Lecture 11 Networked Game Traffic and Transport Protocol

Assignment 1

isolate traffic payload size histogram activity pattern periodic pattern MOBA RTS FPS RPG





UDP

bandwidth in: 40 - 160 kbps out: 15 - 40 kbps payload size in: 100 - 450 bytes out: 50 - 150 bytes

packet rate 25 - 30 packets/seconds gap in between: 30-40ms



TCP

lower bandwidth in: 4.8 kbps out: 6 kbps

an order of magnitude smaller than MOBA!

payload size < 25 bytes

an order of magnitude smaller than MOBA!

packet rate ~10 packets/seconds



















UDP

bandwidth in: 20 - 100 kbps out: 8 - 100 kbps

comparable to MOBA, with higher outgoing throughput

slightly larger payload size in: 50 - 300 bytes out: 30 - 100 bytes

smaller packet rate in: 20 - 120 packets/seconds out: 10 - 90 packets / seconds

I expected this to be smaller

RPG













TCP

much lower bandwidth in: 5 - 16 kbps out: 1 - 8 kbps

larger payload size in: 100 - 300 bytes out: 20 - 160 bytes

smaller packet rate in: 1 - 15 packets/seconds out: 1 - 15 packets / seconds





1.6 packet / second

2 packet / second

What you found:

RPG have smaller packets and smaller update rate.

what about periodicity?

For many games, server updates are periodic. (50 - 200ms interval)



low bandwidth small packets low frequency predictable

Both UDP and TCP are used

TCP or UDP ?

Why use TCP?

- TCP provides reliable, in-order delivery
- TCP goes through most firewalls, UDP does not
- TCP manages connection for us

Why not to use TCP?

- TCP incurs higher latency
- Don't always need reliability and in-order delivery
- High header overhead

position = $10 \longrightarrow 10 \longrightarrow 13 \longrightarrow 13$ position = $13 \longrightarrow 13 \longrightarrow 13$

Updated position not delivered to application until (outdated) lost packet is received

A's position = $10 \longrightarrow$ B's position = $13 \longrightarrow X$ C's position = $15 \longrightarrow$

Some messages need not be delivered in sequence.



Gestures from someone far away need not be received reliably.

TCP header is >= 20 bytes high overhead for small packets (46% in Shenzhou Online)

https:// ()/lsalzman/enet

Example of a library that provides reliability, sequencing, connection managements over UDP

Delivery can be stream-oriented (like TCP) or message-oriented (like UDP)

Supports partial reliability

enet_packet_create ("abc",
4, ENET_PACKET_FLAG_RELIABLE)

Retransmission triggered by timeout-based on RTT

Data in queue are bundled into one packet if there is space

enet

Portable, easy to use, but still, most firewalls block UDP traffic

Need to study the use of TCP for networked games

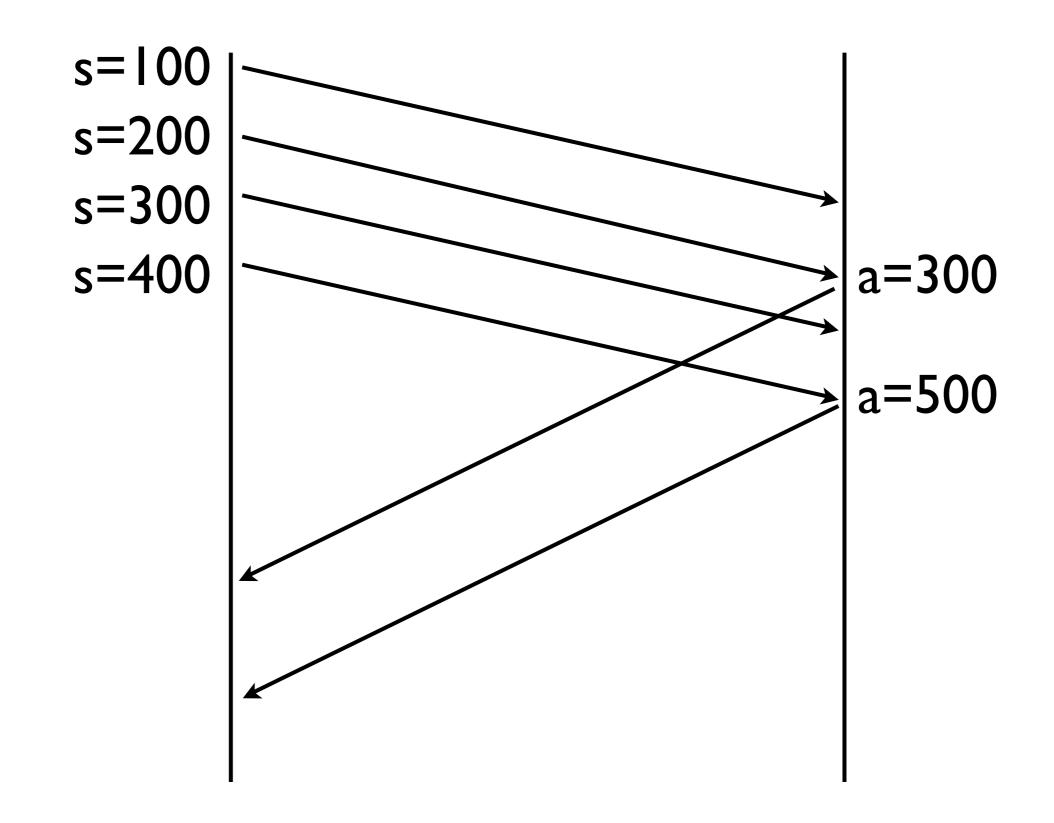
Lessons are still useful to build enet-like UDP library

How slow is TCP, really?

Which part of TCP is the root of slowness?

