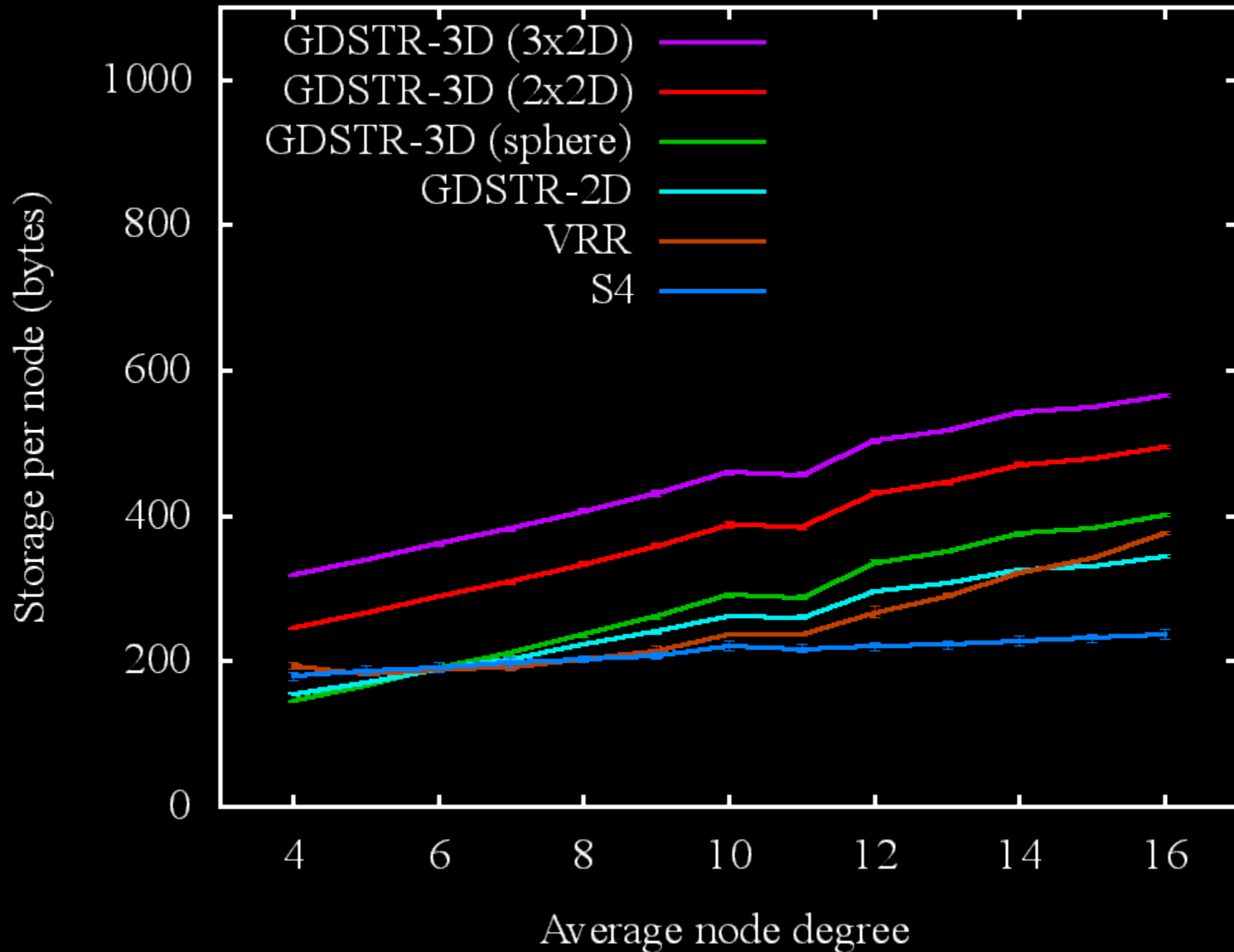


# Hop stretch vs. network density

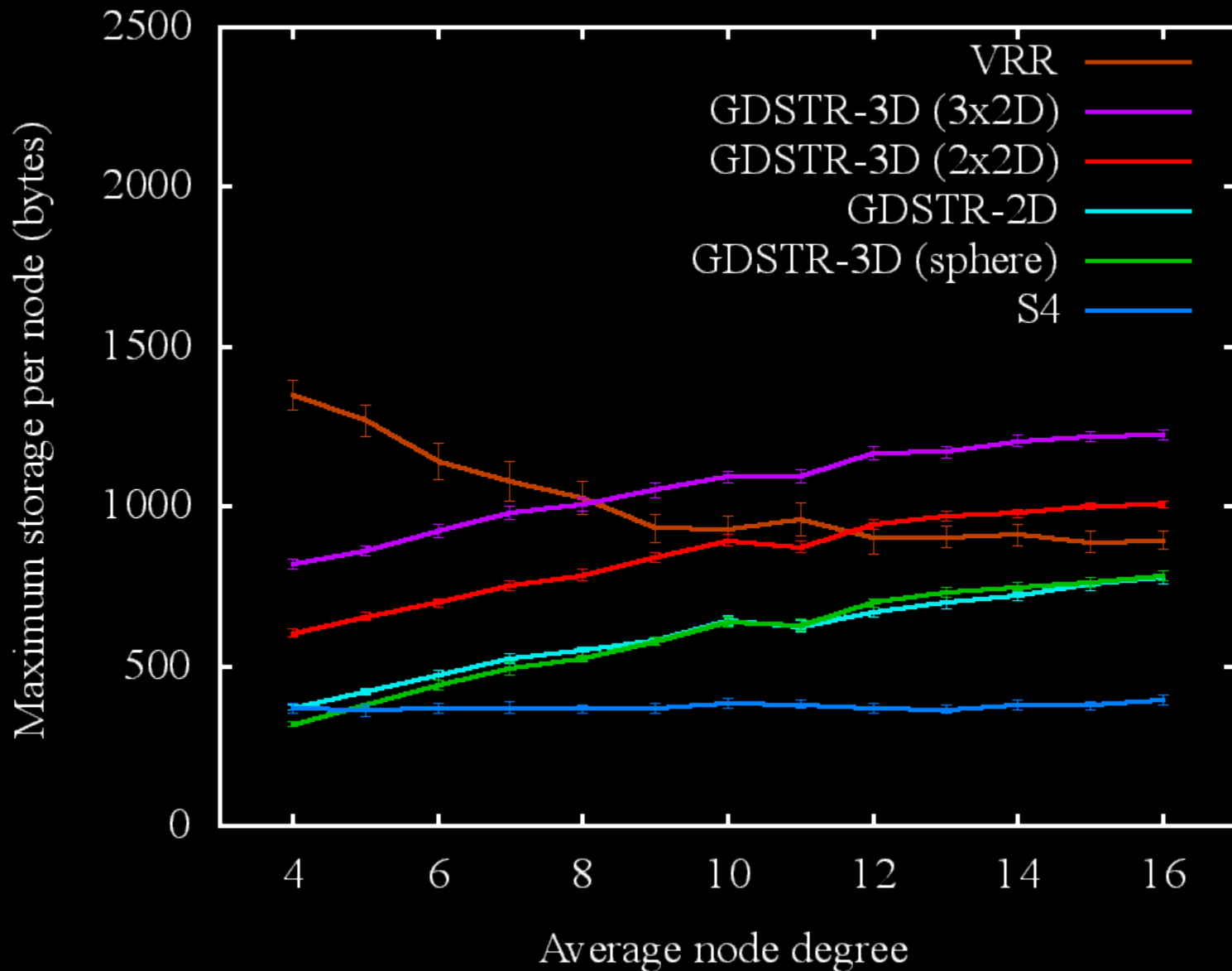




