

## Short Test (Exercise) for CS5232(appx: 60min)

Surname: \_\_\_\_\_ Given Name: \_\_\_\_\_

Matriculation Number: \_\_\_\_\_

**Note: Write your answers in the blank spaces in this answer book only.**

Q1. Decide if the following predicates are true or false (no need to give a reason). Note that  $\mathbb{N}$  is the natural number set starting from 0 and  $\mathbb{N}_1$  is the natural number set starting from 1.

(a)  $\forall n : \mathbb{N}_1 \bullet \exists m, k : \mathbb{N} \bullet nk = m \wedge n < m$

(b)  $\exists n : \mathbb{N} \bullet \forall m : \mathbb{N} \bullet m > n \Rightarrow \exists k : \mathbb{N} \bullet m = 2k$

(c)  $\forall S, T : \mathbb{P}\mathbb{N} \bullet S \subseteq T \vee T \subseteq S$

(d)  $\mathbb{P}\mathbb{P}\mathbb{P}\emptyset = \{\emptyset, \{\emptyset\}, \{\emptyset, \{\emptyset\}\}, \{\{\emptyset\}\}\}$

Q2. Let  $S$  be the set of numbers from 1 to 12 inclusive.

(a) Give a formal definition (in the declaration/predicate style) of the relation  $R : S \leftrightarrow S$  where  $x$  is related to  $y$  exactly when  $y$  is greater than the square of  $x$  but less than the square of  $x + 1$ .

(b) Write down  $R$  as a set of ordered pairs.

Q3. A function *squash* that takes any finite function  $f : \mathbb{N} \rightarrow A$  (for arbitrary set  $A$ ) and converts it into an element in  $\text{seq } A$  by squashing the domain to  $1 \dots \#f$ . The formal definition of *squash* is given as:

$$FN == \{f : \mathbb{N} \rightarrow A \mid f \text{ is finite}\}$$

$$\frac{\text{squash} : FN \rightarrow \text{seq } A}{\forall f : FN \bullet \begin{array}{l} f = \emptyset \Rightarrow \text{squash}(f) = \langle \rangle \\ f \neq \emptyset \Rightarrow \\ \quad \text{let } m = \min \text{ dom } f \bullet \\ \quad \quad \text{squash}(f) = \langle f(m) \rangle \hat{\wedge} \text{squash}(\{m\} \triangleleft f) \end{array}}$$

(a) Write down the value (result) of  $\text{squash}(\{(4, a), (2, b), (3, c)\})$ .

(b) Specify a function *remove* which, given an element in  $A$  and a sequence in  $\text{seq } A$ , removes from that sequence all occurrences of the element. (hint: using the function *squash*.)

Q4. Let  $[People]$  be the set of all possible persons and  $[Subjects]$  the set of all possible subjects. A specification for an *University* has a state schema:

$$\frac{\text{University}}{\begin{array}{l} \text{students} : \mathbb{P} \text{ People} \\ \text{subjects} : \mathbb{P} \text{ Subjects} \\ \text{enrolments} : \text{People} \leftrightarrow \text{Subjects} \end{array}}{\begin{array}{l} \text{dom enrolments} \subseteq \text{students} \\ \text{ran enrolments} \subseteq \text{subjects} \end{array}}$$

For the class *University*, specify following operations. Write your answer on this page, directly in the space provided in the schema definition.