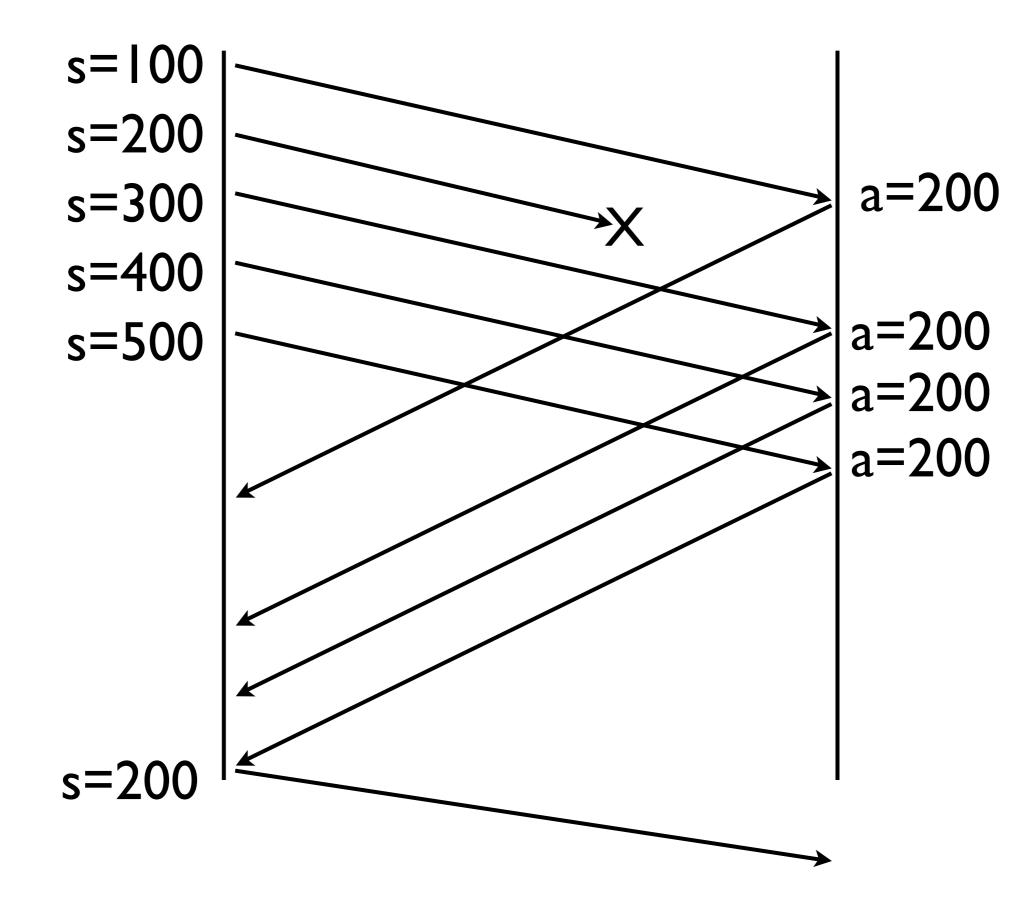
Can we fix TCP?