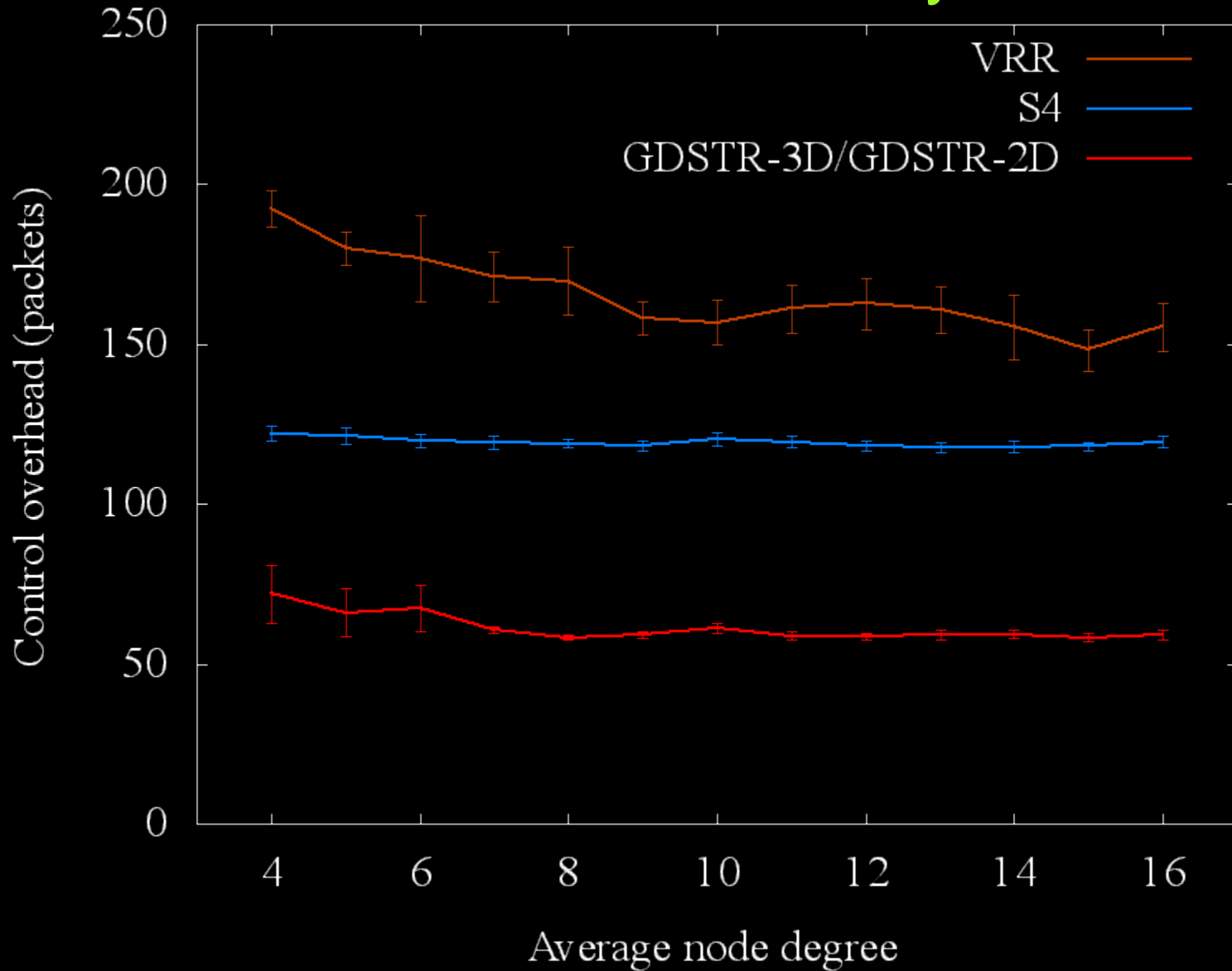
# Average storage vs. network density



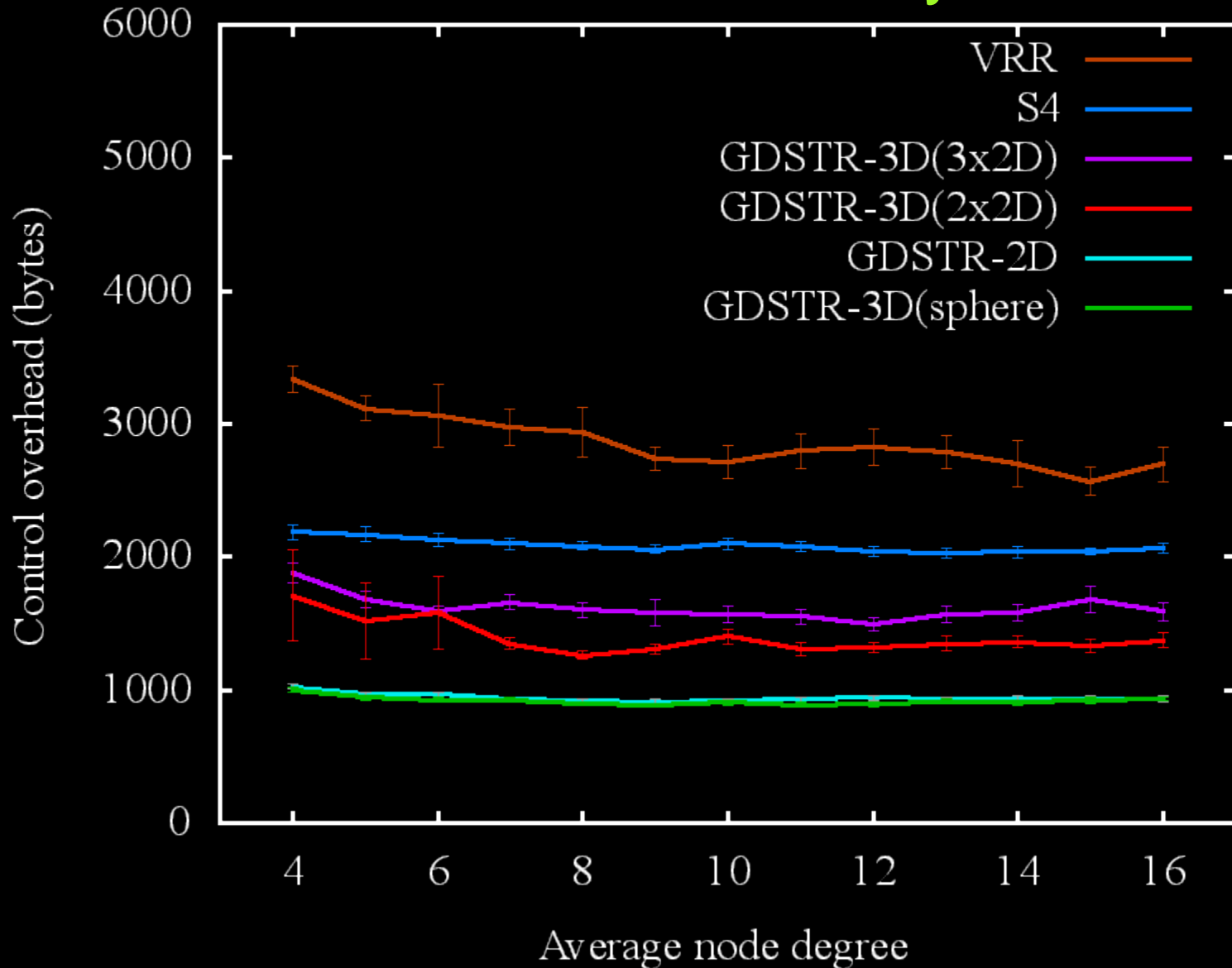
