Selection by Recursively Enumerable Sets*

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Abstract. For given sets A, B and Z of natural numbers where the members of Z are $z_0, z_1, \ldots$ in ascending order, one says that A is selected from B by Z if $A(i) = B(z_i)$ for all i. Furthermore, say that A is selected from B if A is selected from B by some recursively enumerable set, and that A is selected from B in n steps iff there are sets $E_0, E_1, \ldots, E_n$ such that $E_0 = A$, $E_n = B$, and E_i is selected from E_{i+1} for each $i < n$. The following results on selections are obtained in the present paper. A set is ω-r.e. if and only if it can be selected from a recursive set in finitely many steps if and only if it can be selected from a recursive set in two steps. There is some Martin-Löf random set from which any ω-r.e. set can be selected in at most two steps, whereas no recursive set can be selected from a Martin-Löf random set in one step. Moreover, all sets selected from Chaitin's Ω in finitely many steps are Martin-Löf random.

1 Introduction

Post [13] introduced various important reducibilities in recursion theory among which the one-one reducibility is the strictest one; here A is one-one reducible to B iff there is a one-one recursive function F such that $A(x) = B(F(x))$ for all x. In a setting of closed left-r.e. sets, Jain, Stephan and Teutsch [2] investigated a strengthening of one-one reductions were it is required in addition that F is strictly increasing or, equivalently, that F is the principal function of an infinite recursive set. The present paper relaxes the latter notion of reducibility and considers reductions given by principal functions of infinite sets that are recursively enumerable (r.e., for

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short). Recall that the principal function of an infinite set Z is the strictly increasing function F such that Z can be written as $\{F(0), F(1), \ldots\}$. In case for such Z and F some set A is reduced to some set B in the sense that $A(x) = B(F(x))$ for all x, this reduction could also be viewed as retrieving A from the asymmetric join B where the two "halves" of the join are not coded at the even and odd positions, respectively, of the join as usual but are coded into the positions that correspond to members and nonmembers, respectively, of the set Z.

Definition 1. A set A is selected from a set B by a set Z if Z is infinite and for the principal function F of Z it holds that $A(i) = B(F(i))$ for all i. A set A is selected from a set B, $A \sqsubset B$ for short, if A is selected from B by some r.e. set. Furthermore, say that A is selected from B in n steps iff there are sets $E_0, E_1, \ldots, E_n$ such that $E_0 = A$, $E_n = B$ and E_i is selected from E_{i+1} for each $i < n$.

The set B has selection rank n, if n is the maximum number such that some set A can be selected from B in n steps but not in $n - 1$ steps.

It makes sense to consider selection in several steps as the selection relation is not transitive: it follows by Theorems 14 and 15 that there is a Martin-Löf random set from which every recursive set can be selected in two steps but not in one step.

Note that for any infinite set Z, the principal function F of Z depends uniquely on Z. Furthermore, for a selection of A from B in n steps as in Definition 1, where F_m selects E_m from E_{m+1}, one can easily see that the function $\widetilde{F}$ given by

$$\widetilde{F}(y) = F_{n-1}(F_{n-2}(\ldots F_2(F_1(F_0(y)))\ldots))$$

satisfies that $A(i) = B(\widetilde{F}(i))$. However, since the selection relation is not transitive, in general, the range of $\widetilde{F}$ is not recursively enumerable and one cannot use the function $\widetilde{F}$ to select A from B in one step.

Early research in algorithmic randomness was formulated in terms of admissible selection rules. More precisely, given a certain way of selecting a subsequence from the characteristic sequence of a set, a set is called random iff all of its subsequences selected this way satisfy the condition in the law of large numbers that in the limit the frequencies of the symbols 0 and 1 are both equal to 1/2 [6]. Furthermore, van Lambalgen's Theorem [5] states that if one decomposes a set A by selecting along a recursive set and its complement into two infinite halves B_0 and B_1 then A is Martin-Löf random iff B_0 and B_1 are Martin-Löf random relative to each other. From this viewpoint it is natural to ask whether the choice of the selection along a recursive set can be generalised here to the choice along an r.e. set and how this is compatible with randomness notions. For notions whose definition involves the halting problem, in particular for Kurtz random relative to K, Schnorr random relative to K and Martin-Löf random relative to K, it can easily be shown that if B has one of these randomness properties and A is selected from B then A has also the same randomness property and that in the case of Martin-Löf randomness relative to K, one even gets the full equivalent of van Lambalgen's Theorem.

These initial and obvious connections ask for deeper investigation in order to see how far these correspondences go and thus, one of the central questions investigated in this paper is when random sets can be selected from extremely nonrandom ones and vice versa. Furthermore, the notion of ω-r.e. sets — which also play an important role in algorithmic randomness — fits well with the notion of selection by r.e. sets and strong connections are found. Hence, the present work aims at establishing some basic properties of the selection relation and at clarifying its interplay with other established recursion-theoretic notions like Martin-Löf randomness, immunity, enumeration-properties and initial segment complexity.

In the next section it is shown that a set is ω-r.e. if and only if it can be selected from an infinite and coinfinite recursive set in two steps. Every recursive set E has selection rank of at most 2, where the selection rank is 2 if and only if the set is infinite and coinfinite. Furthermore, the truth-table cylinder of the halting problem has selection rank 1. Every set selected from Ω in finitely many steps is Martin-Löf random (but differs from Ω in case at least one of the selections is nontrivial). There are Martin-Löf random sets which behave differently, for example, all ω-r.e. sets can be selected from some Martin-Löf random set.

2 Selection and ω-r.e. sets

Recall that a set A is ω-r.e. iff there is a recursive function f and a sequence of sets $A_0, A_1, \ldots$ such that the A_s form a recursive approximation to A where the number of mind changes is bounded by f, that is,

- the sets A_s are uniformly recursive, that is, the mapping $(x, s) \mapsto A_s(x)$ is a recursive function of two arguments;
- for all x and all sufficiently large stages s, $A(x) = A_s(x)$;
- $A_0 = \emptyset$ and for every x there at most $f(x)$ stages s with $A_s(x) \neq A_{s+1}(x)$.

Note that a set A is r.e. if and only if it is ω-r.e. with a bounding function f as above that satisfies the additional constraint that $f(x) = 1$ for all x. The r.e. sets and ω-r.e. sets have been well-studied in recursion theory [6, 10–12, 15].

The first result is that the ω-r.e. sets are closed downwards under the selection relation.

