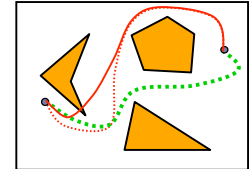


# Path Planning for Point Robots

## Problem

- Input
  - Robot represented as a **point** in the **plane**
  - Obstacles represented as polygons
  - Initial and goal positions
- Output
  - A collision-free path between the initial and goal positions



## Framework

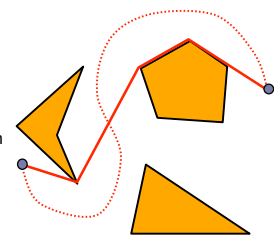
**continuous representation**  
(configuration space formulation)

**discretization**  
(random sampling, processing critical geometric events)

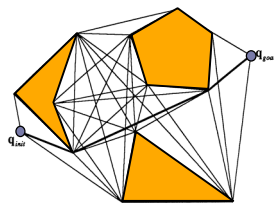
**graph searching**  
(breadth-first, best-first, A\*)

## Visibility graph method

- **Observation:** If there is a collision-free path between two points, then there is a polygonal path that bends only at the obstacle vertices.
- **Why?** Any collision-free path can be transformed into a polygonal path that bends only at the obstacle vertices.
- A **polygonal path** is a piecewise linear curve.



## What is a visibility graph?



- A **visibility graph** is a graph such that
  - Nodes:  $q_{init}$ ,  $q_{goal}$ , or an obstacle vertex.
  - Edges: An edge exists between nodes  $u$  and  $v$  if the line segment between  $u$  and  $v$  is an obstacle edge or it does not intersect the obstacles.

## A simple algorithm for building visibility graphs

**Input:**  $q_{init}$ ,  $q_{goal}$ , polygonal obstacles  
**Output:** visibility graph  $G$

```

1: for every pair of nodes  $u, v$ 
2:   if segment( $u, v$ ) is an obstacle edge then
3:     insert edge( $u, v$ ) into  $G$ ;
4:   else
5:     for every obstacle edge  $e$ 
6:       if segment( $u, v$ ) intersects  $e$ 
7:         go to (1);
8:     insert edge( $u, v$ ) into  $G$ .
```

## Computational efficiency

```

1: for every pair of nodes  $u, v$   $O(n^2)$ 
2:   if segment( $u, v$ ) is an obstacle edge then  $O(n)$ 
3:     insert edge( $u, v$ ) into  $G$ ;
4:   else
5:     for every obstacle edge  $e$   $O(n)$ 
6:       if segment( $u, v$ ) intersects  $e$ 
7:         go to (1);
8:     insert edge( $u, v$ ) into  $G$ .

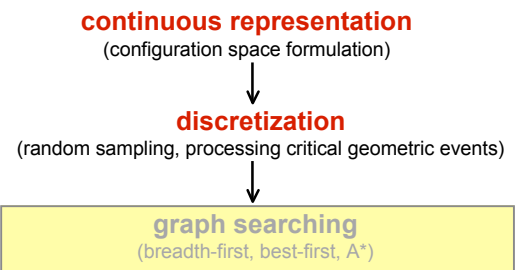
```

- Simple algorithm  $O(n^3)$  time
- More efficient algorithms
  - Rotational sweep  $O(n^2 \log n)$  time
  - Optimal algorithm  $O(n^2)$  time
  - Output sensitive algorithms
- $O(n^2)$  space

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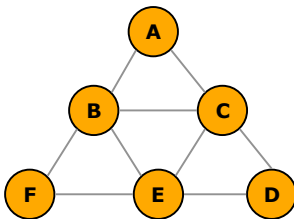
## Framework



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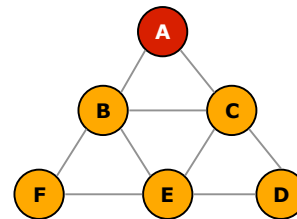
## Breadth-first search



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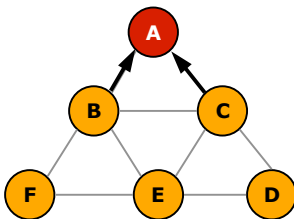
## Breadth-first search



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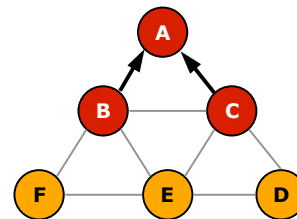
## Breadth-first search



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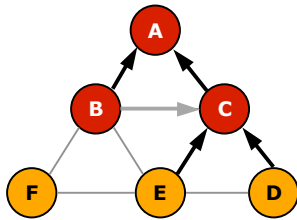
## Breadth-first search



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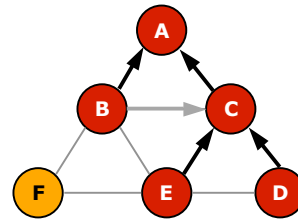
## Breadth-first search



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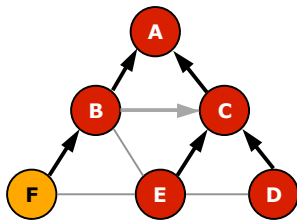
## Breadth-first search