A Quick Review of TCP



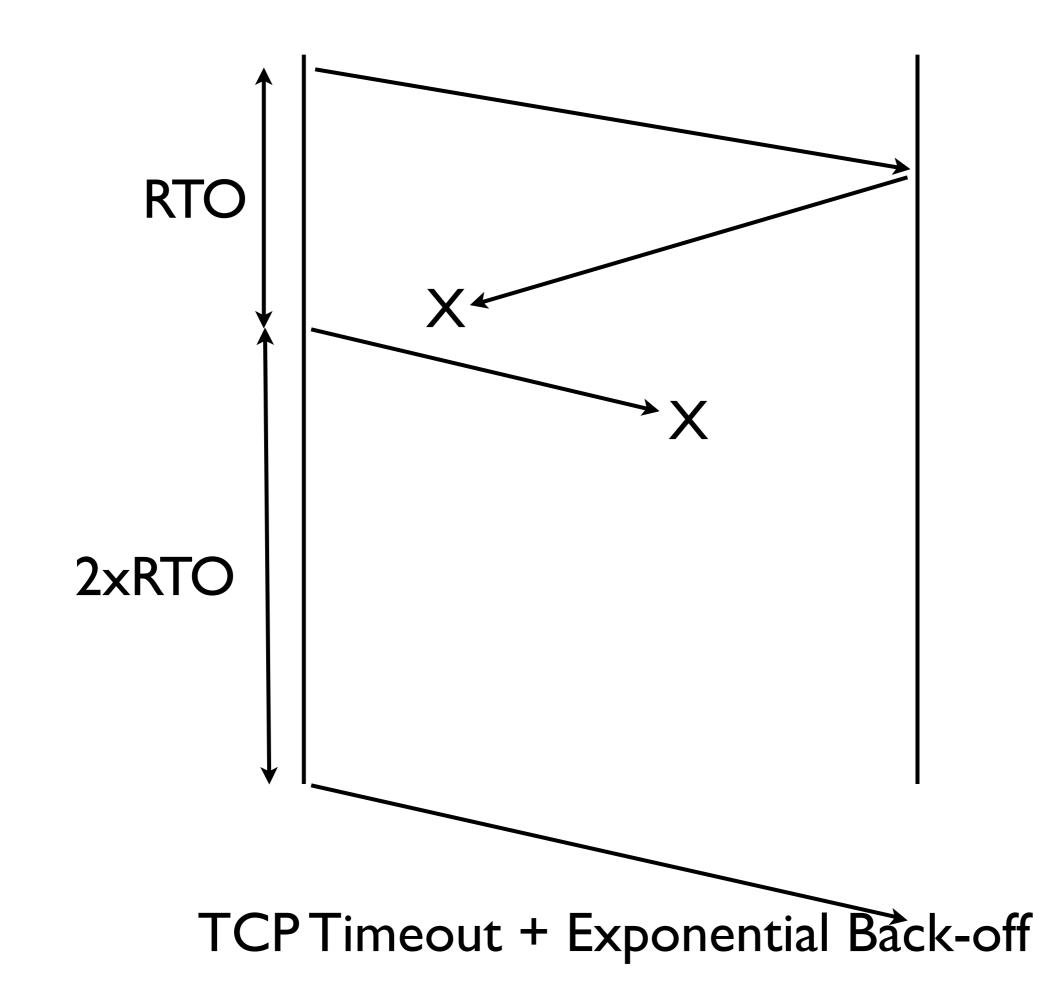
TCP Delayed ACK

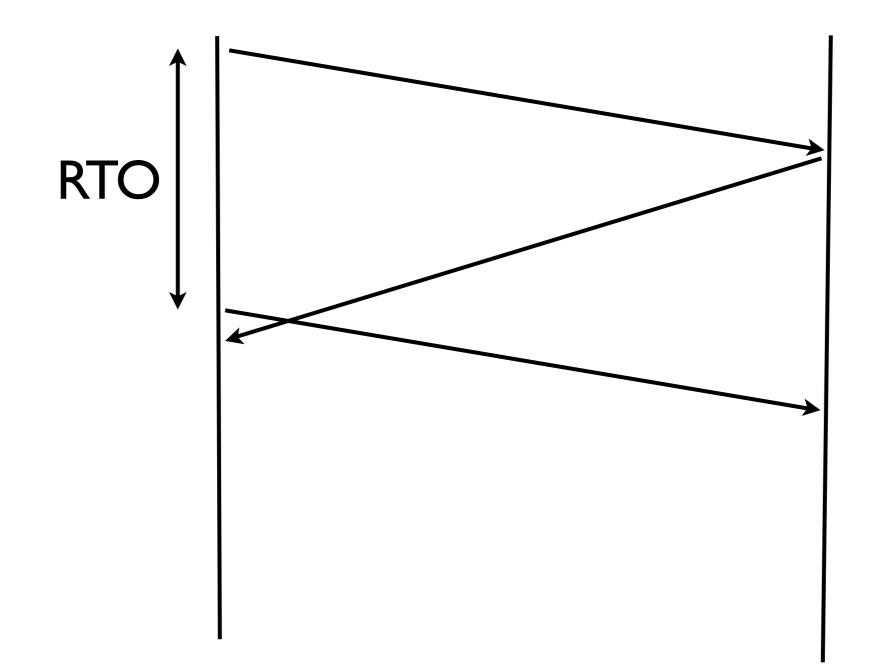
TCP Spec: max **500ms** delay Most implementation: **200ms**



3 dup ACKs within RTO - RTT: TCP Fast Retransmission

Definition of Dup ACKs in 4.4BSD and Stevens: "pure ACK with no data"





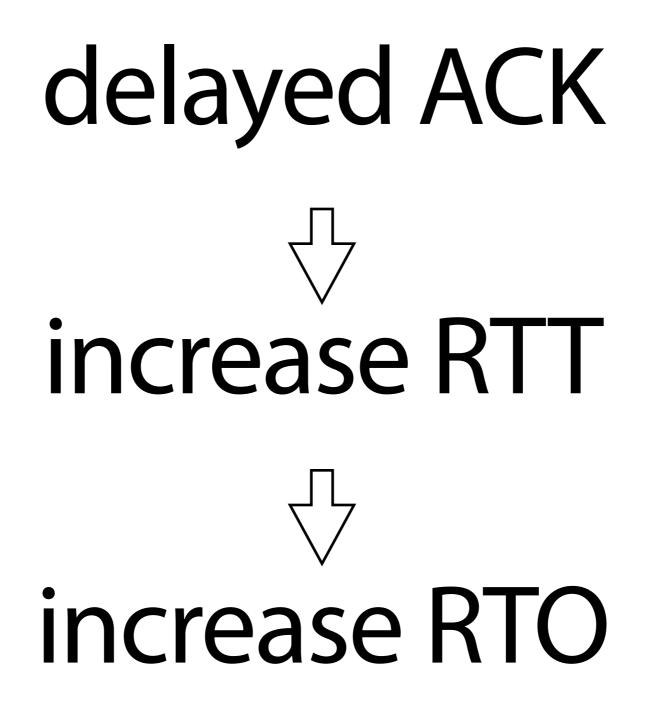
Spurious Retransmission

RTO estimation

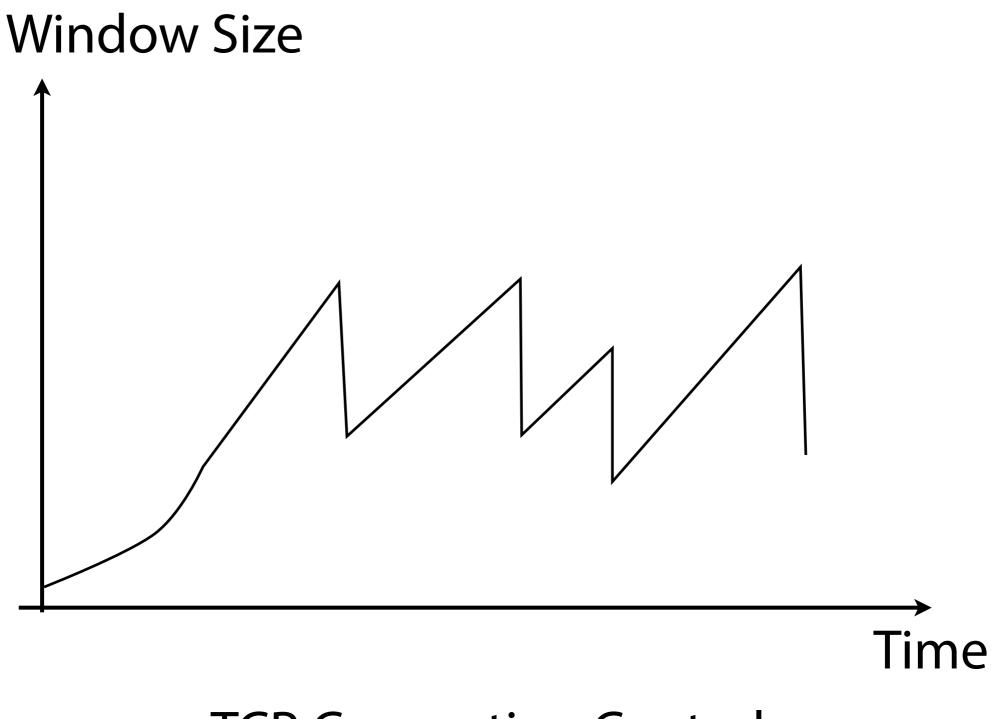
 $E_{i} = 7E_{i-1}/8 + RTT/8$ $V_{i} = 3V_{i-1}/4 + |RTT-E_{i-1}|/4$ $RTO = max(E_{i} + 4V_{i}, 1s)$

