CS2100 MID-TERM TEST ANSWER SHEET

STUDENT NO.:

NOTE: Write your particulars above legibly using PEN.
TUTORIAL GROUP: $\square$
1.

2.

3.

4.

5.
 [2 marks/MCQ]
7.

| C Code | MIPS Code |
| :---: | :---: |
| $\begin{aligned} & \text { int main(void) }\{ \\ & \text { i }=0 ; \\ & A= \text { 'A'; } \\ & a= \text { 'a'; } \\ & Z= \text { 'Z'; } \end{aligned}$ | [ addi \$s1, $\$$ \$zero, 0 ]  <br> [ addi $\$ s 3$, \$zero, 65 ] <br> [ addi $\$ s 4$, \$zero, 97 ]  <br> [ addi $\$ s 5, \$ z e r o, ~ 90$ ]  |
| do \{ | Loop: |
| ```if( str[i] >= 'A' && str[i] <= 'Z' ) {``` | $\left.\begin{array}{llll}\text { add } & \$ s 7, & \$ s 0, & \$ s 1 \\ \text { lb } & \$ t 2, & 0(\$ s 7) & \\ \text { slt } & \$ t 1, & \$ t 2, & \$ s 3 \\ \text { bne } & \$ t 1, & \$ z e r o, & \text { else } \\ {\left[\begin{array}{lll}\text { slt } & \$ t 1, & \$ s 5,\end{array}\right.} & \\ {[\text { bne } 2} & \$ t 1, & \$ z e r o, & \text { else }\end{array}\right]$ |
| func (str + i) ; | $\text { ret: }{ }^{\mathrm{j}} \text { func }$ |
| \} | else: |
| [ i++ ]; | addi \$s1, \$s1, 1 |
| \} while(str[i-1] ! = 0); | [ bne \$t2, \$zero, loop ] |
| $\text { \} return 0; }$ | quit: j exit |
| // Code omitted | \# Code omitted |
| void func(char* str) \{ | func: |
| $\begin{aligned} & {\left[{ }^{*} \text { str }={ }^{*}\right. \text { str + 'a' - 'A' ]; }} \\ & \text { return; } \end{aligned}$ | $\begin{array}{lll} \text { lb } & \$ \mathrm{t} 8, & 0(\$ \mathrm{~s} 7) \\ \text { addi } & \$ \mathrm{t} 8, & \$ \mathrm{t} 8, \\ \text { addi } & \$ \mathrm{t} 8, & \$ \mathrm{t} 8,-65 \\ \text { sb } & \$ \mathrm{t} 8, & 0(\$ \mathrm{~s} 7) \\ \text { j } & \text { ret } & \end{array}$ |
| \} |  |
| // Code omitted | \# Code omitted |
| // End of Program | exit: |

8. 
```
cs2100 is easy!
```

9. 

[2 marks]
$0 \times 15200003$
10.
[CHALLENGING]
[2 marks]
$2^{26}-9$ instructions $[67,108,855$ instructions]
11.
[6 marks]
$-5.84375$
12.
[6 marks]

| Registers File |  |  |  | ALU |  | Data Memory |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| RR1 | RR2 | WR | WD | Opr1 | Opr2 | Address | Write Data |
|  |  | $\$ 31$ | M[R[\$8]-R[\$0]] |  | $\mathrm{R}[\$ 0]$ | $\mathrm{R}[\$ 8]-\mathrm{R}[\$ 0]$ | $\mathrm{R}[\$ 0]$ |
| $\mathbf{\$ 8}$ | $\mathbf{\$ 0}$ | $/$ | $/$ | $\mathrm{R}[\$ 8]$ | $/$ | $/$ | $/$ |
|  |  | $\$ \mathrm{ra}$ | $\mathrm{M}[\mathrm{R}[\$ 8]]$ |  | 0 | $\mathrm{R}[\$ 8]$ | 0 |

13. [CHALLENGING]
```
sw $zero, O($zero) [or sw X, O($zero)]
```


## REFLECTIONS

Any thought about the module?
Share it in the thought bubble on the right. This will not be graded! ©


## WORKINGS:

1. By trial and error:

- $32_{8}=26_{10}=2 \times 13$
- $32_{9}=29_{10}=1 \times 29$ (this is a prime number, not a product of two primes)
- $32_{11}=35_{10}=5 \times 7$

NOTE: prime is still prime in any base. Note, 1 is NOT prime!
2. By tracing the value of $\$ \mathbf{t} 0, \$ \mathrm{t} 1$, and $\$ \mathrm{t} 2$ line by line

- lui $\$ t 0,0 x A A A A$
- srl \$t0, \$t0, 16
- lui $\$ t 0,0 x A 0 A 0$
- ori \$t1, \$zero, 0x5555
- and \$t2, \$t1 , \$t0

| \$t0 | \$t1 | \$t2 |
| :--- | :--- | :--- |
| 0xAAAA0000 | - | - |
| 0x0000AAAA | - | - |
| 0xAOAO0000 | - | - |
| 0xAOA00000 | 0x00005555 | - |
| 0xAOA00000 | 0x00005555 | 0x00000000 |

For Question 3-4, first compute the required number of bits:

- Registers: there are 6 registers, therefore 3-bits
- Addresses: there are 64 addresses, therefore 6-bits
- For Class A:
- 3 registers: 9-bits
- Opcode: 16-9 = $\underline{\text { 7-bits }}$
- For Class B:
- 1 address: 6-bits
- 2 registers: 6-bits
- Opcode: 16-6-6 = 4-bits