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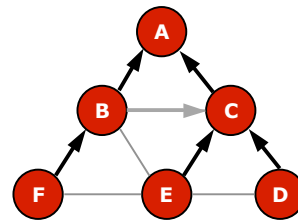
## Breadth-first search



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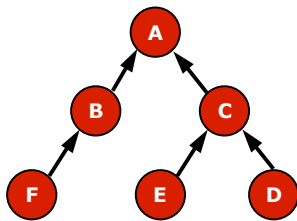
## Breadth-first search



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## Breadth-first search

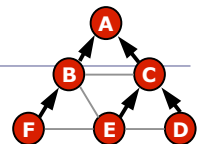


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## Breadth-first search

**Input:**  $q_{init}$ ,  $q_{goal}$ , visibility graph  $G$   
**Output:** a path between  $q_{init}$  and  $q_{goal}$



```

1: Q = new queue;
2: Q.enqueue( $q_{init}$ );
3: mark  $q_{init}$  as visited;
4: while Q is not empty
5:   curr = Q.dequeue();
6:   if curr ==  $q_{goal}$  then
7:     return curr;
8:   for each w adjacent to curr
10:    if w is not visited
11:      w.parent = curr;
12:      Q.enqueue(w)
13:      mark w as visited
  
```

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## Other graph search algorithms

- Depth-first
- Best-first, A\*

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## Framework

continuous representation



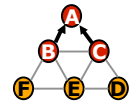
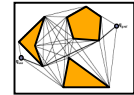
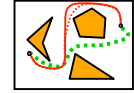
discretization

construct visibility graph



graph searching

breadth-first search



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## Summary

- Discretize the space by constructing visibility graph
- Search the visibility graph with breadth-first search
- How to perform the intersection test?

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## Computational efficiency

- Running time  $O(n^3)$ 
  - Compute the visibility graph
  - Search the graph
  - An optimal  $O(n^2)$  time algorithm exists.
- Space  $O(n^2)$
- Can we do better?

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## Classic path planning approaches

- **Roadmap**  
Represent the connectivity of the free space by a network of 1-D curves
- **Cell decomposition**  
Decompose the free space into simple cells and represent the connectivity of the free space by the adjacency graph of these cells
- **Potential field**  
Define a potential function over the free space that has a global minimum at the goal and follow the steepest descent of the potential function

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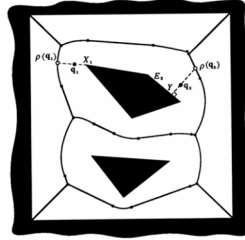
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## Roadmap

- Visibility graph  
Shakey Project, SRI [Nilsson, 1969]

- Voronoi diagram  
Introduced by computational geometry researchers. Generate paths that maximizes clearance.

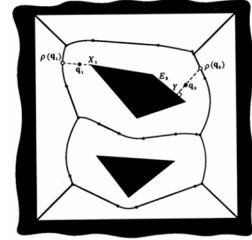


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## Voronoi diagram method

- Space  $O(n)$
- Running time  $O(n \log n)$
- Applicable mostly to 2-D configuration spaces



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## Other roadmap methods

- Silhouette  
First complete general method that applies to spaces of any dimensions and is singly **exponential** in the number of dimensions [Canny, 87]
- Probabilistic roadmaps

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## Classic path planning approaches

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Decompose the free space into **simple** cells and represent the connectivity of the free space by the adjacency graph of these cells
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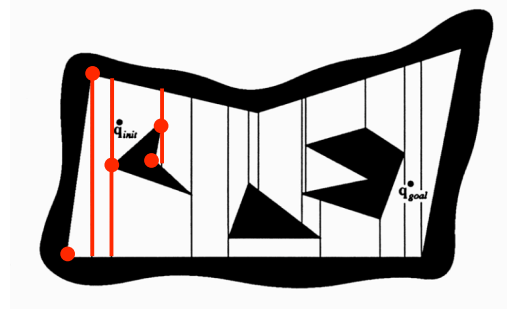
## Cell-decomposition methods

- Exact cell decomposition  
The free space  $F$  is represented by a collection of non-overlapping simple cells whose union **is exactly**  $F$ .
  - Examples of cells: trapezoids, triangles

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## Trapezoidal decomposition



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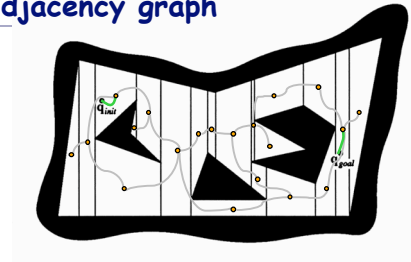
## Computational efficiency

- Running time  $O(n \log n)$  by planar sweep
- Space  $O(n)$
- Mostly for 2-D configuration spaces

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## Adjacency graph

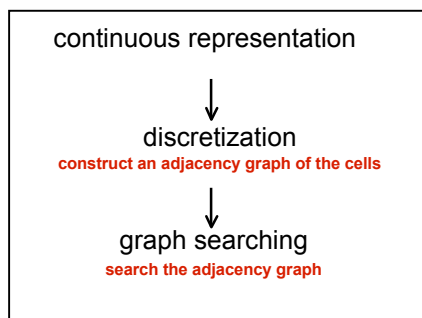


- Nodes:** cells
- Edges:** There is an edge between every pair of nodes whose corresponding cells are adjacent.