Theorem 2. *Assume that A is selected from B and B is an ω-r.e. set. Then A is an ω-r.e. set, too.*

Proof. Let the recursive approximation $B_0, B_1, \ldots$ and the recursive function f witness that B is an ω-r.e. set. Furthermore, let W be an r.e. set selecting A from B. There is a strictly increasing recursive function g such that its range is a recursive subset W_0 of W. Fix some recursive approximation $W_0, W_1, \ldots$ of W with $W_0 \subseteq W_1 \subseteq \ldots$ and let A_s be the set selected from B_s by W_s. Then $A_0 = \emptyset$ because $B_0 = \emptyset$. Furthermore, it can easily be seen that $A_s(n) \neq A_{s+1}(n)$ requires that there is an $x \leq g(n)$ with $W_s(x) \neq W_{s+1}(x)$ or $B_s(x) \neq B_{s+1}(x)$. Since for each such x these two conditions can be true for at most 1 and for at most $f(x)$ stages s, respectively, the total number of stages s where $A_s(n) \neq A_{s+1}(n)$ holds is at most $g(n)+1+f(0)+\ldots+f(g(n))$. Hence A is an ω-r.e. set. $\square$

Theorem 3. *Let O be the set of odd numbers and A be an ω-r.e. set. Then A can be selected from O in two steps, that is, there is a set B such that $A \sqsubset B$ and $B \sqsubset O$.*

Proof. Let the recursive approximation $A_0, A_1, \ldots$ and the recursive function f witness that A is an ω-r.e. set and let $h(n) = n^2 \cdot (1 + f(0) + f(1) + \ldots + f(n))$. It suffices to construct B such that, first, $A(n) = B(h(n))$, that is, the range of h selects A from B and, second, the set B is selected from O by some r.e. set W.

Split the natural numbers into consecutive intervals $I_0, J_0, I_1, J_1, I_2, J_2, \ldots$ where the length of I_n is $1 + f(0) + f(1) + \ldots + f(n)$ and where the length of J_n is $h(0) + 1$ for $n = 0$ and is $h(n) - h(n-1)$ for $n \geq 1$. Let W_0 be the union of all J_n and let $k_0 = 0$. During stages $s = 0, 1, \ldots$, one applies the following updates:

1. Let B_s be the set selected via W_s from O;
2. If $A_s(k_s) \neq B_s(h(k_s))$ then let $W_{s+1} = W_s \cup \{\max(I_{k_s} - W_s)\}$ else let $W_{s+1} = W_s$;
3. Let $k_{s+1} = \min(\{k_s + 1\} \cup \{j : A_s(j) \neq A_{s+1}(j)\})$.

Say a stage s is enumerating in case on reaching its second step the condition of the if-clause is satisfied. First it is shown that for every enumerating stage s, the set $I_{k_s} - W_s$ is nonempty, hence indeed the maximum member of this set is enumerated into W during stage s. Fix n and consider the enumerations of members of I_n into W in step 2. After each such enumeration at some stage s, any further such enumeration requires that at some stage $t > s$, one has $k_{t+1} \leq n < k_t$, which by step 3 in turn requires that the approximation to A has a mind change of the form $A_t(j) \neq A_{t+1}(j)$ where $j \leq n$ and $s < t$. Since for enumerations of distinct members of I_n there must be distinct such pairs (j, t), by choice of f at most $1 + f(0) + f(1) + \ldots + f(n) = |I_n|$ members of I_n are enumerated into W.

Next let $z_0, z_1, \ldots$ be the members of W in ascending order and for all s let $z_0^s, z_1^s, \ldots$ be the members of W_s in ascending order. By choice of the lengths of the intervals J_n and since W_0 was chosen as the union of these intervals, one has $z_{h(n)}^0 = \max J_n$ for all n. Furthermore, at most $|I_0| + \cdots + |I_n| \leq |J_n|$ times at some stage s a number smaller than $z_{h(n)}^s$ is enumerated into W_{s+1}. Hence for all such stages, one has

$$z_{h(n)-1}^s = z_{h(n)}^s - 1, \text{ hence } z_{h(n)}^{s+1} = z_{h(n)}^s - 1 \text{ and } O(z_{h(n)}^{s+1}) = 1 - O(z_{h(n)}^s) .$$

In particular, during each enumerating stage s the value of the previous approximation to the value $B(h(k_s))$ is flipped from $O(z_{h(k_s)}^s)$ to $O(z_{h(k_s)}^{s+1})$, and after step 2 of each stage s, one has $A_s(k_s) = B_{s+1}(h(k_s))$. By induction on stages one can then show as an invariant of the construction that during each stage s at the end of step 2 it holds that

$$A_s(j) = B_{s+1}(h(j)) = O(z_{h(j)}^{s+1}) \quad \text{for all } s \text{ and all } j \leq k_s .$$

This concludes the verification of the construction because k_s goes to infinity by step 3 and because the sets A_s, B_s and W_s converge pointwise to A, B and W, respectively. □

The two preceding theorems give rise to the following corollary.

Corollary 4. *For any set A the following assertions are equivalent:*

1. *A is ω-r.e.;*
2. *A can be selected from O in two steps;*
3. *A can be selected from O in finitely many steps.*

As a further consequence of Theorem 3, there is a set B that is selected from a recursive set but has a logarithmic lower bound on the plain Kolmogorov complexity $C(\sigma)$ of its initial segments $\sigma = B(0)B(1)\dots B(n)$, hence, in particular, the set B is complex [3]. Here the plain Kolmogorov complexity $C(\sigma)$ of a string σ is the length of the shortest program p such that $U(p) = \sigma$ for some fixed underlying universal machine U, see the textbook of Li and Vitányi [6] for more details. Note that the bound in Corollary 5 is optimal up to a constant factor by the proof of Theorem 11 below, which yields as a special case that every set selected from a recursive set has infinitely many initial segments of at most logarithmic complexity.

Corollary 5. *There is a set B selected from O with $\forall^\infty n\,[C(B(0)B(1)\dots B(n)) \geq 0.5 \cdot \log(n)]$.*

Proof. Section 3 provides a closer look at Chaitin's Ω, which is the standard example of a Martin-Löf random left-r.e. set. From these properties it is immediate that for almost all n it holds that $C(\Omega(0)\Omega(1)\dots\Omega(n)) > n - 3\log n$ and that Ω is ω-r.e. with bounding function $f(n) = 2^{n+1}-1$. Applying the construction in the proof of Theorem 3 with A equal to Ω, one has $1+f(0)+f(1)+\dots+f(n) \leq 2^{n+2}$, hence $h(n) \leq 3^n$ for almost all n. So one can retrieve $\Omega(0)\Omega(1)\dots\Omega(n)$ from $B(0)B(1)\dots B(3^n)$. The corollary now follows by some elementary rearrangements. □

Proposition 7 determines the rank of certain sets with rank 0, 1 or 2. The corresponding arguments use again Theorems 3 and 11, together with the following absorption principle for selections by recursive sets.