Draw a possible bit arrangement

3. Maximum is achieved by maximising Class A instruction (hence, minimising Class B ):

- Assign only 1 Opcode_1 to Class B: for instance, 0000 is for Class B
- Assign the rest to Class A:
- Opcode ${ }_{1}$ : 0001 to 1111: $2^{4}-1=15$
- Opcode $2: 2^{3}=8$
- Total: $15 \times 8=120$
- Sum for both classes: $120+1=121$

4. Minimum is achieved by maximising Class $B$ instruction (hence, minimising class $A$ ):

- Assign only 1 Opcode_1 to Class A: for instance, 0000 is for Class A
- Opcode 1 : 0000 only: 1
- Opcode $2: 2^{3}=8$
- Total: $1 \times 8=8$
- Assign the rest to Class B:
- Opcode 1 : 0001 to $1111: 2^{4}-1=15$
- Total: 15
- Sum for both classes: $8+15=23$

5. NOTE: Function call is pass-by-value. Therefore, the array numer inside rational is copied.
6. NOTE: On the other hand, we are passing pointers directly here. Hence, pass-by-reference.
7. IDEA:

- First 4 MIPS code: Refer to the ASCII table for the decimal value.
- Next 2 MIPS code: Perform slt followed by either bne or beq similar to the two lines above this.
- $\operatorname{str}[\mathbf{i}]<=$ 'Z'
- \$t2 <= \$s5

○ ! (\$s5 < \$t2)

- $(\$ \mathrm{~s} 5<\$ \mathrm{t} 2)=0$

○ slt \$t1, \$s5, \$t2 ; b bne \$t1, \$zero, else

- First $1 \mathbf{C}$ code: This is simply $\$ \mathbf{s} 1++$ (or $\mathbf{i + +}$ ) after mapping.
- Next 1 MIPS code: Simply bne \$t2, \$zero, loop (cannot use jloop here as $\mathbf{j}$ is unconditional). NOTE: since there is an $\mathbf{i + +}$ before, $\mathbf{s t r}[\mathbf{i}-\mathbf{1}]$ is exactly $\$ \mathbf{t} \mathbf{2}$.
- Last 1 C code: Start by translating the 4 MIPS code
- lb \$t8, 0(\$s7) $\Rightarrow \$$ t8 $=$ *str
- addi \$t8, \$t8, $97=>\$$ t8 $=\$$ t8 + ' $a \prime \Rightarrow \$ t 8=*^{\prime}$ str + 'a'
- addi $\$$ t8, $\$$ t8, $-65=>\$$ t8 $=\$$ t8 - ' $A$ ' $\Rightarrow \$$ t8 $=* s t r+$ 'a' - 'A'
- sb \$t8, 0(\$s7) $\quad>{ }^{*}$ str $=\$$ t8 $\quad \Rightarrow * s t r=*^{\prime}$ str + 'a' - 'A'

8. TRACE: The program basically converts uppercase to lowercase: "cs2100 is easy!"
9. STEPS:

- Compute immediate value: bne \$t1, \$zero, 3
- Compute registers value: bne \$9, \$0, 3
- opcode: 000101
- rs: 01001
- rt: 00000
- immediate: 0000000000000011
- Binary: 00010101001000000000000000000011
- Hexadecimal: $1 \quad 5 \quad 2 \quad 0 \quad 0 \quad 0 \quad 0 \quad 3$
- Answer: 0x15200003

10. STEPS:

- Compute maximum jump using $\mathbf{j}$ instruction
- j func: between ret label and func label, need to be within 256 MB boundary
- j exit: between top \#Code omitted and exit label, need to be within 256MB boundary
- j ret : between bottom \#Code omitted and ret label, need to be within 256 MB boundary
- Due to overlap between these $\mathbf{j}$ instructions, it implies between ret label and exit label are within 256 MB boundary
- 256 MB boundaries contain $2^{26}$ instructions.
- Subtract 8 instructions already in the region.
- Subtract 1 since we must jump to exit label, which means there must be one instruction there at the exit label. Note that the label is outside the \#Code omitted region.
- Total: $2^{26}-9$


## 11. STEPS:

- Convert to binary: 11000000101110110000000000000000
- Split into region:
- Sign: 1 (negative)
- Exponent: 10000001
- Convert to decimal: 129
- Excess-127: $\quad 129=127+2=>2$
- Mantissa: 01110110000000000000000
- Normalize: $\quad 1.0111011 \times 2^{2}$
- Remove exponent: 101.11011
- Convert to decimal: $5+2^{-1}+2^{-2}+2^{-4}+2^{-5}$

○ Sum: $\quad 5+0.5+0.25+0.0625+0.03125=5.84375$

- Answer: -5.8435

12. TRACE: Note that the Control signal is different. Be careful with which values are selected in multiplexer as well as the behaviour of each component (e.g., Register File and Data Memory). WR is the first 5 bits, so it will be $\$ 31$ (note the use of $\$$ sign to indicate register).
13. NOTE: Need to ensure that ALU operation results in ALUresult $=0$ to force is0? $=\mathbf{1}$. This can be guaranteed by having sw \$? ? 0(\$zero). Note that \$? ? can be any register. This is a branch, since the PCSrc is set to 1 , but it branch to the next instruction.