- (a) an operation schema *newStudent* such that a person (*pers?*) becomes a student:

*newStudent* \_\_\_\_\_

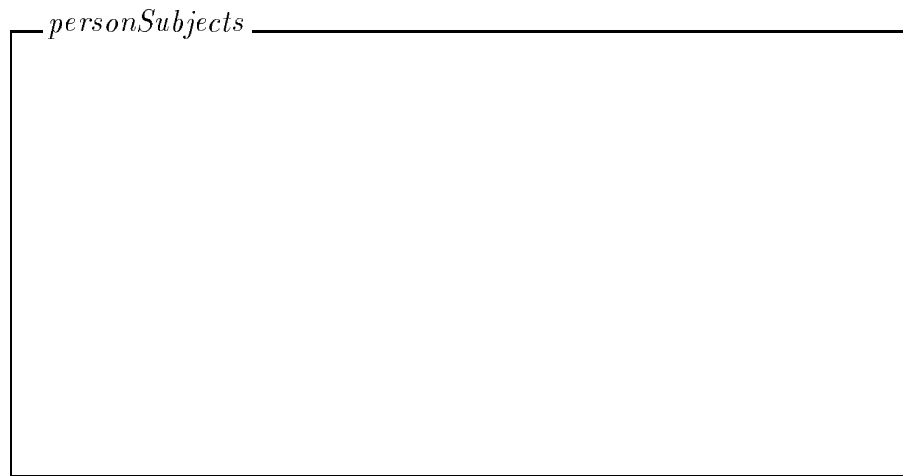
- (b) an operation schema *studentLeaves* such that a person (*pers?*), not enrolled in any subject, ceases to be a student:

*studentLeaves* \_\_\_\_\_

- (c) an operation schema *cancel* such that a person (*pers?*) cancels enrolment in a subject (*subj?*):

*cancel* \_\_\_\_\_

- (d) an operation schema  $personSubjects$  which outputs the set of subjects ( $subjs!$ ) in which the person ( $pers?$ ) is enrolled:



## Solution

Q1.

- (a) True.  
given any strictly positive  $n$ , let  $m = 2n$  and  $k = 2$ ; then  $nk = m$  and  $n < m$ .
- (b) False.  
take  $m = 2n + 1$ ; then  $m > n$  but as  $2n + 1$  is odd there is no  $k$  with  $m = 2k$ .
- (c) False.  
take  $S = \{1, 2\}$  and  $T = \{3, 4\}$ .
- (d) True.

Q2.

$$\left| \begin{array}{l} R : S \leftrightarrow S \\ \hline \forall x, y : \mathbb{N} \bullet \\ x \underline{R} y \Leftrightarrow y > x^2 \wedge y < (x + 1)^2 \end{array} \right.$$

$$R = \{(1, 2), (1, 3), (2, 5), (2, 6), (2, 7), (2, 8), (3, 10), (3, 11), (3, 12)\}$$

Q3.

- (a)  $\langle b, c, a \rangle$
- (b)

$$\frac{\text{remove} : A \times \text{seq } A \rightarrow \text{seq } A}{\forall a : A; s : \text{seq } A \bullet \text{remove}(a, s) = \text{squash}(s \triangleright \{a\})}$$

Q4.

(a)

$$\frac{\text{newStudent} \quad \Delta(\text{students}) \quad \text{pers?} : \text{People}}{\text{pers?} \notin \text{students} \quad \text{students}' = \text{students} \cup \{\text{pers?}\}}$$

(b)

$$\frac{\text{studentLeaves} \quad \Delta(\text{students}) \quad \text{pers?} : \text{People}}{\text{pers?} \in \text{students} \quad \text{pers?} \notin \text{dom enrolments} \quad \text{students}' = \text{students} \setminus \{\text{pers?}\}}$$

(c)

$$\frac{\text{cancel} \quad \Delta(\text{enrolments}) \quad \text{pers?} : \text{People} \quad \text{subj?} : \text{subjects}}{\text{pers?} \text{ enrolments } \text{subj?} \quad \text{enrolments}' = \text{enrolments} \setminus \{(\text{pers?}, \text{subj?})\}}$$

(d)

$$\frac{\text{personSubjects} \quad \text{pers?} : \text{People} \quad \text{subjs!} : \mathbb{P} \text{ subjects}}{\text{pers?} \in \text{students} \quad \text{subjs!} = \text{ran}(\{\text{pers?}\} \triangleleft \text{enrolments})}$$