Theory of Computation Additional Exercises

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Additional Exercises A–B

Each student can write up one of these exercises in the Forum to get 2 additional marks; please check before writing that no other student already wrote it up in the Forum. Maximum marks for Exercises is 10.

Exercise A.

Provide a regular expression for all words which have either the subword 110 or the subword 120 over the ternary alphabet $\{0, 1, 2\}$.

Exercise B.

Use structural induction to define $L^{mi} = \{w^{mi} : w \in L\}$ for all regular sets L. So first define this for all sets containing words up to length 1 and then explain, how one goes on with concatenation, union, Kleene star and Kleene plus. Sets consisting of words of length at least 2 can be obtained from the base cases this way.

Additional Exercises C–E

Exercise C.

Provide a DFA accepting those binary numbers which are multiples of 3: 11, 110, 1001, 1100, 1111, 10010, 10101, 11000, 11011 and so on.

Exercise D.

Construct over the ternary alphabet $\{0, 1, 2\}$ a DFA which accepts all words which contain each of the digits an odd number of times.

Exercise E.

Provide a non-deterministic finite automaton of thirteen states which accepts all decimal numbers which are not a multiple of 210; for this note that 0 is also a multiple of 210 and should be rejected. The numbers accepted should not having leading zeroes.

Additional Exercises F–H

Exercise F.

Construct a context-free grammar in Chomsky Normal Form for the language $\{0^n 1^m 2^n : n, m \in \mathbb{N}\}$.

Exercise G.

Construct a context-free grammar in Greibach Normal Form for the language $\{0^n 1^m 2^n : n, m \in \mathbb{N}\}$.

Exercise H.

Use the grammar from F and the Cocke Kasami Younger algorithm to check whether F generates the following words: 00122 and 00112.

Additional Exercises I–K

Exercise I.

Recall that a language L satisfies the weakest form of the Pumping Lemma iff there is a constant k such that all words of length at least k in L can be split into parts xyz with $y \neq \varepsilon$ and $\{x\} \cdot \{y\}^* \cdot \{z\} \subseteq L$. Which of the following choices for L satisfy this pumping lemma:

- 1. L = $\{0^n 1^m 2^n : n, m \in \mathbb{N}\};$
- 2. $\mathbf{L} = \{\mathbf{0^n1^m0^n}: \mathbf{n}, \mathbf{m} \in \mathbb{N}\};$
- 3. L = $\{0^n 1^m 2^k : n + k \neq m\}$?

Exercise J.

Which of the languages in Exercise I have a linear grammar?

Exercise K.

Use Ogden's Lemma to prove that the language $\{0^n1^m2^k : n \neq m \land n \neq k \land m \neq k\}$ is not context-free.

Additional Exercises L–O

Let h, k be homomorphisms with h(a) = a, $h(b) = \varepsilon$, $k(a) = \varepsilon$ and k(b) = b for $a \in \{0, 1, 2, 3, 4\}$ and $b \in \{5, 6, 7, 8, 9\}$. Prove the following statements for $L \subseteq \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}^*$.

Exercise L. If $\mathbf{k}(\mathbf{L})$, $\mathbf{h}(\mathbf{L})$ are both regular then \mathbf{L} is regular.

Exercise M. If $\mathbf{k}(\mathbf{L})$ is deterministic context-free and $\mathbf{h}(\mathbf{L})$ regular then \mathbf{L} is deterministic context-free.

Exercise N. There is an L for which h(L), k(L) are both context-free but L is not.

Exercise O. If $\mathbf{k}(\mathbf{L})$, $\mathbf{h}(\mathbf{L})$ are context-sensitive languages not containing ε , so is \mathbf{L} .

Additional Exercises P–S

Do for $\mathbf{L} = {\mathbf{w} \in {000, 111, 222}^+ : \mathbf{w} \text{ is a palindrome}}$ the following exercises.

Exercise P. Which is the least constant k such that every word $w \in L$ can be split into three parts x, y, z with w = xyz and $1 \leq |y| \leq k$ and $xy^*z \subseteq L$. Give the answer ∞ if there is no such constant k. Prove the answer.

Exercise Q. Provide a context-free grammar in Chomsky Normal Form for L.

Exercise R. Provide a context-free grammar in Greibach Normal Form for L.

Exercise S. Provide a PDA accepting by state for L. Can this PDA be made deterministic? Give a short reason for the answer.

Additional Exercises T–W

Prove that the following functions are primitive recursive.

- Exercise T. Function $f(n) = n^2$.
- Exercise U. Function $g(n) = n^n$.
- Exercise V. Function $h(m, n) = \binom{m+n}{m}$.
- Exercise W. Function $\mathbf{k}(\mathbf{m}, \mathbf{n}) = \mathbf{m}! + \mathbf{n}!$.

Additional Exercises X-Z

Exercise X. Is it decidable to check whether a polynomial with integer coefficients and one input variable x takes on some input the value 1024? For example, for input function $f(x) = x^2 + 1$, one wants to check whether there is an integer x with f(x) = 1024. Prove the answer.

Exercise Y. Let $\varphi_0, \varphi_1, \ldots$ be an acceptable numbering of all partial-recursive functions. Is $\mathbf{L} = \{\mathbf{e} : \varphi_{\mathbf{e}}(\mathbf{x}) \text{ is undefined for some } \mathbf{x}\}$ recursively enumerable? Prove the answer using Rice's Theorem.

Exercise Z. Let $\varphi_0, \varphi_1, \ldots$ be an acceptable numbering of all partial-recursive functions. Is $\mathbf{H} = \{\mathbf{e} : \text{the domain of } \varphi_{\mathbf{e}} \text{ is the range of a primitive recursive function}\}$ (a) decidable, (b) recursively enumerable and undecidable, (c) not recursively enumerable? Prove the answer using Rice's Theorem.