Tutorial 5 (and 6)

Show that the following problems are NP-complete.

(a) Vertex Cover

Input: a (undirected) graph (V, E), and a number k.

Question: Is there a subset V' of V such that V' has at most k elements, and for every edge (u, v) in E, at least one of u, v belongs to V'?

(b) Clique

Input: a (undirected) graph (V, E) and a number k.

Question: Is there a subset V' of V of size k such that for each distinct u, v in V', (u, v) is in E?

(c) X3C

Input: A finite set A, with |A| = 3q, and a collection C of 3-element subsets of A.

Question: Does C contain an exact cover for A, that is, a subcollection $C' \subseteq C$ such that every element of A occurs in exactly one member of C'.

(d) 3-Colorability.

Input: An undirected graph G = (V, E).

Question: Is there a mapping $color: V \to \{1, 2, 3\}$ such that for all $(v, w) \in E$, $color(v) \neq color(w)$?

(e) Not-All-Equal SAT (NAESAT).

Input: A set of variables V, and a set C of clauses, each having exactly three literals.

Question: Is there a truth assignment to the variables so that each clause has at least one true literal and at least one false literal?

(f) MAX2SAT

Input: A set of variables V, a number k and a set of clauses C, where each clause has at most two literals.

Question: Is there a truth assignment to the variables such that at least k of the clauses become true?

(g) Hamiltonian Circuit

Input: A graph G(=(V, E))

Output: Does G contain a Hamiltonian Circuit, that is, an ordering (v_1, \ldots, v_n) of the vertices in V such that $(v_i, v_{i+1}) \in E$, for $1 \leq i < n$, and $(v_n, v_1) \in E$.