Proposition 6. *Let A be selected from E by the r.e. set W and let E be selected from B by the recursive set V. Then A is selected from B.*

Proof. Let $v_0, v_1, \dots$ and $w_0, w_1, \dots$ be the members of V and W, respectively, in ascending order. Note that $n \in A$ iff $w_n \in E$ iff $v_{w_n} \in B$, hence A is selected from B by the r.e. set $\{v_{w_0}, v_{w_1}, \dots\}$. □

Proposition 7. *Exactly the sets $\emptyset$ and $\mathbb{N}$ have selection rank 0. Every finite and every cofinite set that differs from $\emptyset$ and $\mathbb{N}$ has selection rank 1. Every recursive set that is infinite and coinfinite has selection rank 2.*

Proof. The only set that can be selected from the empty set is the empty set itself, hence the empty set has rank 0; the same argument works for $\mathbb{N}$. Next consider a finite set B. In case B differs from $\emptyset$ and $\mathbb{N}$, some set different from B can be selected from A, thus the selection rank of B is at least 1. However, in case a set A is selected from B in several steps, then A and all the intermediate sets are finite and all selecting sets can be taken to be recursive. Then B can be selected from A in a single step according to Proposition 6, hence the selection rank of A is at most 1. Again, an almost identical argument works for the symmetric case of a coinfinite set.

The selection rank of O is at most 2 by Corollary 2, and is at least 2 because by Theorems 3 and 11 the ω-r.e set Ω can be selected from O in two steps but not in one step. Given any infinite and coinfinite recursive set B, the set B can be selected from O by a recursive set and vice versa. By the absorption principle in Proposition 6, from B and O exactly the same sets can be selected in exactly the same number of steps, hence B and O share the same selection rank. □

Recall the following concepts from recursion theory [10, 11, 15]:

- A is truth-table reducible to B iff there are recursive functions f, g where f maps pairs of numbers and strings to bits and g computes the bound of the reduction which is given as $A(x) = f(x, B(0), B(1), \ldots, B(g(x)))$ for every x.
- A is weakly truth-table reducible to B iff the above definition holds with f being only partial recursive while g has to be recursive; furthermore, $f(x, B(0)B(1) \ldots B(g(x)))$ is always defined and provides the right answer $A(x)$; however, if one would replace B by another oracle, the corresponding values of f could be undefined.
- A is one-one reducible to B iff there is a one-one recursive function f with $A(x) = B(f(x))$ for all x.
- A set B is a truth-table cylinder iff there are three recursive functions pad, and, neg such that $\mathrm{pad}(x) > x$ for all x, $B(\mathrm{pad}(x)) = B(x)$, $B(\mathrm{neg}(x)) = 1 - B(x)$ and $B(\mathrm{and}(x, y)) = B(x) \cdot B(y)$. A truth-table cylinder B has the property that whenever a set A is truth-table reducible to B then the set A is also one-one reducible to B via a strictly increasing function.

Proposition 8. *The truth-table cylinder* K^{tt} *of the halting problem has selection rank* 1 *and every* ω*-r.e. set can be selected from it by a recursive set* W*.*

Proof. Let B be the truth-table cylinder of K. It follows from the definition of a truth-table cylinder that every set A which is truth-table reducible to B is already one-one reducible to B; due to the padding requirement one can also achieve that the reduction f is strictly monotonically increasing. Now the recursive set $\{f(0), f(1), f(2), \ldots\}$ selects A from B in the same way as the reduction does.

Furthermore, the truth-table cylinder of the halting problem K is ω-r.e.: For every element x one can compute a formula how $B(x)$ depends on $K(y)$ for some selected values y and then one can estimate that the value of the truth-table changes at most as often as one of these y gets enumerated into K (plus one mind change if the initial value is 1). Hence an upper bound on the number of mind changes of $B(x)$ can be computed from x and B is ω-r.e.; by Theorem 2 it then also follows that every A selected from B is ω-r.e., even if A is selected in several steps from B. Thus, only ω-r.e. sets can be selected from B and every ω-r.e. set can be selected from B in one step, so B has rank 1. □

3 Selection and Ω

Chaitin's Ω is a standard example for a Martin-Löf random set which is in addition also an ω-r.e. set [1]. It will turn out that Ω has various special properties and some but not all of them are shared by Martin-Löf random sets in general. The following gives an overview about Martin-Löf randomness.

Using a characterisation of Schnorr [14], one can say that a set A is Martin-Löf random [7] iff no r.e. martingale M succeeds on A. In this context, a martingale is a function from binary strings to nonnegative real numbers such that $M(\sigma) = (M(\sigma 0)+M(\sigma 1))/2$. M succeeds on a set A iff the set $\{M(\sigma) : \sigma \preceq A\}$ has the supremum ∞, where $\sigma \preceq A$ means that $\sigma(x) = A(x)$ for all x in the domain of σ; similarly one can compare strings with respect to $\preceq$. Furthermore, M is called r.e. iff $\{(\sigma, q) : \sigma \in \{0,1\}^*, q \in \mathbb{Q}, M(\sigma) > q\}$ is recursively enumerable; M is recursive iff the just defined set is recursive. Without loss of generality, one can take a recursive martingale to be $\mathbb{Q}$-valued and can show that whenever some recursive martingale succeeds on A then some $\mathbb{Q}$-valued recursive martingale succeeds on A where in addition the function $\sigma \mapsto M(\sigma)$ is a recursive mapping which returns on input σ the canonical representation (as a pair of numerator and denominator) of $M(\sigma)$. This also holds relativised to oracles. Furthermore, one can say that a partial-recursive martingale M succeeds on A iff for every $\sigma \preceq A$, $M(\sigma), M(\sigma 0), M(\sigma 1)$ are all defined, for every $\sigma \preceq A$ the relation $M(\sigma) = (M(\sigma 0) + M(\sigma 1))/2$ holds and the supremum of $\{M(\sigma) : \sigma \preceq A\}$ is ∞. It is known that if a recursive martingale succeeds on A, then also a partial-recursive martingale succeeds on A; furthermore, if a partial-recursive martingale succeeds on A then a r.e. martingale succeeds on A. A further characterisation by Zvonkin and Levin [16] is that A is Martin-Löf random iff there is no partial-recursive function G compressing A. Here G compresses A iff G maps strings to strings, the domain of G is prefix free – that is whenever $G(p)$ is defined then $G(pq)$ is undefined for all $p, q \in \{0,1\}^*$ with $q \neq \varepsilon$ – and there are infinitely many n for which there is a p of length at most n with $G(p) = A(0)A(1)\ldots A(n)$. The interested reader is referred to standard textbooks on algorithmic randomness for more information [6, 10].

