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**Q1**: Let  $\text{Infinite}_{\mathsf{DFA}} = \{ \langle A \rangle : A \text{ is a DFA and } L(A) \text{ is an infinite language} \}$ . Show that Infinite\_{\mathsf{DFA}} is decidable.

**Q2**: Let  $A = \{\langle M \rangle : M \text{ is a DFA which doesn't accept any string containing an odd number of 1s}. Show that A is decidable.$ 

**Q3**: Let A and B be two disjoint languages. Say that language C separates A and B if  $A \subseteq C$  and  $B \subseteq \overline{C}$ . Show that any two disjoint co-Turing-recognizable languages are separable by some decidable language.

**Q4**: Let  $\text{PAL}_{\mathsf{DFA}} = \{ \langle M \rangle : M \text{ is a DFA that accepts some palindrome} \}$ . Show that  $\text{PAL}_{\mathsf{DFA}}$  is decidable.

**Q5**: Let A be a Turing-recognizable language consisting of descriptions of Turing machines  $\{\langle M_1 \rangle, \langle M_2 \rangle, \ldots\}$ , where every  $M_i$  is a decider. Prove that some decidable language D is not decided by any decider  $M_i$  whose description appears in A.

**Q6**: Let *B* be a Turing-recognizable language consisting of TM descriptions. Show that there is a decidable language *C* consisting of TM descriptions such that every machine described in *B* has an equivalent machine in *C* and vice versa.

**Q7**: Show that  $\{\langle G \rangle : G \text{ is a } \mathsf{CFG} \text{ over } \{0,1\} \text{ and } 1^* \subseteq L(G) \}$  is decidable.