# Maximum storage vs. network density



# Message overhead(packets) vs. network density

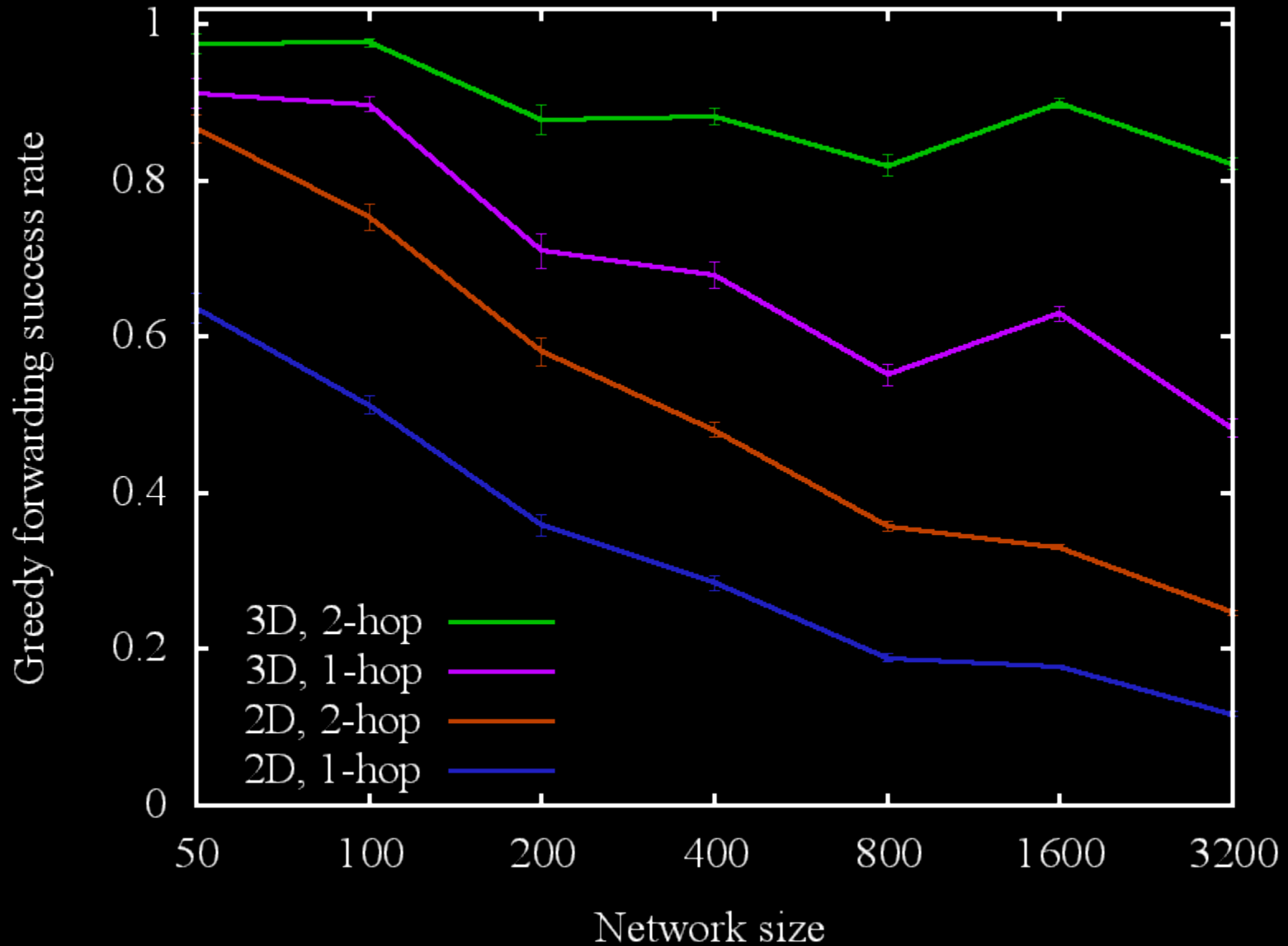


# Message overhead(bytes) vs. network density

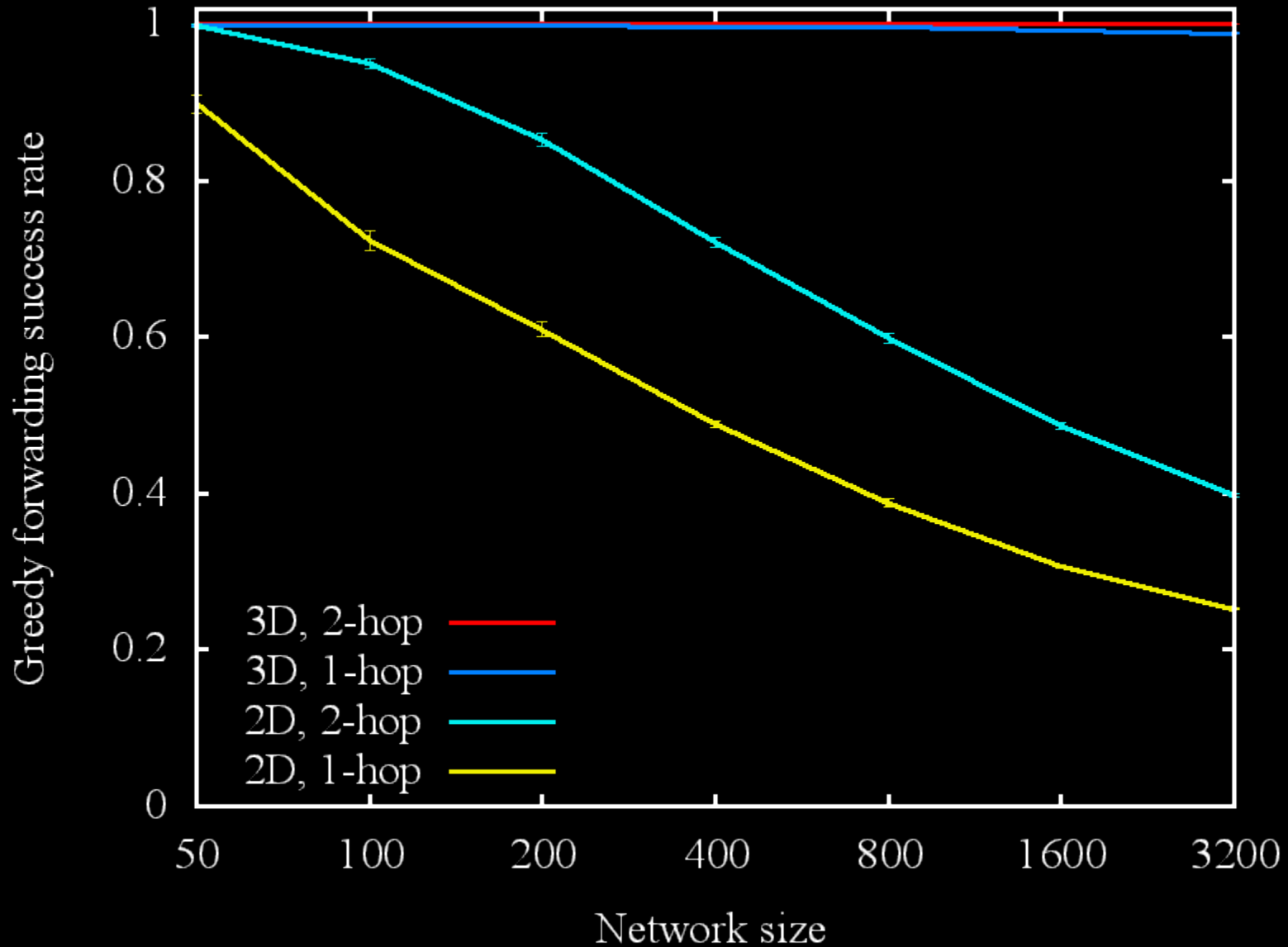


