

**CS1102 Semester 2 AY 2003/2003**  
**Annotated Solutions for Lab 4A**  
**Problem 40: Augmented Binary Search Tree**

**Introduction**

This problem is meant to be a easy problem, that could be completed within two hours.

To solve this problem, you need to understand how the given code works. Drawing pictures of BSTs, and thinking about how insertion/deletion can change the five data fields are important. Trying to solve this problem straight away before understanding the code or how insert/delete can affect the tree would cause wasted hours of debugging. It also requires knowledge of how to delete and insert from a doubly linked list.

The given source code is shown in the appendix. The unrelated part of the source code has been removed for simplicity. Therefore the line numbers referred to in this document does not match those in the java file. Let's look at how each data field can be affected by insertion and deletion. We will begin by looking at the simplest one, the parent.

**Parent**

The parent of a node N should either point to itself, if N is a root, or points to the parent. Therefore, the parent of node N changes, if either N becomes a new root, or if some other node sets its left or right child to N.

When can a node becomes a new root? This happens when the member "root" of BinarySearchTree changed and only happens in line 2 to 10.

We can reset the parent pointer here:

```
2  public void insert(KeyedItem newItem)
3  {
4      root = insertItem(root, newItem);
      root.parent = root;
5  }
6
7  public void delete(Comparable searchKey) throws TreeException
8  {
9      root = deleteItem(root, searchKey);
      if (root != null) root.parent = root;
10 }
```

Note that we do not have to check if root is null after insertItem.

The other case where a parent pointer could change is when some other node has changed its child pointer. This happens in line 135 to 146.

We can update the parent pointer here:

```
135 public void setLeft(TreeNode left)
136 {
137     // Sets the left child reference to left.
138     leftChild = left;
      if (left != null) left.parent = this;
139 }
140
141 public void setRight(TreeNode right)
142 {
143     // Sets the right child reference to right.
144     rightChild = right;
```

```

        if (right != null) right.parent = this;
145     }
146 }

```

## Size and Height

Initialization of size and height are easy. When a new `TreeNode` is created, just set both of them to 1.

Since size and height of a node *N* indicate the number of nodes and the height of the subtree rooted at *N*, these two data fields need to be updated whenever the subtree of *N* is changed. The subtree rooted at a node *N* can only be changed when an item is inserted/deleted from one of its subtree. For insertion, this happens between line 24 and 34 after `insertItem()` is called on either the left or right subtree.

```

23     // search for the insertion position
24     if (newItem.getKey().compareTo(nodeItem.getKey()) < 0) {
25         // search the left subtree
26         newSubtree = insertItem(tNode.getLeft(), newItem);
27         tNode.setLeft(newSubtree);
28         tNode.updateSizeAndHeight();
29         return tNode;
30     }
31     else { // search the right subtree
32         newSubtree = insertItem(tNode.getRight(), newItem);
33         tNode.setRight(newSubtree);
34         tNode.updateSizeAndHeight();
35         return tNode;
36     }

```

For deletion, this happens in four different places, whenever `deleteItem` or `deleteLeftmost` is called.

```

49     // else search for the item
50     else if (searchKey.compareTo(nodeItem.getKey()) < 0) {
51         // search the left subtree
52         newSubtree = deleteItem(tNode.getLeft(), searchKey);
53         tNode.setLeft(newSubtree);
54         tNode.updateSizeAndHeight();
55     }
56     else { // search the right subtree
57         newSubtree = deleteItem(tNode.getRight(), searchKey);
58         tNode.setRight(newSubtree);
59         tNode.updateSizeAndHeight();
60     }
61 }
62
63 // retrieve and delete the inorder successor
64 else {
65     replacementItem = findLeftmost(tNode.getRight());
66     tNode.setItem(replacementItem);
67     tNode.setRight(deleteLeftmost(tNode.getRight()));
68     tNode.updateSizeAndHeight();
69 }
70
71 return tNode;
72 }
73 }
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99
100 protected TreeNode deleteLeftmost(TreeNode tNode)
101 {
102     if (tNode.getLeft() == null) {
103         return tNode.getRight();
104     }
105     else {
106         tNode.setLeft(deleteLeftmost(tNode.getLeft()));
107         tNode.updateSizeAndHeight();
108     }
109 }