In this section, the relations between Ω and $\sqsubset$ are investigated. First, Theorem 9 shows every set selected from Ω is Martin-Löf random. Second, the next result shows that one cannot select Ω nontrivially in several steps from itself, that is, there are no sets $E_0, E_1, \ldots, E_n$ such that $E_m \sqsubset E_{m+1}$ via an $W_{e_m} \neq \mathbb{N}$ for all $m < n$ and $E_0 = E_n = \Omega$.

Note that for this section, for an r.e. set W with recursive enumeration $W_0, W_1, \ldots$ (that is, the W_s are uniformly recursive, $W = \bigcup_s W_s$ and $W_0 \subseteq W_1 \subseteq \ldots$), one defines the convergence module $c_W(x)$ is the first stage $s \geq x$ such that $W_s(y) = W(y)$ for all $y \leq x$. Furthermore, one fixes an approximation $\Omega_0, \Omega_1, \ldots$ from the left for Ω, that is, this approximation satisfies the following three conditions:

- the Ω_s are uniformly recursive;
- for all x and almost all s, $\Omega(x) = \Omega_s(x)$;
- whenever $\Omega_{s+1} \neq \Omega_s$, then the least element x in the symmetric difference satisfies $x \in \Omega_{s+1} - \Omega_s$.

Now one defines the convergence module of Ω at x as $c_\Omega(x) = \min\{s \geq x : \forall y \leq x\,[\Omega_s(y) = \Omega(y)]\}$. Note that c_Ω, due to Ω being Martin-Löf random, grows much faster than c_W for any given r.e. set W; in particular there is a constant c such that, for all $x > 0$, $c_\Omega(x-1)+c \geq c_W(x)$. This is used in several of the proofs below, in particular as martingales working on Ω and currently having the task to predict $\Omega(x)$, can from the already known values $\Omega(0)\ldots\Omega(x-1)$ figure out which $y \leq x$ are in finitely many fixed r.e. sets and therefore reconstruct the nature of reductions up to x.

Theorem 9 answers an open question by Kjos-Hanssen, Stephan and Teutsch [4, Question 6.1] on whether a set selected from Ω by an r.e. set is Martin-Löf random; the corresponding question with respect to selections by co-r.e. sets also mentioned there is still open.

Theorem 9. *If A is selected from Ω in finitely many steps then A is Martin-Löf random.*

Proof. Assume that there are a number n and sets $E_0, E_1, \ldots, E_n$, $A = E_0$, $\Omega = E_n$ and for all $m < n$ there is an increasing function F_m with $x \in E_m \Leftrightarrow F(x) \in E_{m+1}$ and the range of F_m being an r.e. set. Note that one knows $F_m(0), F_m(1), \ldots, F_m(y)$ at time s iff all elements of the range of F_m below $F_m(y)$ are enumerated within s time steps (with respect to some given recursive enumeration of the range of F_m). As the convergence module of Ω dominates the convergence module of every r.e. set there is a constant c such that one can, for every $x > 0$ and every $m < n$, compute $F_m(y)$ for all y with $F_m(y) \leq x$ within time $c_\Omega(x-1)+c$.

Now assume by way of contradiction that A is not Martin-Löf random. Miller [9] showed that there is an oracle B which is low for Ω and PA-complete; that is, B satisfies that Ω is Martin-Löf random relative to B and that every partial-recursive $\{0,1\}$-valued function has a total B-recursive extension. It is known that every set which is not Martin-Löf random is not recursively random relative to such an oracle B; hence there is a B-recursive martingale M which succeeds on A.

Now it is shown that M can be transformed into a partial B-recursive martingale N succeeding on Ω in contradiction to the choice of B; this N will be defined inductively and the $N(\sigma a)$ will be defined for all $\sigma \preceq \Omega$ and all $a \in \{0,1\}$. This is done by inductively defining sequences $\Phi(\sigma)$ from σ for some partial-recursive function Φ and then letting $N(\sigma) = M(\Phi(\sigma))$. As a starting point, let $\Phi(\varepsilon) = \varepsilon$ and hence $N(\varepsilon) = M(\varepsilon)$. Inductively, $\Phi(\sigma a)$ is defined from $\Phi(\sigma)$ and hence $N(\sigma a)$ from $N(\sigma)$.

Now for any given σ where $\Phi(\sigma)$ and $N(\sigma)$ are defined, one does for $a = 0, 1$ the following: Let $s = t + c$ for the first time $t \geq |\sigma|$ with $\sigma \preceq \Omega_t$; if this time t does not exist then $N(\sigma 0)$ and $N(\sigma 1)$ are undefined. Having s, one computes approximations $F_{m,0}, F_{m,1}, \ldots$ to F_m where $F_{m,s}(y)$ is the y-th element of the set of strings enumerated into the range of F_m within s steps with respect to some recursive enumeration. Let

$$\widetilde{F}_s(y) = F_{n-1,s}(F_{n-2,s}(\ldots(F_{1,s}(F_{0,s}(y)))\ldots))$$

and $\widetilde{F}(y)$ be the limit of $\widetilde{F}_s(y)$. Note that when $\sigma \preceq \Omega$ then $\widetilde{F}_s(y) = \widetilde{F}(y)$ for all y with $\widetilde{F}(y) \leq |\sigma| + 1$ because of the above domination properties; note that the t there would be $c_\Omega(|\sigma|)$. Furthermore, for all y, either $\widetilde{F}_s(y)$ is undefined or $\widetilde{F}_s(y) \geq \widetilde{F}(y)$.

If there is a y such that $\widetilde{F}_s(y) = |\sigma|$ then let $\Phi(\sigma a) = \Phi(\sigma)a$ else let $\Phi(\sigma a) = \Phi(\sigma)$. Furthermore, $N(\sigma a) = M(\Phi(\sigma a))$.