# Scaling up

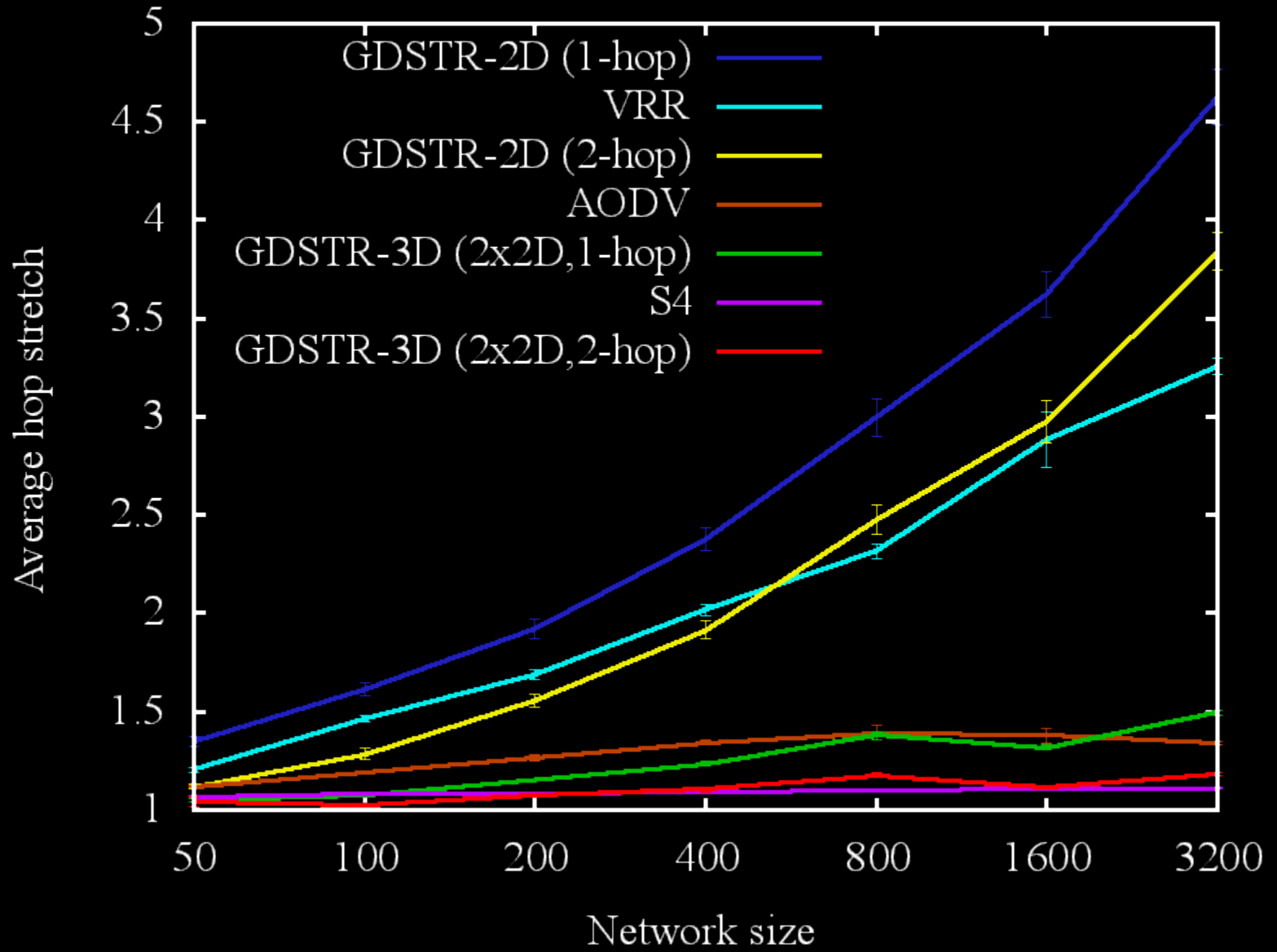
# Greedy forwarding success rate vs. network size (low density)