```

```

106     return tNode;
107 }
108 }
109

```

Method `updateSizeAndHeight()` can be written as a method in class `TreeNode` as follows:

```

private void updateSizeAndHeight()
{
    if (leftChild == null && rightChild == null) {
        size = height = 1;
    } else if (leftChild == null) {
        size = rightChild.size + 1;
        height = rightChild.height + 1;
    } else if (rightChild == null) {
        size = leftChild.size + 1;
        height = leftChild.height + 1;
    } else {
        size = leftChild.size + rightChild.size + 1;
        height = Math.max(leftChild.height, rightChild.height) + 1;
    }
}
}

```

Someone observed that we can update the size and height in `setLeft()` and `setRight()`. This is correct for this particular implementation. However, for other implementation, it might not be the case that any modification to the subtree rooted at node `N` results in a change to the left or right child of `N`.

### Predecessor and Successor

The predecessor and successor fields are not related to the structure of the tree, but to the values of the items stored in the tree. An important observation is that these two references form a doubly linked list. Thus, updates of predecessor and successor when insert/delete is performed on the tree are the same as updates of next/prev references in a doubly linked list.

To insert a new item `N`, we first need to find out the position of `N` in the doubly linked list. This is easy once you made the observation that the parent of `N` must be either the successor (if it is inserted to the left) or the predecessor (if inserted to the right) of `N`. The next question is, when to update the predecessor and successor? Obviously this must be done when a new node is created, that is, after line 18. But at line 18, we have no access to the parent of the newly created `tNode`. We can solve this by passing in the parent node to `insertItem`.

```

12     protected TreeNode insertItem(TreeNode tNode, KeyedItem newItem,
                                   TreeNode parent)
13     {
14         TreeNode newSubtree;
15         if (tNode == null) {
16             // position of insertion found; insert after leaf
17             // create a new node
18             tNode = new TreeNode(newItem, null, null);
19             // if parent is larger than tNode,
20             //     insert tNode before parent
21             // else
22             //     insert tNode after parent
23             return tNode;
24         }
25     }

```

Finally, we consider how to update successor and predecessor when a node `N` has been deleted. Deletion of a node is done in `deleteNode`, in which the argument `tNode` is removed from the tree. Case 1, 2 and 3 are straight forward – since we are removing `tNode` from the tree, we just remove it from the double linked list formed by the successor and predecessor reference. In case 4, `tNode` is not deleted, but instead, was replaced by its successor. Hence, we do not remove `tNode`, but we remove its successor instead.

```

63  protected TreeNode deleteNode(TreeNode tNode)
64  {
65      // Algorithm note: There are four cases to consider:
66      // 1. The tNode is a leaf.
67      // 2. The tNode has no left child.
68      // 3. The tNode has no right child.
69      // 4. The tNode has two children.
70      // Calls: findLeftmost and deleteLeftmost
71      KeyedItem replacementItem;
72
73      // test for a leaf
74      if ( (tNode.getLeft() == null) &&
75          (tNode.getRight() == null) ) {
76          // remove tNode from doubly linked list
77          return null;
78      }
79      // test for no left child
80      else if (tNode.getLeft() == null) {
81          // remove tNode from doubly linked list
82          return tNode.getRight();
83      }
84      // test for no right child
85      else if (tNode.getRight() == null) {
86          // remove tNode from doubly linked list
87          return tNode.getLeft();
88      }
89      // there are two children:
90      // retrieve and delete the inorder successor
91      else {
92          replacementItem = findLeftmost(tNode.getRight());
93          tNode.setItem(replacementItem);
94          tNode.setRight(deleteLeftmost(tNode.getRight()));
95          // remove tNode's successor from doubly linked list
96          return tNode;
97      }

```

One interesting note is that we can simplify the code for case 4 in deletion. We do not really need to traverse the tree to find the replacement item using `findLeftmost()` anymore, because the successor of `tNode` points to the replacement item directly.