Linux's RTO estimation

 $E_i = 7E_{i-1}/8 + RTT/8$ $V_i = 3V_{i-1}/4 + |RTT-E_{i-1}|/4$ $W_i = min(V_i, 50ms)$ $RTO = max(200ms, E_i+W_i)$



Congestion Control



TCP Congestion Control

Congestion window resets to 2 after an idle period (> RTO)

What does real game traffic look like?

low packet rate **small** packet size

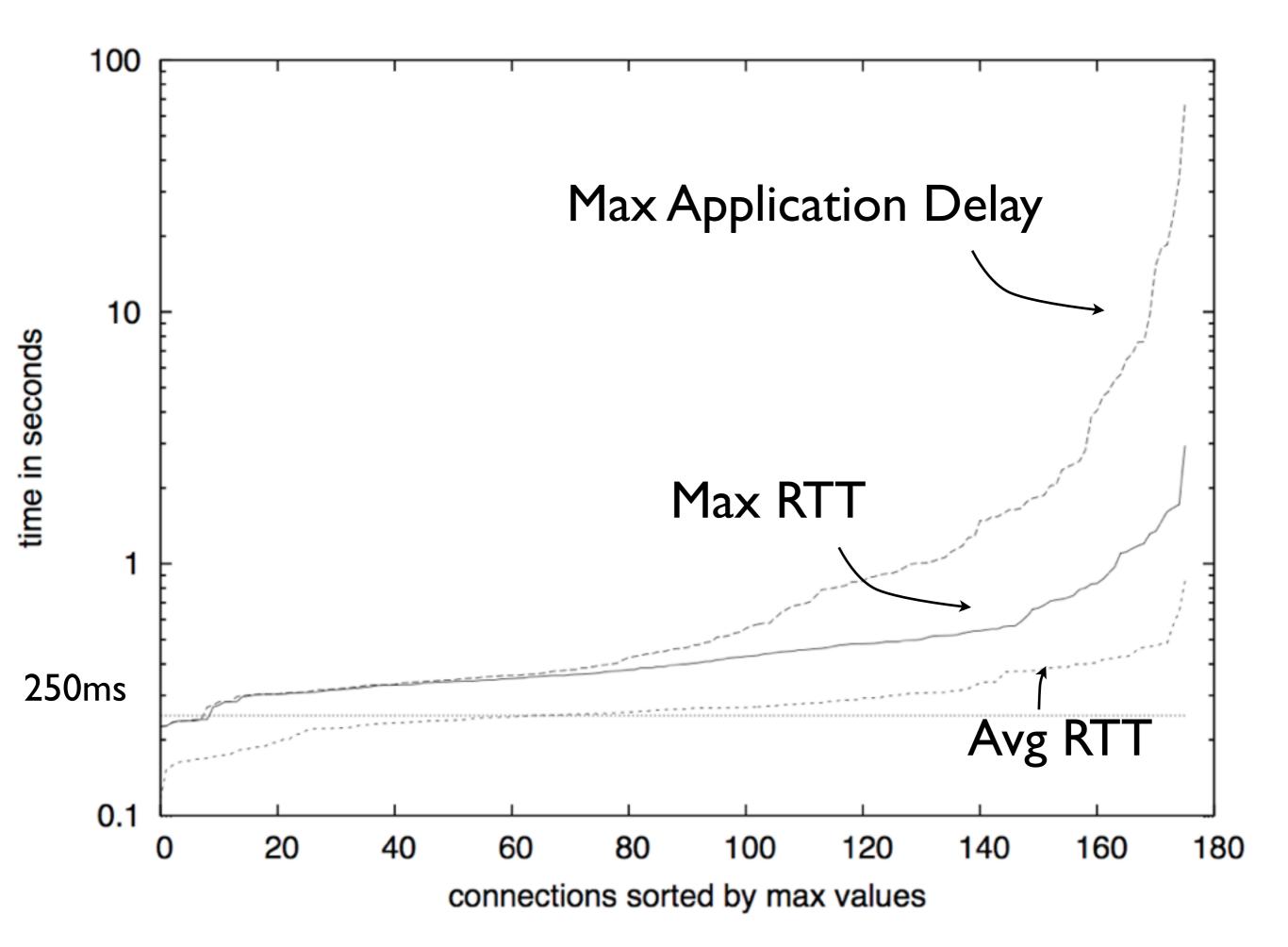
"Thin Streams"



