

## 20—Meta-circular Evaluator

CS1101S: Programming Methodology

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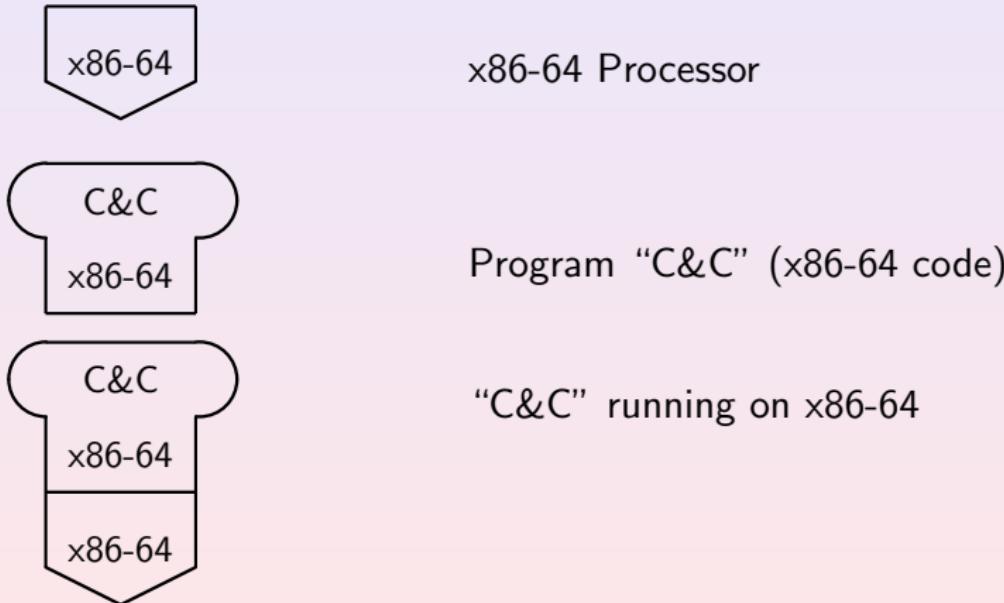
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# T-Diagrams



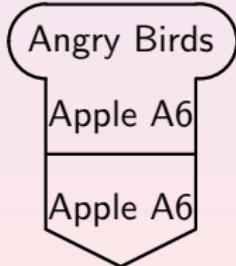
# Running an app on iPhone 5



Apple A6, the processor of iPhone 5



Program “Angry Birds” (A6 build)



“Angry Birds” running on iPhone 5

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# Interpreter

- Interpreter is program that executes another program
- The interpreter's *source language* is the language in which the interpreter is written
- The interpreter's *target language* is the language in which the programs are written which the interpreter can execute

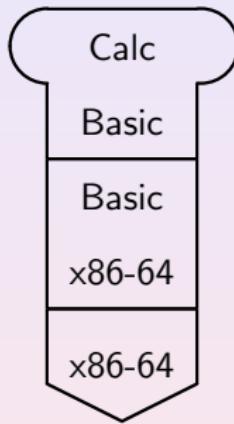
# Interpreters

Basic

x86-64

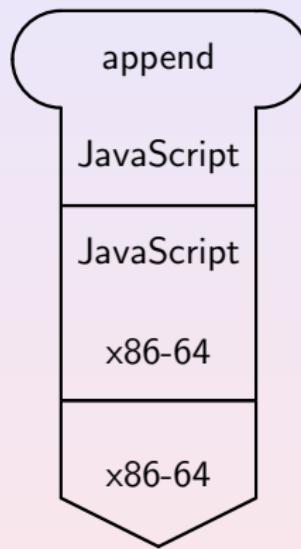
Interpreter for Basic, written in x86-64 machine code

# Interpreting a Program



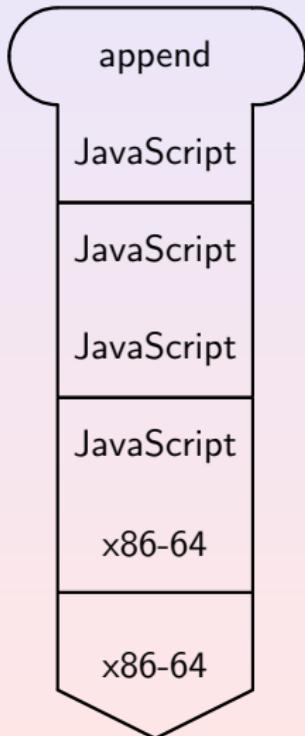
Basic program “Calc”  
running on x86-64 using interpretation

# Running JavaScript in Firefox



JavaScript “append” program running in Firefox

# Running the meta-circular interpreter in Firefox



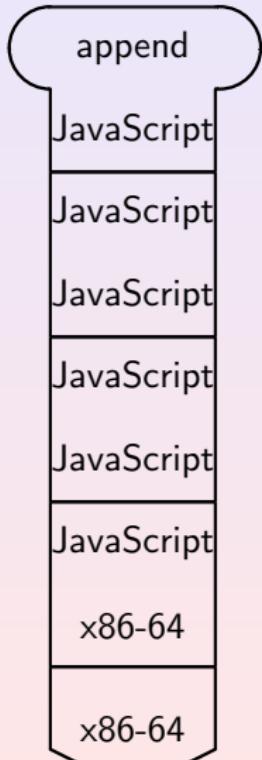
JavaScript “append” program running in Firefox using `parse_and_evaluate ("... append (...);")` in Web Console.

# Gaps between JavaScript and Meta-circular Interpreter

- Assignment to undeclared variables
- Variable declarations in function bodies
- Treatment of `undefined` in Firefox
- Missing features: `new`, `while`, etc

Ambition: Add enough features to be able to interpret the interpreter

# Meta-meta



JavaScript “append” program  
running in Firefox using

```
parse_and_evaluate(  
    "function_evaluate(...);..."+  
    "parse_and_evaluate(" +  
    "... append(...); ''"+  
    " );"  
);
```

in Web Console.

# Assignment to undeclared variables

## Idea

Modify `set_variable_value` such that the variable is added at toplevel (global environment)

# Variable declarations in function bodies

## Idea

Whenever a function is being defined, collect the variable declarations of the body, and save the list of declared variables in the function value

## Modify function application

When functions are applied, extend the environment by a binding of the declared variables to `undefined`.

# Recall Streams

## Central idea of streams

Place data in head; represent the rest of the data (tail) by a function that is called when the data is needed

## Why stop at streams?

We could aim for a general principle of “laziness”: Compute values only when they are needed.

# Lazy Functional Programming

## Central idea

Execute applications of user-defined functions such that argument expressions are not evaluated, but placed in functions

## When will the functions ever be evaluated?

When we apply a primitive operation, such as “+” to them, or when we “force” evaluation to get the overall result of the program.

# Implementation

```
function list_of_values(exp, env) {
    if (no_operands(exp)) return [];
    else return pair(make_thunk(
        first_operand(exp), env),
        list_of_values(
            rest_operands(exp), env)));
}
```

# Implementation

```
function evaluate(stmt, env) {  
    ...  
    else if (is_variable(stmt)) {  
        var val=lookup_variable_value(  
            variable_name(stmt),  
            env);  
        if (is_thunk(val)) {  
            return force_thunk(val);  
        } else {  
            return val;  
        }    }    }}
```

# Background

## Lazy languages exist

A number of programming languages use lazy functional programming as the general evaluation mechanism: Haskell, Miranda

## Lambda-calculus

The semantics of such languages is defined using a strategy to evaluate lambda calculus expressions, called *normal order reduction*