07—Application of SAT Solving

CS 3234: Logic and Formal Systems

Martin Henz

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Propositional Logic: Application of SAT Solving

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- there should be at least 7 dates distance between first leg and return match.
- To achieve this, we assume a fixed mirroring between dates: (1,8), (2,9), (3,12), (4,13), (5,14), (6,15) (7,16), (10,17), (11,18)

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- Every team except FSU has a traditional rival. The rival pairs are Clem-GT, Duke-UNC, UMD-UVA and NCSt-Wake. In the last date, every team except FSU plays against its rival, unless it plays against FSU or has a bye.

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- UNC plays Clem in the second date.

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- UNC plays Duke in last date and date 11.
- UNC plays Clem in the second date.
- Duke has bye in the first date 16.

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- Wake does not play home in date 17.
- Wake has a bye in the first date.
- Clem, Duke, UMD and Wake do not play away in the last date.
- Clem, FSU, GT and Wake do not play away in the fist date.

- Wake does not play home in date 17.
- Wake has a bye in the first date.
- Clem, Duke, UMD and Wake do not play away in the last date.
- Clem, FSU, GT and Wake do not play away in the fist date.
- Neither FSU nor NCSt have a bye in the last date.

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- Neither FSU nor NCSt have a bye in the last date.
- UNC does not have a bye in the first date.

Background

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- From then onwards, Henz, Walser and Zhang use different techniques to solve the problem

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 - Generate schedules from pattern sets
- Output: all feasible solutions, from which the organizers can choose the most suitable one

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- Zhang Hantao uses SAT solving, turn-around time of 2 seconds, see "Generating College Conference Basketball Schedules using a SAT Solver"
- Different approach: In 1998, J.P. Walser described a local-search based method for finding some (not all) solutions, without using 3 phases



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- For every pair of distinct teams *s* and *t*, we have:

$$(p_{s,t,1} \wedge \neg p_{s,t,2} \wedge \cdots \wedge \neg p_{s,t,18}) \vee (\neg p_{s,t,1} \wedge p_{s,t,2} \wedge \neg p_{s,t,3} \wedge \cdots \wedge \neg p_{s,t,18}) \vee \\ \vdots \\ (\neg p_{s,t,1} \cdots \wedge \neg p_{s,t,17} \wedge p_{s,t,18})$$

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Convert formula into CNF, and use a complete SAT solver



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- Phase 2: 38 · 9 · 3 = 1026 propositional atoms, 569300 clauses, taking 0.60 seconds, resulting in 17 pattern sets
- Phase 3: 9 · 9 + 9 · 8 · 18 = 1377 propositional atoms, hundreds of thousands of clauses, taking less than 2 seconds, resulting in 179 solutions

Conclusion

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- The approach takes advantage of the effort that the designers of SAT solvers such as SATO spent in order to optimize the solver.
- This works well, because the solver is independent of the application domain; it can be used without modification across application domains.