A States

1 11

<07



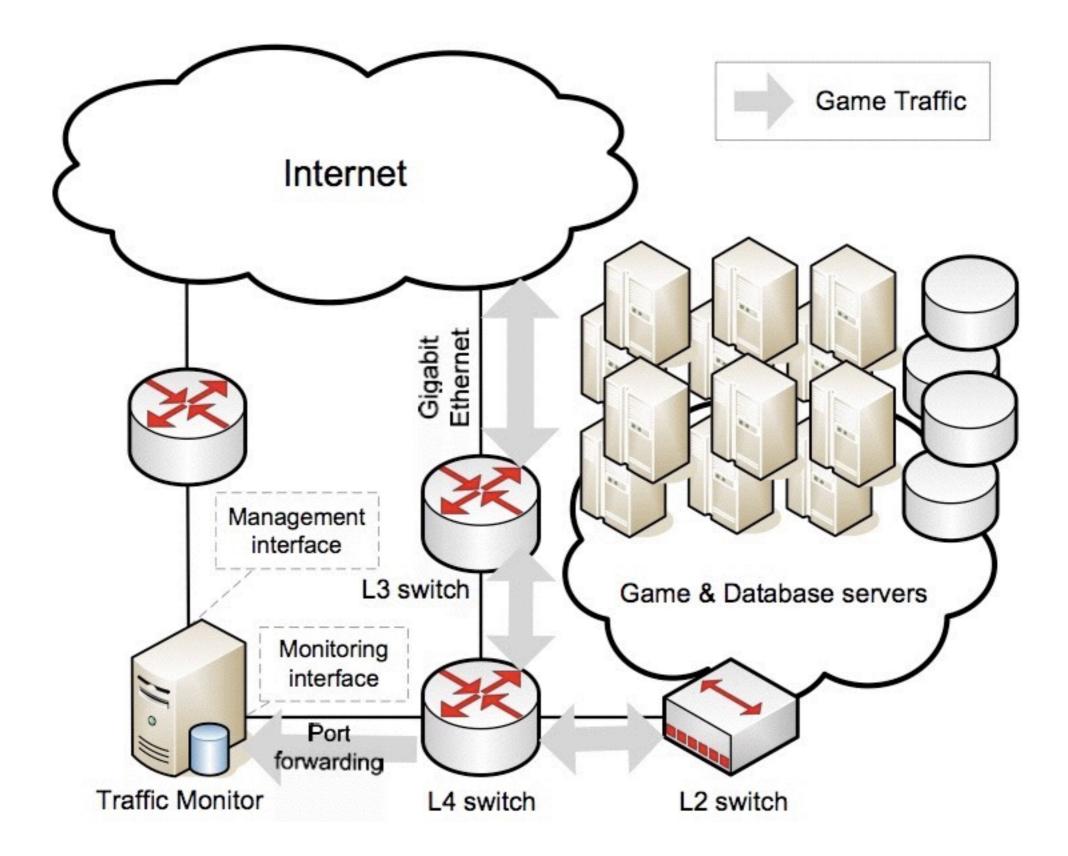
About 4 packets / sec

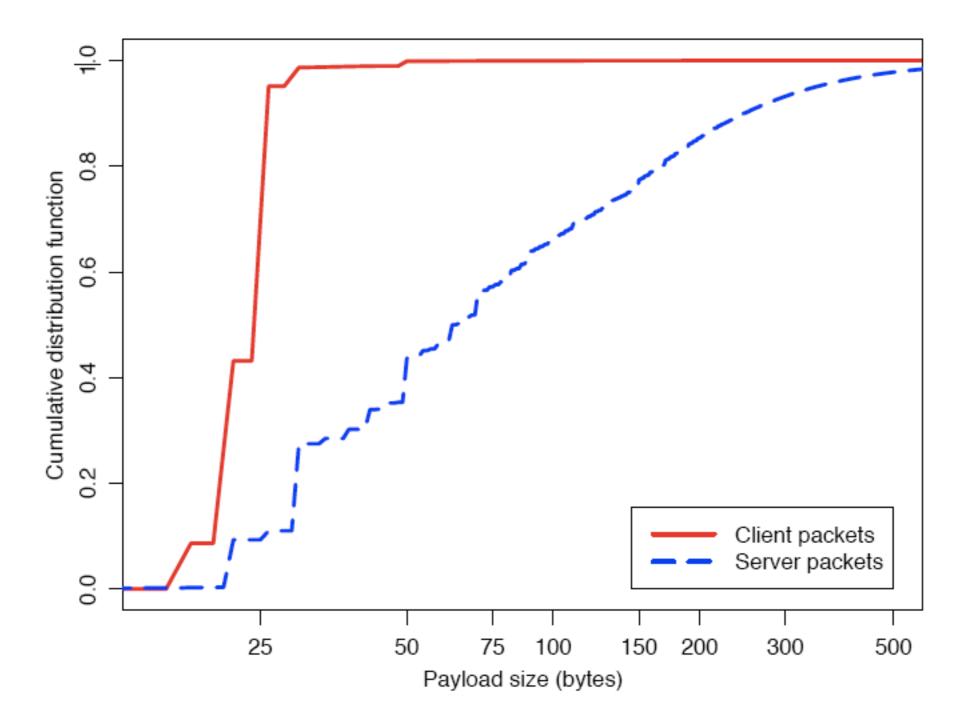
Average Payload: 100 Bytes

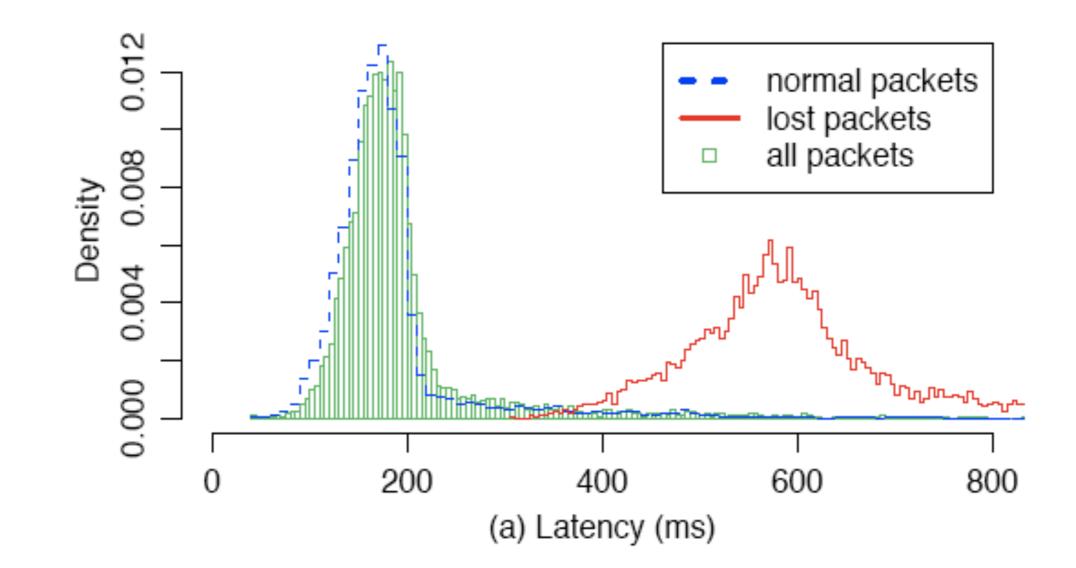
Loss Rate 1%

But some experience 6 retransmissions



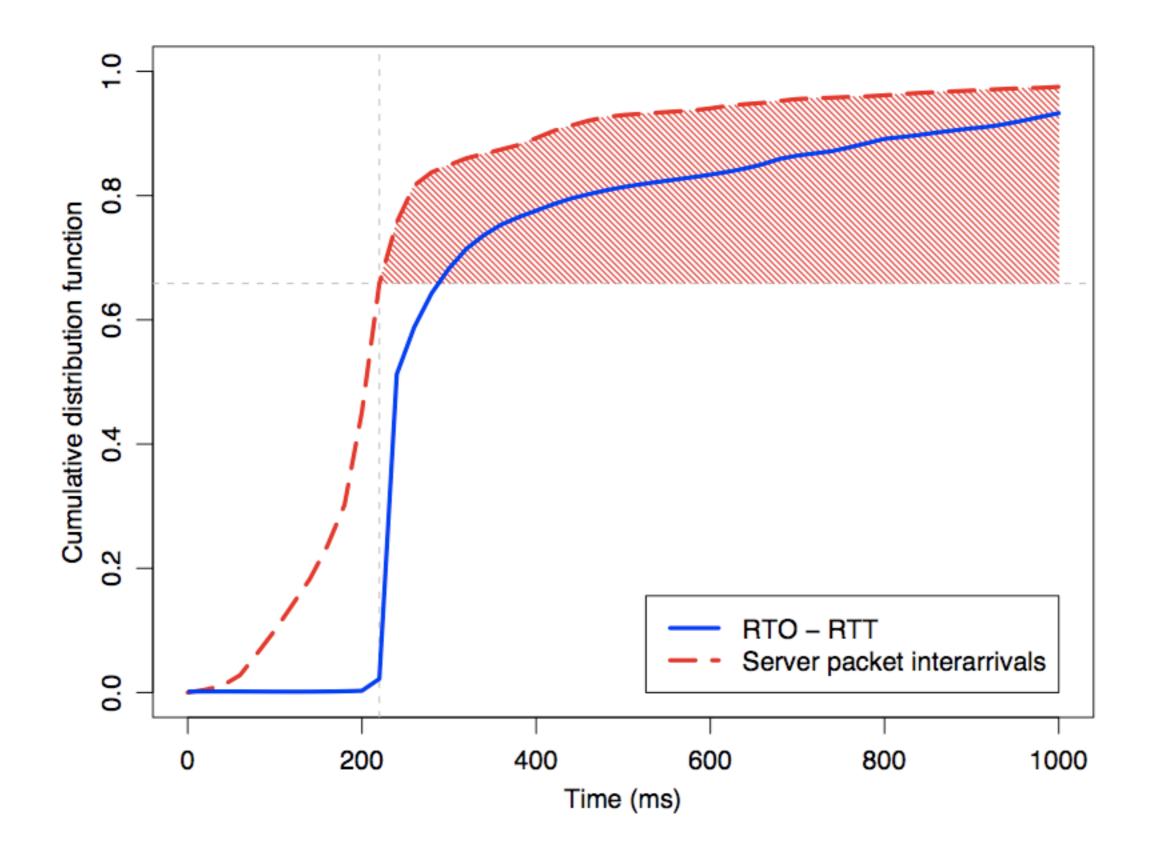






