

Relational Algebra Equivalences

- Allow us to choose different join orders and to `push` selections and projections ahead of joins.
- Rules on joins, cross products and union

$$R \bowtie S = S \bowtie R$$
$$(R \bowtie S) \bowtie T = R \bowtie (S \bowtie T)$$

$$R \times S = S \times R$$

$$(R \times S) \times T = R \times (S \times T)$$

$$R \cup S = S \cup R$$

$$R \cup (S \cup T) = (R \cup S) \cup T$$

Rules: Selects

$$\sigma_{p_1 \wedge p_2}(R) = \sigma_{p_1} [\sigma_{p_2}(R)]$$

$$\sigma_{p_1 \vee p_2}(R) = [\sigma_{p_1}(R)] \cup [\sigma_{p_2}(R)]$$

Rules: Project

Let: X = set of attributes

Y = set of attributes

$XY = X \cup Y$

$$\pi_{xy}(R) = \pi_x[\pi_y(R)]$$

Rules: Project

Let: X = set of attributes

Y = set of attributes

$XY = X \cup Y$

$$\pi_{xy}(R) = \pi_x[\pi_y(R)]$$

$$\pi_x(R) = \pi_x[\pi_y(R)] \text{ if } y \text{ contains } x$$

Rules: $\sigma + \bowtie$ combined

Let P = predicate with only R attribs

Q = predicate with only S attribs

M = predicate with only R,S attribs

$$\sigma_p (R \bowtie S) = [\sigma_p(R)] \bowtie S$$

$$\sigma_q (R \bowtie S) = R \bowtie [\sigma_q(R)]$$

More Rules

Let x = subset of R attributes

z = attributes in predicate P
(subset of R attributes)

$$\pi_x[\sigma_p(R)] = \pi_x \{ \sigma_p [\overset{\pi_{xz}}{\cancel{\pi_x}}(R)] \}$$

More Rules

Let x = subset of R attributes

y = subset of S attributes

z = intersection of R,S attributes

$\pi_{xy} (R \bowtie S) =$

$$\pi_{xy} \{ [\pi_{xz} (R)] \bowtie [\pi_{yz} (S)] \}$$

More Rules

$\pi_{xy} \{ \sigma_P (R \bowtie S) \} =$

$$\pi_{xy} \{ \sigma_P [\pi_{xz'} (R) \bowtie \pi_{yz'} (S)] \}$$

$$z' = z \cup \{ \text{attributes used in } P \}$$