

## MA 3219 – Computability Theory

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**Assignment for 06.04.2005.** Can be corrected on request, it is not obligatory to hand the homework in.

**1. Many-One Reduction.** Assume that  $A \leq_m C$  and  $B \equiv_m C$ . Which of these properties are then inherited to  $A$  or to  $B$  or to both:

- (a)  $C$  is recursively enumerable;
- (b)  $C$  is simple;
- (c)  $C$  is creative;
- (d)  $C$  is semirecursive;
- (e)  $C$  is recursive.

Here the set  $C$  is semirecursive iff there is a recursive linear ordering  $\sqsubset$  such that whenever  $x \sqsubset y$  and  $x \in C$  then also  $y \in C$ .

**2. Turing Reduction.** Which two of the following sets are Turing equivalent? Prove this equivalence:

- (a)  $\emptyset$ ;
- (b)  $\{e : \phi_e(2345) \downarrow = 8\}$ ;
- (c)  $\{e : \phi_e(0) \uparrow \vee \phi_e(1) \uparrow\}$ ;
- (d)  $\{e : \phi_e \text{ is total}\}$ ;
- (e)  $\{e : \phi_e \text{ is infinitely often undefined}\}$ .

**3. Special Cases.** Prove that  $\emptyset$ ,  $\mathbb{E}$  and  $\mathbb{N}$  Turing equivalent but not many-one equivalent;  $\mathbb{E}$  is the set of even natural numbers.

**4. Classifying Reductions.** Which of the following reductions are many-one reductions and Turing reductions; if a reduction is both then state that it is a many-one reductions:

- (a)  $x \in A \Leftrightarrow x \notin B$ ;
- (b)  $x \in A \Leftrightarrow 2x \in B \wedge 2x + 1 \notin B$ ;
- (c)  $x \in A \Leftrightarrow 256^x \in B$ ;
- (d)  $x \in A \Leftrightarrow \forall y (2^x \cdot 3^y \in B)$ ;
- (e)  $x \in A \Leftrightarrow y$  is even for the least  $y$  with  $2^x \cdot 3^y \in B$ .

At (e) it is assumed that for every  $x$  there is a  $y$  with  $2^x \cdot 3^y \in B$ .

**5. Growth of functions.** Let  $A, B$  be r.e. sets and assume that for every total  $A$ -recursive function  $f$  there is a total  $B$ -recursive function  $g$  such that  $\forall x (f(x) \leq g(x))$ . Prove that then  $A \leq_T B$ .