Findings 1: Fast retransmission rarely triggered

In ShenZhou Online traces, fail to trigger fast retransmission because insufficient dup ACK (50%) interrupted by data (50%)



Findings 2: Delay due mostly to timeout

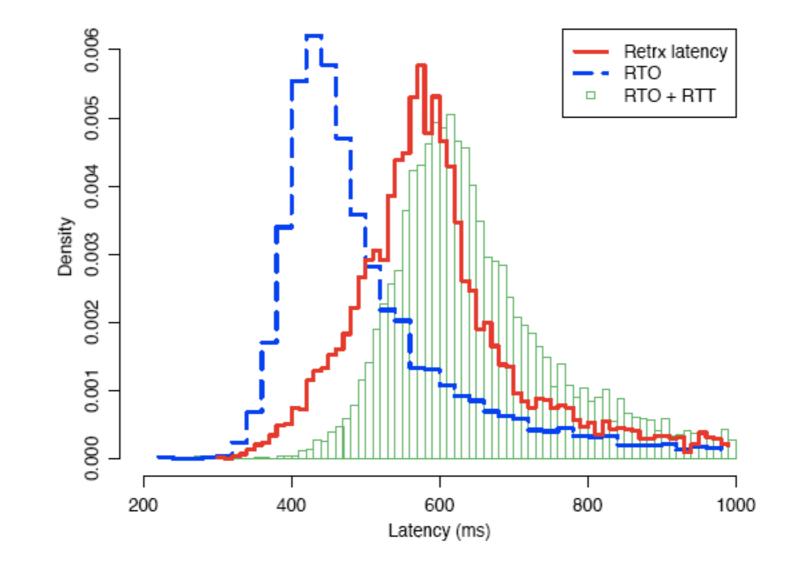
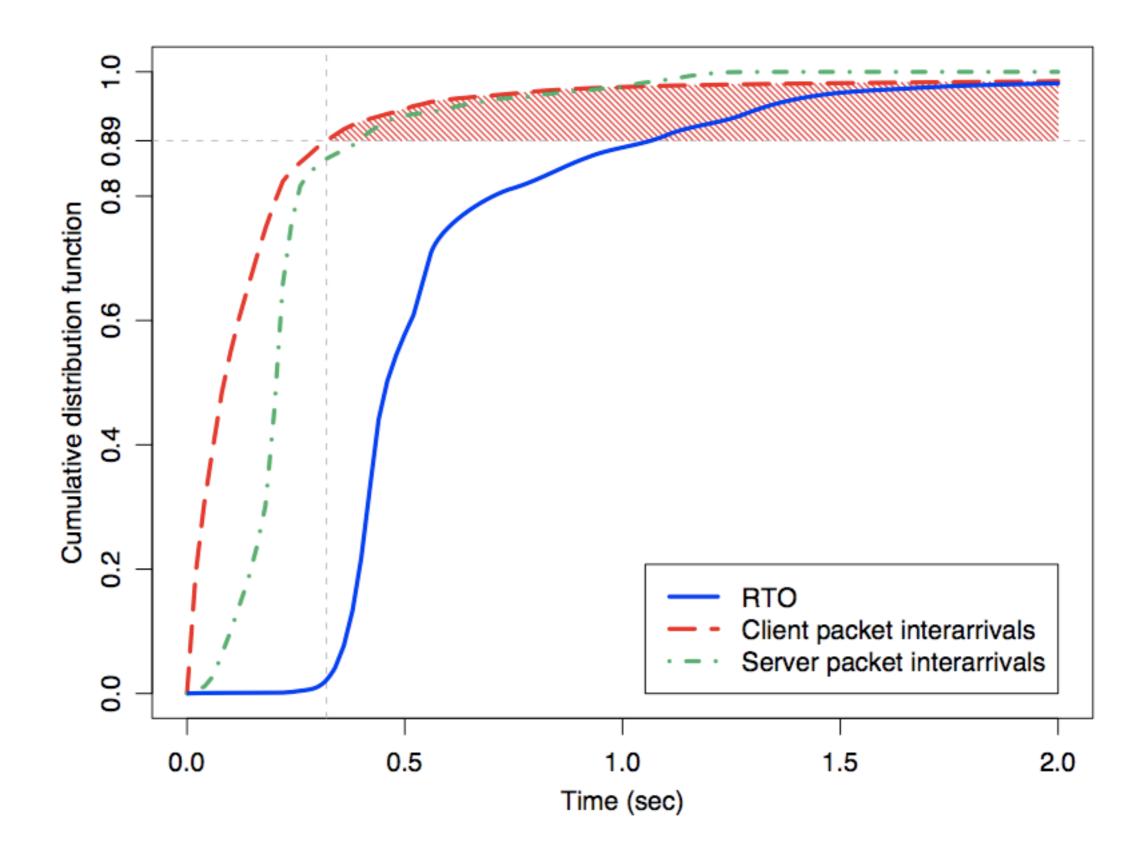