# Greedy forwarding success rate vs. network size (high density)

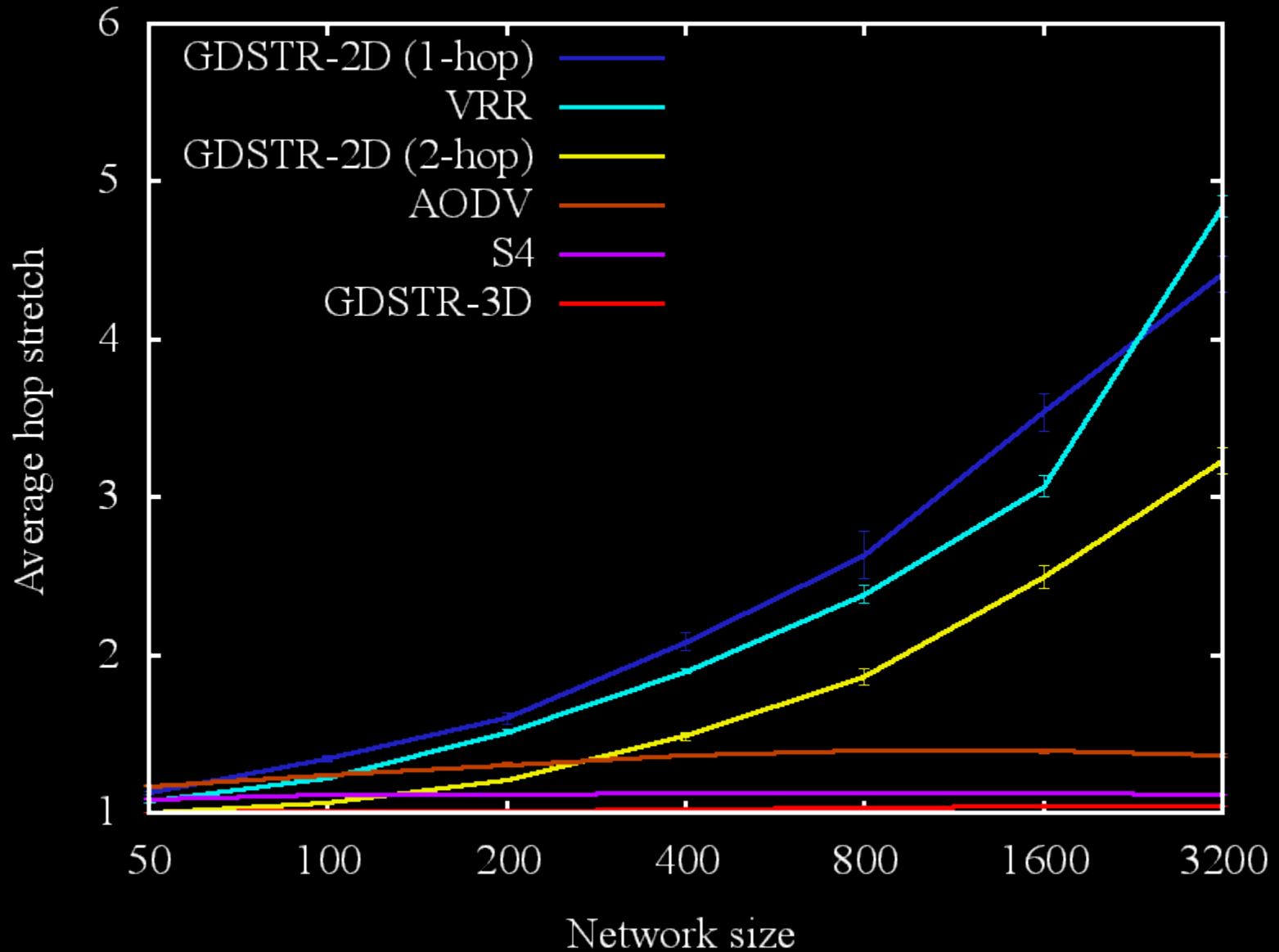


# Hop stretch vs. network size (low