### Appendix: Original Source Code for Binary Search Tree.

```

1  class BinarySearchTree extends BinaryTreeBasis {
2      public void insert(KeyedItem newItem)
3      {
4          root = insertItem(root, newItem);
5      }
6
7      public void delete(Comparable searchKey) throws TreeException
8      {
9          root = deleteItem(root, searchKey);
10     }
11
12     protected TreeNode insertItem(TreeNode tNode, KeyedItem newItem)
13     {
14         TreeNode newSubtree;
15         if (tNode == null) {
16             // position of insertion found; insert after leaf
17             // create a new node
18             tNode = new TreeNode(newItem, null, null);
19             return tNode;

```

```

20     }
21     KeyedItem nodeItem = (KeyedItem)tNode.getItem();
22
23     // search for the insertion position
24     if (newItem.getKey().compareTo(nodeItem.getKey()) < 0) {
25         // search the left subtree
26         newSubtree = insertItem(tNode.getLeft(), newItem);
27         tNode.setLeft(newSubtree);
28         return tNode;
29     }
30     else { // search the right subtree
31         newSubtree = insertItem(tNode.getRight(), newItem);
32         tNode.setRight(newSubtree);
33         return tNode;
34     }
35 }
36
37 protected TreeNode deleteItem(TreeNode tNode, Comparable searchKey)
38 {
39     TreeNode newSubtree;
40     if (tNode == null) {
41         throw new TreeException("TreeException: Item not found");
42     }
43     else {
44         KeyedItem nodeItem = (KeyedItem)tNode.getItem();
45         if (searchKey.compareTo(nodeItem.getKey()) == 0) {
46             // item is in the root of some subtree
47             tNode = deleteNode(tNode); // delete the item
48         }
49         // else search for the item
50         else if (searchKey.compareTo(nodeItem.getKey()) < 0) {
51             // search the left subtree
52             newSubtree = deleteItem(tNode.getLeft(), searchKey);
53             tNode.setLeft(newSubtree);
54         }
55         else { // search the right subtree
56             newSubtree = deleteItem(tNode.getRight(), searchKey);
57             tNode.setRight(newSubtree);
58         }
59     }
60     return tNode;
61 }
62
63 protected TreeNode deleteNode(TreeNode tNode)
64 {
65     // Algorithm note: There are four cases to consider:
66     // 1. The tNode is a leaf.
67     // 2. The tNode has no left child.
68     // 3. The tNode has no right child.
69     // 4. The tNode has two children.
70     // Calls: findLeftmost and deleteLeftmost
71     KeyedItem replacementItem;
72
73     // test for a leaf
74     if ( (tNode.getLeft() == null) &&
75         (tNode.getRight() == null) ) {
76         return null;
77     }
78
79     // test for no left child
80     else if (tNode.getLeft() == null) {
81         return tNode.getRight();
82     }
83
84     // test for no right child
85     else if (tNode.getRight() == null) {
86         return tNode.getLeft();
87     }
88

```

```

89     // there are two children:
90     // retrieve and delete the inorder successor
91     else {
92         replacementItem = findLeftmost(tNode.getRight());
93         tNode.setItem(replacementItem);
94         tNode.setRight(deleteLeftmost(tNode.getRight()));
95         return tNode;
96     }
97 }
98
99 protected TreeNode deleteLeftmost(TreeNode tNode)
100 {
101     if (tNode.getLeft() == null) {
102         return tNode.getRight();
103     }
104     else {
105         tNode.setLeft(deleteLeftmost(tNode.getLeft()));
106         return tNode;
107     }
108 }
109
110 }
111
112
113 class TreeNode {
114     private Object item;
115     private TreeNode leftChild;
116     private TreeNode rightChild;
117
118     public TreeNode(Object newItem)
119     {
120         // Initializes tree node with item and no children.
121         item = newItem;
122         leftChild = null;
123         rightChild = null;
124     }
125
126     public TreeNode(Object newItem, TreeNode left, TreeNode right)
127     {
128         // Initializes tree node with item and
129         // the left and right children references.
130         item = newItem;
131         leftChild = left;
132         rightChild = right;
133     }
134
135     public void setLeft(TreeNode left)
136     {
137         // Sets the left child reference to left.
138         leftChild = left;
139     }
140
141     public void setRight(TreeNode right)
142     {
143         // Sets the right child reference to right.
144         rightChild = right;
145     }
146 }

```