Figure 9: Average latency of dropped packets

Findings 3: Congestion window reset is frequent



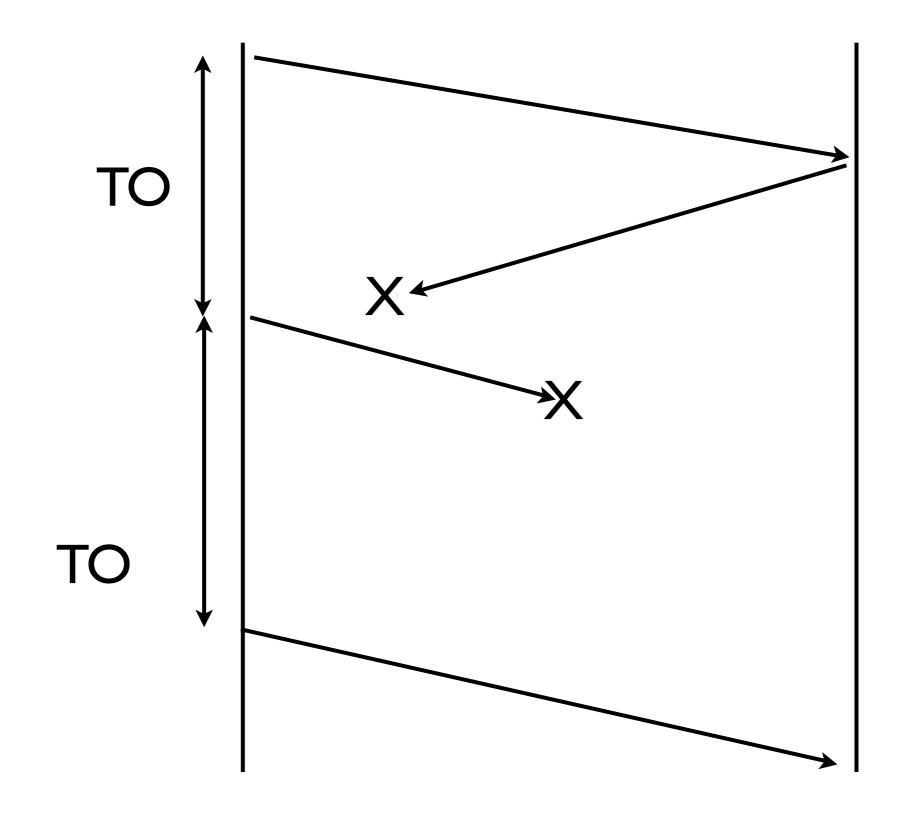
12% - 18% of packets faces window reset

think.. think.. think.. click (tank attack here) \longrightarrow click (missile launch there) \longrightarrow click (charge soldiers) \longrightarrow

The last command is delayed as congestion window = 2

How to make TCP (or, transport protocol) go faster in these games?