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## Framework



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## A brief summary

- Discretize the space by constructing an adjacency graph of the cells
- Search the adjacency graph

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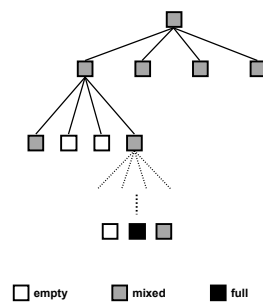
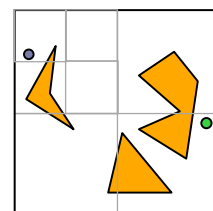
## Cell-decomposition methods

- Exact cell decomposition
- Approximate cell decomposition**  
The free space  $F$  is represented by a collection of non-overlapping cells whose union is **contained** in  $F$ .
  - Cells usually have simple, regular shapes, e.g., rectangles, squares.
  - Facilitate hierarchical space decomposition

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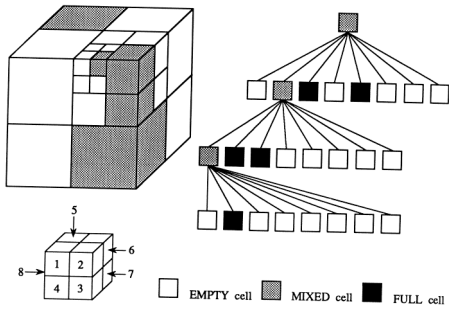
## Quadtree decomposition



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## Octree decomposition



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## Sketch of the algorithm

1. Decompose the free space  $F$  into cells.
2. Search for a sequence of **mixed** or **free** cells that connect the initial and goal positions.
3. Further decompose the mixed.
4. Repeat (2) and (3) until a sequence of **free** cells is found.

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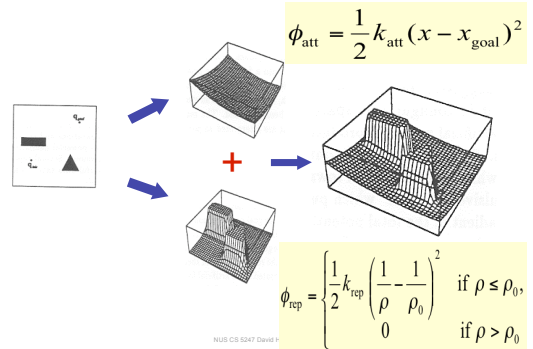
## Classic path planning approaches

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## Algorithm in pictures



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## Attractive & repulsive fields

$$F_{\text{att}} = -\nabla \phi_{\text{att}} = -k_{\text{att}} (x - x_{\text{goal}})$$

$$F_{\text{rep}} = -\nabla \phi_{\text{rep}} = \begin{cases} k_{\text{rep}} \left( \frac{1}{\rho} - \frac{1}{\rho_0} \right) \frac{1}{\rho^2} \frac{\partial \rho}{\partial x} & \text{if } \rho \leq \rho_0 \\ 0 & \text{if } \rho > \rho_0 \end{cases}$$

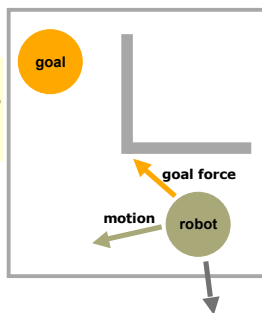
$k_{\text{att}}, k_{\text{rep}}$ : positive scaling factors

$x$ : position of the robot

$\rho$ : distance to the obstacle

$\rho_0$ : distance of influence

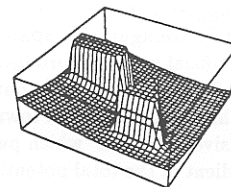
[Khatib, 1986]



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## Local minima



- What can we do?
  - Escape from local minima by taking random walks
  - Build an ideal potential field – navigation function – that does not have local minima

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## Sketch of the algorithm

- Place a regular grid  $G$  over the configuration space
- Compute the potential field over  $G$
- Search  $G$  using a best-first algorithm with potential field as the heuristic function

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## Potential field

- Initially proposed for real-time collision avoidance [Khatib, 1986]. Hundreds of papers published on this topic.
- A potential field is a scalar function over the free space.
- To navigate, the robot applies a force proportional to the negated gradient of the potential field.
- A **navigation function** is an ideal potential field that
  - has global minimum at the goal
  - has no local minima
  - grows to infinity near obstacles
  - is smooth

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## Question

- Can such an ideal potential field be constructed efficiently in general?

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## Completeness

- A **complete motion planner** always returns a solution when one exists and indicates that no such solution exists otherwise.
  - Is the visibility graph algorithm complete? Yes.
  - How about the exact cell decomposition algorithm and the potential field algorithm?

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