density)

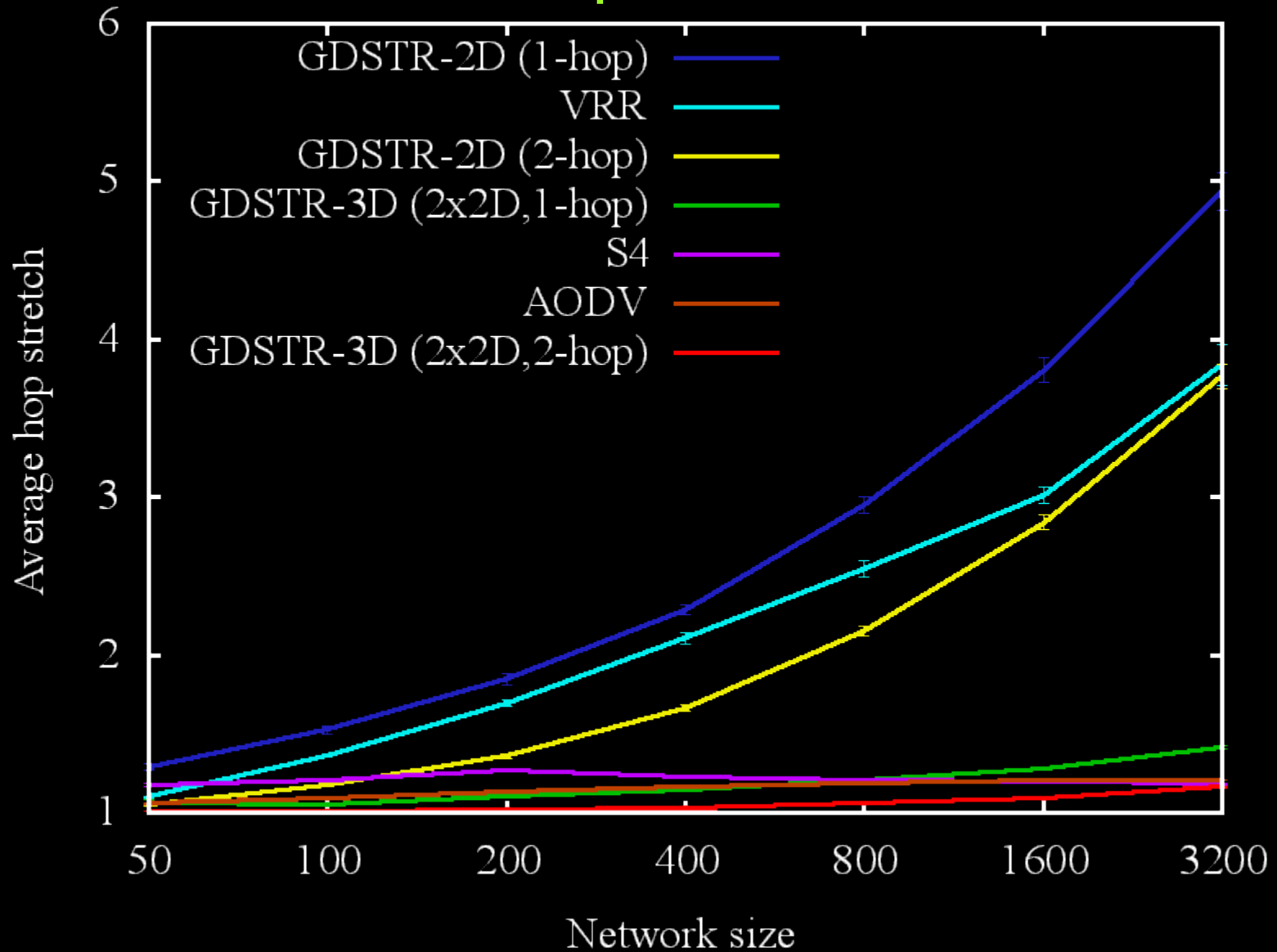




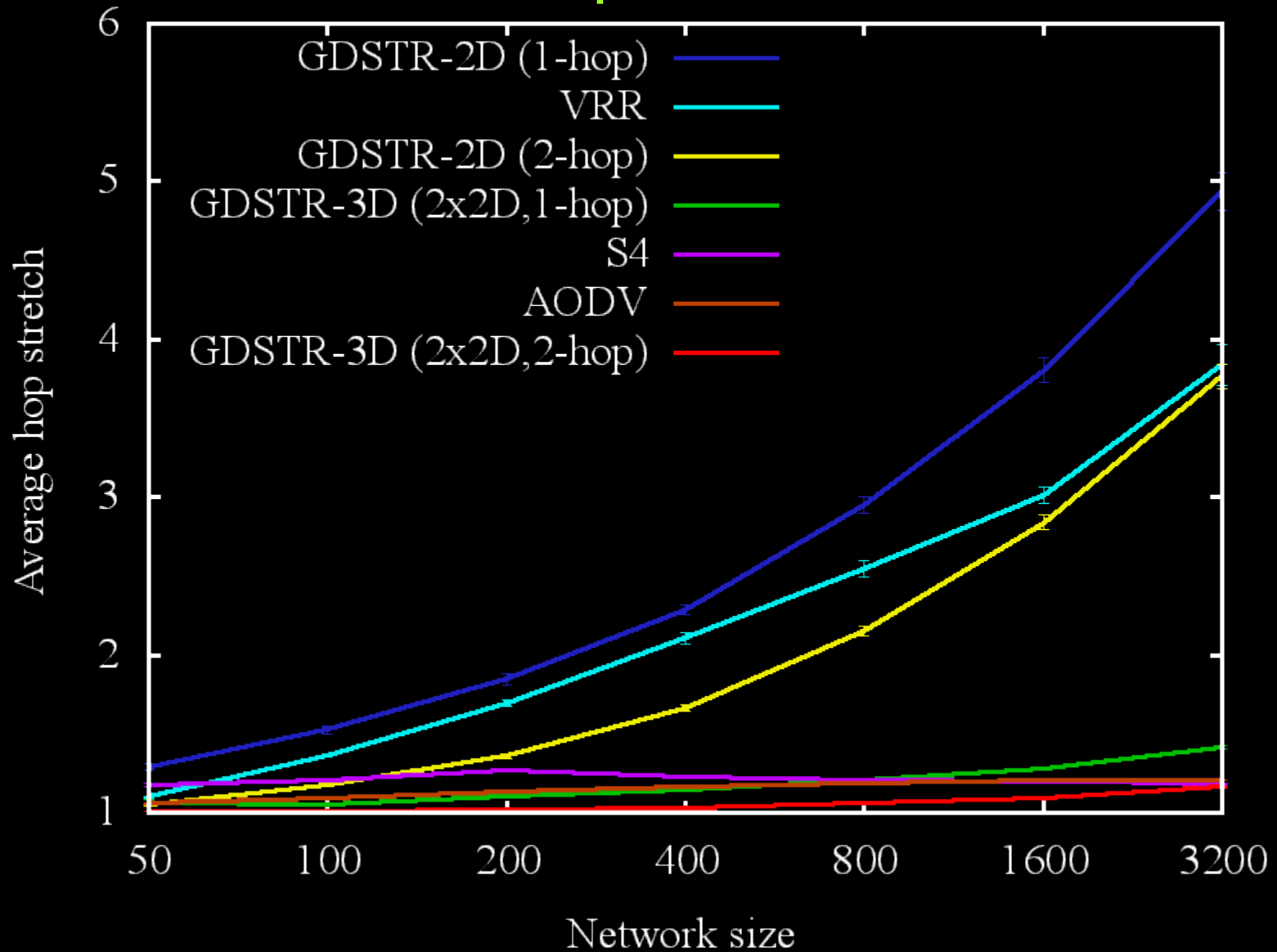
# Hop stretch vs. network size (high density)



