

## MA 5219 - Logic and Foundations of Mathematics 1

Homework due in Week 3, Tuesday.

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Office hours Thursday 14.00-15.00h

Hand in each starred homework; 1 mark per homework (if it is correct), up to 10 marks in total for homework.

**3.1 Conjunctive and Disjunction Normalform.** Make formulas in conjunctive normal form and disjunctive normal form which state that exactly 2 of the atoms  $p_1, p_2, p_3, p_4$  are true.

**3.2\* Worlds.** Recall that a world is an entity which assigns a truth-value to every atom.  $W \models A$  means that the world makes the formula  $A$  true and  $W \models X$  means that the world makes all formulas in the set  $X$  of formulas true.  $X \models A$  means that every world which makes all formulas in  $X$  true, also makes the formula  $A$  true.

Let  $V$  and  $W$  be two different worlds, let  $X = \{A : V \models A \text{ or } W \models A\}$  and let  $Y = \{A \in X : \forall B \in X [A \wedge B \in X]\}$ . Show the following.

- (a)  $X \models A$  for every formula  $A$ .
- (b) There are formulas  $A, B \in X$  with  $A \wedge B \notin X$ .
- (c) If  $A, B, C \in X$  then at least one of the formulas  $A \wedge B, A \wedge C, B \wedge C$  is in  $X$ .
- (d) If  $Y \models A$  then  $A \in Y$ .

**3.3 Logical Implication.** (a) Assume that  $W \models X$  for all worlds  $W$  which make only finitely many atoms true. Show that  $W \models X$  for all worlds and that  $X$  contains only tautologies.

(b) Construct a set  $X$  of formulas such that  $W \models X$  is true iff  $W$  makes at most two atoms true and all others false.

**3.4 Proof System.** Assume that only  $(A + B)$  is permitted to connect formulas  $A, B$ , which are built from the atoms  $p_0, p_1, \dots$  and the logical constants 0 and 1. Is there a set of rules which permits to prove  $A$  from  $X$  whenever  $X \models A$  and  $X$  is a set of formulas of the above form and  $A$  is a formula of the above form? If so, give the set of rules; if not, explain why a set of such rules cannot exist.