Now one analyses the behaviour of N on Ω. Note that whenever $\sigma a \preceq \Omega$ and $\widetilde{F}_s(y) \in \mathrm{dom}(\sigma a)$

then $\widetilde{F}_s(y) = \widetilde{F}(y)$ where the s is as above. As a consequence, one has for the maximal y with $\widetilde{F}_s(y) \in \mathrm{dom}(\sigma a)$ that $\Phi(\sigma a) = \Omega(\widetilde{F}(0))\Omega(\widetilde{F}(1))\ldots\Omega(\widetilde{F}(y))$ and hence $\Phi(\sigma a) \preceq A$. It follows that N works on Ω like a delayed version of M on A; in particular as M takes on A arbitrarily large values, so does N on Ω. This would mean that N succeeds on Ω in contradiction to the assumption that Ω is Martin-Löf random relative to the oracle B. Thus, against the assumption, the set A has to be Martin-Löf random. □

The following theorem shows that although one can only select Martin-Löf random sets from $\mathit{\Omega}$ in any number of steps, none of these sets is equal to $\mathit{\Omega}$, except for the case that one selects with $\mathbb{N}$.

Theorem 10. *One cannot select Ω nontrivially in several steps from itself, that is, there are no $n > 0$ and no sets $E_0, E_1, \ldots, E_n$ such that $E_m \sqsubset E_{m+1}$ via an $W_{e_m} \neq \mathbb{N}$ for all $m < n$ and $E_0 = E_n = \Omega$.*

Proof. Assume by way of contradiction that the number $n > 0$ and the sets $E_0, E_1, \ldots, E_n$ and r.e. sets $W_{e_0}, W_{e_1}, \ldots, W_{e_{n-1}}$ exist in contradiction to the statement of the theorem. Let $\Omega_0, \Omega_1, \ldots$ be a recursive approximation to Ω from the left and let $c_\Omega(x)$ be the first stage s such that $\Omega_s(y) = \Omega(y)$ for all $y \leq x$. Note that the convergence module c_Ω dominates the convergence module of any r.e. set. Hence, there is a constant c such that, for all x and all $m < n$, $x \in W_{e_m}$ iff $x \in W_{e_m, c_\Omega(x-1)+c}$. This permits to check, using $c_\Omega(x-1)$, whether the iterative nontrivial selection-process maps $y < x$ to x so that $\Omega(y) = \Omega(x)$ could be established by looking at the sets $W_{e_0}, W_{e_1}, \ldots, W_{e_{n-1}}$ only, that is, there are numbers $k_0, k_1, \ldots, k_{n-1} \leq x$ such that x is the k_{n-1}-th element of $W_{e_{n-1}}$, k_{n-1} is the k_{n-2}-th element of $W_{e_{n-2}}$, …, k_1 is the k_0-th element of W_{e_0} and $k_0 = y$. Note that for infinitely many x there is such an $y(x) < x$ found in this way and that $\Omega(y(x)) = \Omega(x)$ whenever such an $y(x)$ exists.

Now one can consider the following partial-recursive martingale M: On input σ with $x = |\sigma|$, if s is the first stage where $\sigma \preceq \Omega_s$ and $y_{s+c}(x)$ is the approximation of $y(x)$ using $W_{e_m, s+c}$ in place of W_{e_m} for all $m < n$. If s does not exist then $M(\sigma 0), M(\sigma 1)$ are undefined else do the following: If $y_s(x)$ exists then $M(\sigma \cdot \sigma(y_s(x))) = 2 \cdot M(\sigma)$ and $M(\sigma \cdot (1 - \sigma(y(x)))) = 0$ else $M(\sigma 0) = M(\sigma 1) = M(\sigma)$.

Note that for $\sigma \preceq \Omega$, the s always exists and $y_s(x)$ does not exist if $y(x)$ does not exist and $y_s(x) = y(x)$ if $y(x)$ exists. Hence, M works on Ω correctly and its capital goes to infinity in contradiction to the fact that no partial-recursive martingale succeeds on Ω. Hence there cannot be any nontrivial selection of Ω itself from Ω in finitely many steps. □

Furthermore, one can also show the following: interchange even and odd positions by letting $\widetilde{\Omega}(2n) = \Omega(2n+1)$ and $\widetilde{\Omega}(2n+1) = \Omega(2n)$; the set $\widetilde{\Omega}$ cannot be selected from Ω in any number of steps.

4 Selection and Martin-Löf random sets in general

After having investigated relations between selection and the special Martin-Löf random set Ω, the focus is now on relations between selection and Martin-Löf random sets in general. First, Theorems 11 and 12 exhibit classes of sets from which no Martin-Löf random set can be selected in one step. Furthermore, Theorems 15 and 14 assert that there is a Martin-Löf random set from which one can select all ω-r.e. sets in up to two steps, whereas no recursive set can be selected from any Martin-Löf random set in one step.

Theorem 11. *Assume that B is Turing reducible to a Turing-incomplete r.e. set. Then no set selected from B is Martin-Löf random.*

Proof. Let B be stated as in the theorem. Recall that a sufficient criterion for a set A to be not Martin-Löf random is that there are infinitely many n such that the plain Kolmogorov complexity of $A(0)A(1)\ldots A(n)$ is bounded proportionally to $\log(n)$. Indeed, in the following it is shown that there are a constant c and infinitely many n such that $C(A(0)A(1)\ldots A(n)) \leq 2\cdot\log(n)+c$.

Consider any $A \sqsubset B$ and let W be the r.e. set with $A(n) = B(w_n)$ for the n-th element w_n of W in ascending order. Let $b_0, b_1, b_2, \ldots$ be a recursive one-one enumeration of W and let $e_0 = 0$ and e_{n+1} be the first number $d > e_n$ such that $b_{e_n} < b_d$. Note that the mapping $m \mapsto b_{e_m}$ is recursive. Now given any m, let n be the number with $b_{e_m} = w_n$, note that $m \leq n$. Knowing m and n, one can compute $w_0, w_1, \ldots, w_n$.