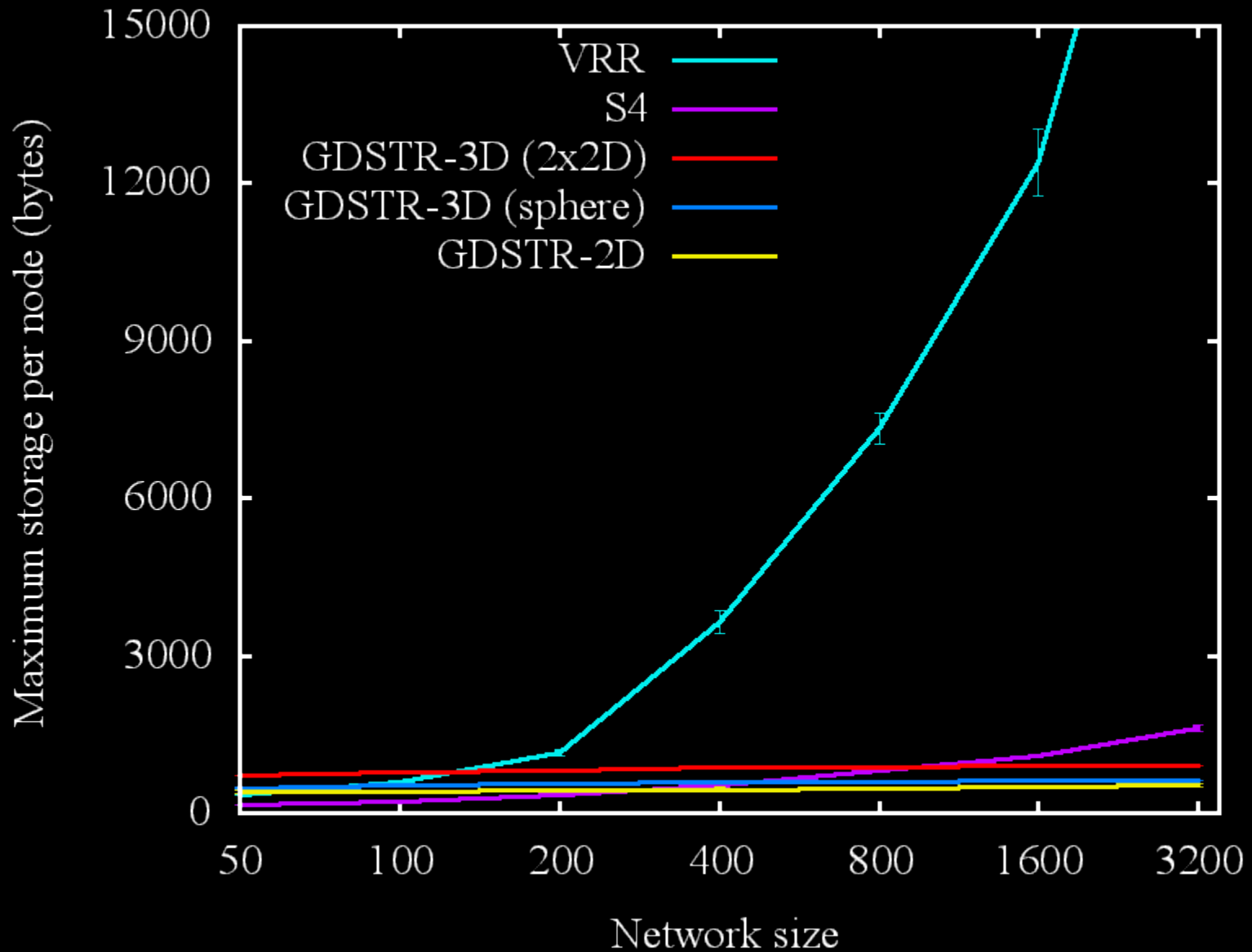
# Hop stretch vs. network size with multiple obstacles



# Hop stretch vs. network size with multiple obstacles



# Maximum storage vs. network size



# Message overhead(bytes) vs. network size

