

# UIT2206 The Importance of Being Formal

Midterm, 19/3/2014

This test has nine questions on eight pages. Please answer all questions in the spaces provided after each (sub)question.

Enter your matriculation number here:

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## Traditional Logic

**Question 1:** (6 marks)

(a) Consider the following categorical proposition:

**No Mermaids are Mermaids.**

Give a model in which this proposition holds or a reason why there is no such a model.

(b) Consider the following categorical proposition:

**Some Werewolves are not Werewolves.**

Give a model in which this proposition holds or a reason why there is no such a model.

**Question 2:** (6 marks) Consider the following six categorical propositions:

- Some Men are Rich.
- Some Men are not Rich.
- Some Rich are Happy.
- Some Rich are not Happy.
- Some Men are Happy.
- Some Men are not Happy.

Give a model in which all six propositions hold.



## Propositional Logic

**Question 4:** (5 marks) Indicate which of the following formulas are valid. You do not need to give reasons for your judgment.

(a)  $\neg\top \vee \top$

(b)  $\neg\perp \wedge \perp$

(c)  $\neg\top \wedge \top$

(d)  $\neg\top \rightarrow \perp$

(e)  $\top \rightarrow \neg\top$

**Question 5:** (4 marks)

Is every invalid formula unsatisfiable? Give a short explanation of your answer, based on the definition of the adjectives involved.

**Question 6:** (5 marks) Assume the following entailment holds:

$$\phi \models \psi$$

Which one of the following statements is correct? You do not need to give reasons for your judgment.

(a) If  $\psi$  is unsatisfiable, then  $\phi$  is unsatisfiable.

(b) If  $\psi$  is valid, then  $\phi$  is valid.

(c) If  $\phi$  is unsatisfiable, then  $\psi$  is unsatisfiable.

(d) If  $\phi$  is valid, then  $\psi$  is valid.

(e) If  $\phi$  is invalid, then  $\psi$  is unsatisfiable.

## Predicate Logic

**Question 7:** (8 marks) Let  $P$  be a predicate symbol of arity 0,  $Q$  be a predicate symbol of arity 2 and  $f$  be a function symbol of arity 2. Consider the formula

$$\phi = \exists w((P \vee (\forall v Q(f(x, v), w))) \wedge (\forall z Q(f(w, x), z)))$$

(a) Draw the parse tree of  $\phi$  and indicate the free and bound variables in this parse tree.

(b) Give the result of the substitution  $[x \Rightarrow f(x, x)]\phi$ .

**Question 8:** (4 marks) Consider the following model  $\mathcal{M}$  for predicate logic with the binary relation  $R$ .

- Let  $U^{\mathcal{M}} = \{a, b, c, d\}$
- Let  $R^{\mathcal{M}} = \{(a, b), (b, c), (c, d), (d, a)\}$

For the following formulas  $\phi$ , indicate (without explanation), whether  $\mathcal{M} \models \phi$  or  $\mathcal{M} \not\models \phi$ .

- (a)  $\exists x(R(x, x))$
- (b)  $\forall x(\neg R(x, x))$
- (c)  $\exists x(\forall y(R(x, y)))$
- (d)  $\forall x(\exists y(R(x, y)))$

**Question 9:** (4 marks) In predicate logic, the universe of a model cannot be empty. If we remove this restriction, the validity of some formulas changes. Give an example for such a formula and a short explanation.