There is a recursive approximation $B_0, B_1, \ldots$ to B such that the convergence module g of this approximation does not permit to compute the diagonal halting problem K. In particular there are infinitely many $m \in K$ such that m is enumerated into K at a stage s larger than $g(b_{e_m})$ and all w_k with $k \leq n$ satisfy $B_s(w_k) = B(w_k)$. Hence, for these m and the corresponding n, $A(0)A(1)\ldots A(n)$ can be described by m and n using the time s when m is enumerated into K and the members $w_0, w_1, \ldots, w_n$ of W obtained from m, n and conjecturing that $A(k) = B_s(w_k)$ for $k = 0, 1, \ldots, n$. For the right parameters, the s exists and the corresponding data can be computed and the resulting string is correct. As one can describe m and n by two numbers of $\log(n)$ binary digits (the number of digits must be the same for permitting to separate out the digits from m from those for n), $C(A(0)A(1)\ldots A(n)) \leq 2\cdot\log(n)+c$ for some constant c and infinitely many n. It follows that A is not Martin-Löf random. $\Box$

Theorem 12. *Every truth-table degree contains a set B such that no set selected from B is Martin-Löf random.*

Proof. This proof is mainly based on the fact that every truth-table degree contains a retraceable set; here a set B is retraceable iff there is a partial-recursive function ψ which returns for every $x \in B$ a canonical index of the set $\{y \leq x : y \in B\}$; on $x \notin B$, ψ can either be undefined or return any information, either wrong or right. For example, if E is a given set then the set $B = \{x_0, x_1, \ldots\}$ with $x_0 = 1$ and $x_{n+1} = 2x_n + E(n)$ for all n is a retraceable set of the same truth-table degree as E. So fix such B and ψ with B being inside the given truth-table degree. The proof follows now in general the proof of Theorem 11 with the adjustment that it is shown that for each $A \sqsubset B$ there are a constant c and infinitely many n such that the plain Kolmogorov complexity of $A(0)A(1)\ldots A(n)$ is bounded by $3\cdot\log(n)+c$, which then gives that A is not Martin-Löf random.

Consider any $A \sqsubset B$ and let W be the r.e. set with $A(n) = B(w_n)$ for the n-th element w_n of

W in ascending order. Without loss of generality, $0 \in A$. Let $b_0, b_1, b_2, \ldots$ be a recursive one-one enumeration of W and let $e_0 = 0$ and e_{n+1} be the first number $d > e_n$ such that $b_{e_n} < b_d$. Note that the mapping $m \mapsto b_{e_m}$ is recursive. Now given any m, let n be the number with $b_{e_m} = w_n$, note that $m \leq n$. Knowing m and n, one can compute $w_0, w_1, \ldots, w_n$. Furthermore, let k be such that w_k is the maximal of the $w_0, w_1, w_2, \ldots, w_n$ with $w_k \in B$. Note that $k \leq n$ and k exists as $0 \in A \wedge w_0 \in B$.

Hence, for each n and the corresponding $m, k \leq n$, one can compute $w_0, w_1, \ldots, w_n$ from m, n and use $\psi(w_k)$ to find out which of these numbers are in B. Hence $A(0)A(1)\ldots A(n)$ can be computed from n, m, k. One can code m, n, k as 3 binary numbers of $\log(n)$ digits each and gets therefore that there are a constant c and infinitely many n such that $C(A(0)A(1)\ldots A(n)) \leq 3 \cdot \log(n) + c$. Hence the set A is not Martin-Löf random. $\square$

Proposition 17 below shows that if B is hyperimmune and $A \sqsubset B$ then A is hyperimmune; as there are uncountably many hyperimmune sets and these are never Martin-Löf random, one has the following result similar to the previous one.

Fact 13. *There are uncountably many sets B such that no set that is selected from B in one or several steps is Martin-Löf random.*

Theorem 14 below shows that one cannot select a recursive set from a Martin-Löf random set in one step. Note that the proof of Theorem 14 indeed shows that it is not possible to select from a Martin-Löf random set any set that obeys certain upper bounds on the complexity of its initial segments.

Theorem 14. *It is not possible to select a recursive set from a Martin-Löf random set.*

Proof. Assume that $A \sqsubset B$ via an r.e. set W and A is recursive and F is the function which lists W in ascending order (F is not recursive). So $A(x) = B(F(x))$ for all x and $W = \{F(0), F(1), \ldots\}$. Let $u_0, u_1, \ldots$ be an ascending recursive enumeration of a recursive subset of W which is selected such that W has at least 2^{n+2} elements below a given u_n. Now one shows that there is a partial-recursive function G with prefix-free domain which compresses B, that is, for which there are infinitely many $p \in \mathrm{dom}(G)$ with $G(p)$ being a prefix of B which is longer than $|p|$; this would then be an alternative way to prove that B is not Martin-Löf random.

On input $p = 0^n 1 0^m 1 b_0 b_1 \ldots b_m c_0 c_1 \ldots c_k$, $G(p)$ first checks whether $k + 1 = u_n - d$ where d is the binary value of $b_0 b_1 \ldots b_m$. In the case that this is true, $G(p)$ enumerates the W until d many elements at places $\widetilde{F}(0), \widetilde{F}(1), \ldots, \widetilde{F}(d)$ have been enumerated into W with $\widetilde{F}(0) < \widetilde{F}(1) < \ldots < \widetilde{F}(d) = u_n$. If this is eventually achieved and if $d \geq 2^{n+2}$, then G outputs a string $\sigma \in \{0,1\}^{u_n+1}$ which is obtained by letting $\sigma(\widetilde{F}(d')) = A(d')$ for all $d' \leq d$ and by filling the remaining missing $k+1$ values in σ below the position u_n according to the string $c_0 c_1 \ldots c_k$. This results in a string of length u_n which is computed from a p of length $n + 2m + k + 4$; by taking m as small as possible, one has that $m \leq \log(d) + 1$ and $n \leq \log(d)$, thus one has a length bounded by $u_n + 3\log(d) - d$ which is, for all sufficiently large n and the corresponding d smaller than u_n.

One has now to show that one can always choose d, m, $b_0 b_1 \ldots b_m$ and $c_0 c_1 \ldots c_k$ such that the corresponding output $G(p)$ is $B(0)B(1)\ldots B(u_n)$. To see this, let d be the number of strings in W up to u_n (which is larger than n) and $m = \log(d)$ and $b_0 b_1 \ldots b_m$ be the binary representation of d. Furthermore, let $k = u_n - d - 1$. One gets that $\widetilde{F}(d') = F(d')$ for all $d' \leq d$. Now one chooses $c_0 c_1 \ldots c_k$ such that the missing positions in σ which are not covered by $F(0), F(1), \ldots, F(d)$ are covered with the corresponding bits of B. Hence one has that for the so selected p that $G(p)$ equals $B(0)B(1)\ldots B(u_n)$. It is furthermore easy to verify that the domain of G is prefix-free. $\square$

The following result stands in contrast to Theorem 9, which says that one cannot select any nonrandom set from Ω in arbitrarily many steps. Note that the resulting set B is like Ω also an ω-r.e. Martin-Löf random set.

Theorem 15. *There is an ω-r.e. Martin-Löf random set B such that some set selected from B is not Martin-Löf random and every ω-r.e. set can be selected from B in two steps.*

Proof. First one gives a construction of a Martin-Löf random set B from a given ω-r.e. set A such that A can be selected from B in two steps. Second, one chooses A such that every ω-r.e. set can be selected from B in two steps which fixes the set B. Third, one will show that therefore some nonrecursive set can be selected in one step from B.

