

## MA 5219 - Logic and Foundations of Mathematics 1

Course-Webpage <http://www.comp.nus.edu.sg/~fstephan/mathlogic.html>

Homework due in Week 10, Tuesday 22 October 2013.

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Hand in each homework which you want to be checked; 1 mark per each correct starred homework; up to 10 marks in total for homework - there will be more than 10 starred homeworks, so you have several chances to try.

**10.1\* Birkhoff rules.** Adjust the axioms of the Birkhoff rules to incorporate binary associative operations like  $\circ$  as follows (where the fourth rule is given for a function  $f$  of arity 3 but applies for functions of all arities):

$$\begin{array}{c} \frac{\emptyset}{t = t} \quad \frac{s = t}{t = s} \\ \frac{s = t, t = r}{s = r} \quad \frac{s_1 = t_1, s_2 = t_2, s_3 = t_3}{f(s_1, s_2, s_3) = f(t_1, t_2, t_3)} \\ \frac{\emptyset}{(r \circ s) \circ t = r \circ s \circ t} \quad \frac{\emptyset}{r \circ (s \circ t) = r \circ s \circ t} \\ \frac{s = t}{s \circ r = t \circ r} \quad \frac{s = t}{r \circ s = r \circ t} \\ \frac{s = t}{q \circ s \circ r = q \circ t \circ r} \quad \frac{s = t}{s^\sigma = t^\sigma} \text{ for all global substitutions } \sigma \end{array}$$

In the following let  $x, y, z$  be variables and  $d, e$  be constants.

(a) Assume now that the additional axioms  $x \circ x \circ x \circ x \circ x \circ x = e$ ,  $e \circ x = x$ ,  $x \circ e = x$  and  $d \circ d \circ d \circ d \circ d = e$  are given. Prove that  $d = e$ .

(b) Assume now that the additional axioms  $x \circ x \circ x \circ x \circ x \circ x = e$ ,  $e \circ x = x$ ,  $x \circ e = x$  and  $d \circ d \circ d = e$  are given. Prove that one can neither derive  $d = e$  nor derive  $x \circ y = y \circ x$  by producing a model in which these two formulas are false.

(c) Assume now that the additional axioms  $x \circ d = d$ ,  $x \circ e = x$ ,  $e \circ x = x$ ,  $f(x \circ y) = f(y) \circ f(x)$ ,  $f(f(x)) = x$  and  $x \circ x \circ f(x) = x$  are given, where  $f$  is a unary function symbol. Derive  $f(x) \circ x \circ x = x$ . Furthermore, provide a model which shows that  $f(x) = x$  cannot be derived.

**10.2\* Second Order Logic.** Second order language permits to quantify over sets. Using that there are uncountably many subsets of the natural numbers, give a set  $X$  of formulas such that every second order model of  $X$  is uncountable.  $X$  should of course be satisfiable, that is, have at least one model. It is permitted to introduce constant symbols, predicate symbols and function symbols; furthermore, the symbol  $\in$  in the formula  $a \in B$  is true iff  $a$  is an element of  $B$ ; lower case variables range over elements and upper case variables range over sets.