First, let T be an infinite binary recursive tree such that every infinite branch of the tree is Martin-Löf random; it is well-known that such trees exist [6, 10] and they come out naturally from an universal Martin-Löf test [7]. Now one can for every level n and every m find, recursively in n and m, a level $f(n,m)$ such that every infinite C either has between n and $f(n,m)$ the subwords 0^m and 1^m or C is cut off from the recursive tree below level $f(n,m)$. This is done by exhaustive search and the termination of the procedure is due to compactness. Let $T_{n,m,r}$ be the class of all subsets of $\mathbb{N}$ which are either cut off from T below r or which have both subwords 0^m and 1^m between the lengths n and r. Then every infinite set C is in some $T_{n,m,r}$, as it is either eventually cut off from T or has infinitely many copies of 0^m and 1^m in its characteristic function; as the Cantor space is compact and the $T_{n,m,r}$ form an ascending union of open classes covering the Cantor space, there is an r such that $T_{n,m,r}$ contains the full Cantor space and one can set $f(n,m) = r$ for this r. Note that $f(n,m) \geq n + 2m$, as otherwise there would not be the space for the subwords 0^m and 1^m between the positions of n and $f(n,m)$ on any infinite branch of T.

Now let A be an ω-r.e. set such that $A(x)$ is approximated by a recursive sequence $(A_t(x))$ using strictly less than $g(x)$ mind changes and, without loss of generality, $g(x+1) \geq 2 \cdot g(x) \geq 1$ for all x. Furthermore, one will construct the set B such that B is the limit of finite strings σ_n of length $h(n)$ such that each σ_n makes at most $k(n) = k(n-1) + g(n) + 2^{h(n)}$ many mind changes with respect to its value; to accomodate the mind changes and nevertheless being able to satisfy all constraints, the recursive function h will be taken such that

$$h(n+1) = \tilde{f}(h(n), h(n)+1, k(n)+g(n)+1)$$

where $\tilde{f}(a,b,0) = a$ and $\tilde{f}(a,b,c+1) = f(\tilde{f}(a,b,c),b)$ for all a,b,c; in other words, $\tilde{f}(a,b,c)$ says that one starts with a and iterates the replacement $a = f(a,b)$ then c times in order to make sure that there are c copies of 0^b and 1^b on the infinite strings on T between a and $\tilde{f}(a,b,c)$. More precisely, one has a copy of 0^b and 1^b between each values $\tilde{f}(a,b,c)$ and $\tilde{f}(a,b,c+1)$.

Now the inductive construction of the σ_n is that $\sigma_0 = \varepsilon$ and given σ_n, σ_{n+1} is chosen at stage t together with values s_t and r_t and an r.e. set W such that the following constraints are satisfied:

- r_t is the number of mind changes done so far by either the approximation of $A(n)$ or the approximation of σ_n;
- s_t is the largest position of an $A_t(n))^{h(n)+1}$ on some $\tau \in T \cap \{0,1\}^{h(n+1)+t}$ with $\sigma_n \preceq \tau$ and $s_t \leq \tilde{f}(h(n), h(n)+1, k(n)+g(n)+1-r_t) - h(n) - 1$;
- σ_{n+1} is the prefix of length $h(n)+1$ of the lexicographically least τ satisfying the previous item at the choice of s_t;
- $W_t \cap \{h(n), h(n)+1, h(n)+2, \ldots, h(n+1)-1\} = \{s_t, s_t+1, \ldots, h(n+1)-1\}$.

Note that each mind change of σ_{n+1} is either caused by a mind change of σ_n or a mind change of $A(n)$ or of all branches through σ_{n+1} being eventually cut off on T. Thus $k(n+1)$ can be chosen as $k(n)+g(n+1)+2^{h(n+1)}$ as indicated above; note that the $g(n+1)$ here counts also, as an upper bound, all the mind changes of A below n, as $g(n+1) \geq 2 \cdot g(n)$ for all n. Furthermore, the values r_t go, for each σ_{n+1}, only upwards as they count mind changes and the values s_t go down; between two changes of r_t, these values are inside the interval I from $\tilde{f}(h(n), h(n)+1, k(n)+g(n)-r_t)$ to $\tilde{f}(h(n), h(n)+1, k(n)+g(n)+1-r_t)$ and they take there always the largest possibility such that there is an extension τ of σ_n having at s_t a run of $h(n)+1$ values $A_t(n)$ without overshooting the upper bound $\tilde{f}(h(n), h(n)+1, k(n)+g(n)+1-r_t)$. As with growing t there are less and less possible choices for σ_{n+1}, the values of s_t go down inside the interval I as long as σ_n has some extension which qualifies; if all these extensions are cut off, then σ_n itself will make a mind change and therefore r_t will also go up and the next lower interval will be used by the next choice for the r_t belonging to the definition of σ_{n+1}. Note that enough provision is given in order to be able to deal with all these mind changes and that also, as s_t goes only downward, the set W grows monotonically with t and is thus recursively enumerable. Although the definition of W is given only for one interval between $h(n)$ inclusively and $h(n+1)$ exclusively, the construction is carried out for all n and therefore the definition of W extends to the full set of natural numbers.

When selecting a set C from B through W, it will be that for each n, there will be $h(n)+1$ numbers in one continuous run be selected from B starting before $h(n)$ which are all equal to $A(n)$. The reason is that there is a run of $(A(n))^{h(n)+1}$ at B starting with the limit of the s_t in the corresponding interval, as W selects in the interval from $h(n)$ to $h(n+1)-1$ from B starting with the position of an interval of length $h(n)+1$ of symbols $A(n)$. It follows that A can be selected from B by doing a second step via selecting with the recursive set $\{h(1), h(2), \ldots\}$.

Furthermore, note that $B(n)$ is represented as a bit of σ_n and that σ_n makes at most $k(n)$ mind changes, thus B is ω-r.e. with the bound $k(n)$.

Second, now one chooses A as the truth-table cylinder of K. By this choice, every further ω-r.e. set can be selected in one step from A via a recursive set. As the second step of the selection of A is via a recursive set, one concatenates the two selections via recursive sets into one, as outlined in Proposition 6. Thus every ω-r.e. set can be selected from B in two steps; note that one of these steps, namely the first one, has indeed to employ a nonrecursive r.e. set, as otherwise one would only be able to select Martin-Löf random sets from the Martin-Löf random set B.

Third, one considers the case that the set of odd numbers O is selected in two steps from B, that is, that there is an A such that $A \sqsubset B$ and $O \sqsubset A$. Now, by Theorem 14, the set A cannot be Martin-Löf random; hence a set A which is not a Martin-Löf random set can be selected from B in one step. □

5 Closure under Selection

The previous sections had shown that ω-r.e. sets are closed under selections via r.e. sets while Martin-Löf random sets are not. Now the corresponding question is investigated for a few other recursion-theoretic notions. As one uses selection via infinite r.e. sets, a selection is always an ascending one-one reduction relative to the oracle K. As the K-recursive, the 2-random and the 2-generic sets are each closed under such reducibilities, one has directly the following result.

Fact 16. *Assume that B is selected from A in one step via an infinite r.e. set. If B is 2-random, so is A; if B is 2-generic, so is A; if B is Turing reducible to K, so is A.*

Furthermore, one can establish that hyperimmune sets are closed under r.e. selection. Here recall that a set B is hyperimmune iff for every recursive function f there are infinitely many n such that $\{n, n+1, \ldots, n+f(n)\} \subseteq B$.

Proposition 17. *Let B be a hyperimmune set. Then every set $A \sqsubset B$ is hyperimmune.*

Proof. Given A, B as above, let W be an infinite r.e. set such that A is selected from B via W. W has an infinite recursive subset V. For every recursive function f, there is a recursive increasing function g such that, between n and $n+g(n)$ there are always at least $f(m)+2$ elements of V for each $m \leq n$. Now there are infinitely many n such that $\{n, n+1, \ldots, n+g(n)\} \subseteq B$. For each such n there is a least m such that the m-th element w_m of W in ascending order is at least n. Now, for all k with $m \leq k \leq m+f(m)$ it holds that $n \leq w_k \leq n+g(n)$ and thus $w_k \in B$, thus it follows that $\{m, m+1, \ldots, m+f(m)\} \subseteq A$. Thus A is hyperimmune. □

While selection via r.e. sets preserves classes of sets defined by notions which use the oracle K and also some notions like hyperimmune, this is no longer true for weaker notions. For example, if one only asks for a hyperimmune Turing degree instead of the set itself being hyperimmune, then one could consider the join of $\mathbb{N}$ with a hyperimmune set and one can easily select $\mathbb{N}$ from this join. Furthermore, the recursive sets and the sets of hyperimmune-free Turing degree and the sets of incomplete Turing degree (that is, a Turing degree not above the halting problem K) are not closed under selection, as one can select from the set O of odd numbers in two steps the set K, see Theorem 3. Furthermore, Martin-Löf random sets are biimmune and one can select from the Martin-Löf random set B in Theorem 15 some of these sets in two steps the set of odd numbers and thus a set which is not biimune can be selected from some set which is biimune; similarly for the notions of immune and coimmune. The next theorem completes the picture of this type of results by showing that the 1-generic sets are not closed under selection via r.e. sets.

Theorem 18. *There is a* 1*-generic set* B *and a set* $A \sqsubset B$ *which is not even weakly* 1*-generic.*

Proof. There are 1-generic sets B which are trivial for prefix-free Kolmogorov complexity. These trivial sets satisfy that one can, for almost all n and for $m = 2^{2^n}$, enumerate up to 2^{2n} possible vectors $\sigma \in \{0,1\}^m$ with one of them coinciding with B on the first m bits; these strings enumerated can be chosen as those with $C(\sigma) \leq 2^{2n}$. Now one shows that one can choose an r.e. set W such that the $A \sqsubset W$ satisfies that it alnost never meets the extension function G which appends to σ the extension $(01)^k$ for $k = 2^{2^{|\sigma|+10}}$.

Now one let W contain all even numbers and some odd numbers enumerated as below. For this, one defines the notion of η selected from τ on interval I via W iff $|\tau| = |I|$ and η is the substring of τ which is taken by using the members of I as indices of τ and then selecting those bits for which the index is in $I \cap W$. The main idea is the following: If A satisfies the extension function G then there must be an interval of the form $I_n = \{m : 2^{2^n} \leq m < 2^{2^{n+1}}\}$ such that W selects a string η in the set $E = \{\varepsilon, 1\} \cdot \{01\}^* \cdot \{\varepsilon, 0\}$ from the restriction τ of B to I. On one hand, W contains all even numbers and $|\eta| \geq 2^{2^{n+1}-1} - 2^{2^n-1}$; on the other hand, $|\eta| \leq |\tau| = 2^{2^{n+1}} - 2^{2^n}$. Now the construction is the following:

> For each interval I_n, one first puts all even numbers into W on I_n. Whenever one finds a string τ with index I_n such that the current W_s maps τ to a member of E and $C(\tau) < 2n$, then one looks for the least odd number w in I_n not yet enumerated into W_s and if w exists and is not the maximum of I_n then one puts w into W_{s+1} else $W_{s+1} = W_s$ (on I_n).

Note that if τ is mapped by W to a member of E and w exists then $\tau(w-1) \neq \tau(w+1)$ and therefore $\tau(w)$ coincides with one of these two values and for all $t \geq s+1$, the η obtained from τ by W_t contains either 00 or 11 as a subword, thus is not in E. As there are only 2^{2n} such strings τ but more than $2^{2^{n-2}}$ odd numbers in $I_n - \max(I_n)$, this construction will, for all sufficiently large n, always be able to find an odd number as required; the finitely many exception n can be ignored, as one only wants to avoid that the extension function is met infinitely often by B, but it might be met finitely often.

Thus, for all but finitely many n, B restricted to the interval I_n will not be mapped to a member of E and thus B is not weakly 1-generic. ☐

6 Conclusion

The present paper focussed on the question when a set A is one-one reducible to B via the principal function of an r.e. set and generalised this notion also to reductions in several steps, as this reducibility is not transitive. The investigations show that there is a rich relation between this type of reducibility and ω-r.e. sets and Martin-Löf random sets. Future work might in particular address the question for which numbers $n \in \{0, 1, 2, \ldots, \infty\}$ there are sets A of selection rank n; for $n = 0, 1, 2$, examples are given within this paper and all of these examples are ω-r.e. sets.

Indeed, any recursive set A has selection rank 0 if and only if A is equal to $\emptyset$ or $\mathbb{N}$, has selection rank 1 if and only if A is finite or cofinite and differs from $\emptyset$ and $\mathbb{N}$, and has selection rank 2 if and only if A is infinite and coinfinite. Some Martin-Löf random set has selection rank 2 and permits to select a set in finitely many steps iff it is an ω-r.e. set.

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