1. Remove exponential backoff



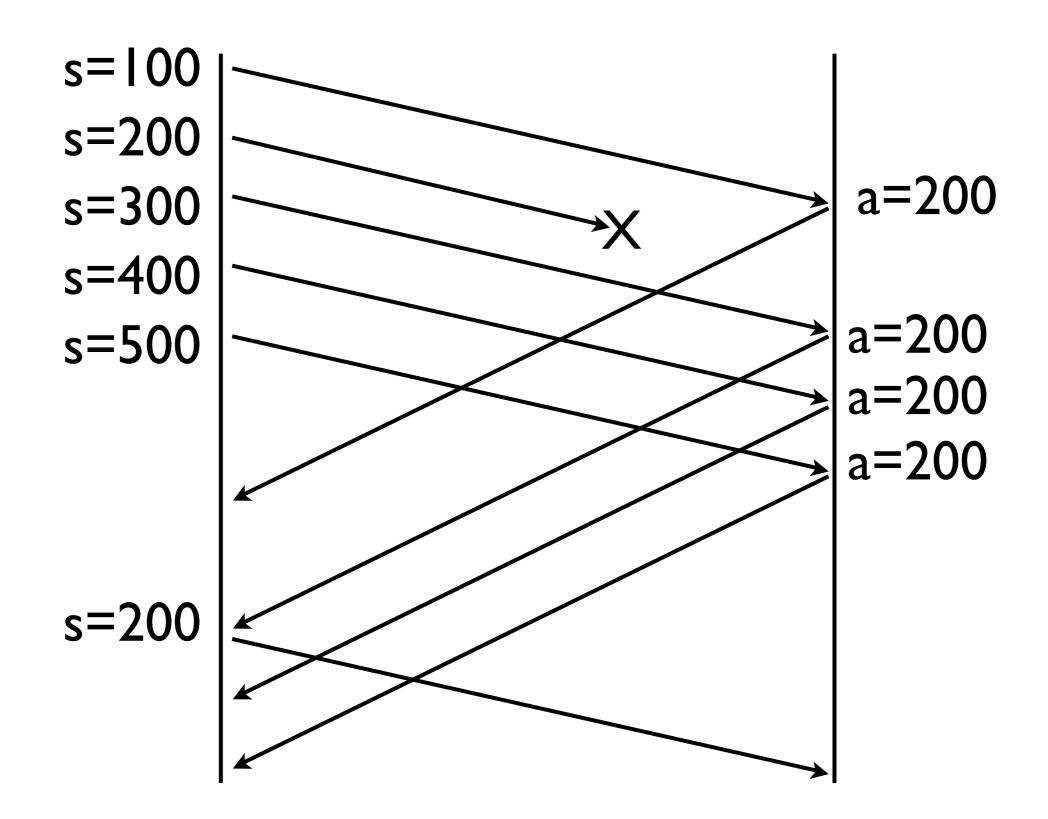
TCP Timeout

2. Make RTO Smaller

make sure minimum RTO is not 1s

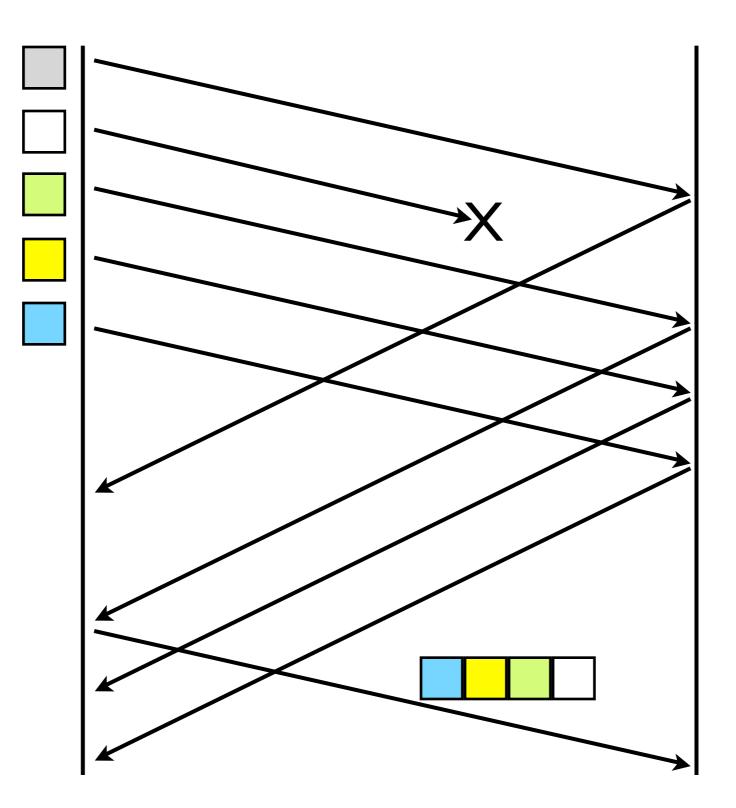
spurious retransmission is not disastrous

3. Make Fast Retransmit Faster



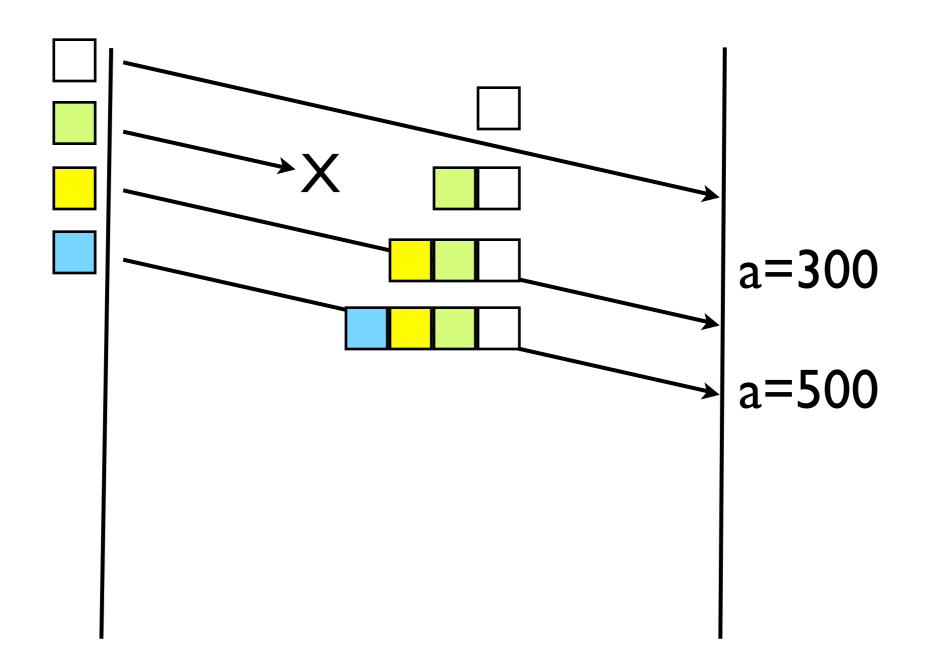
Retransmit after one duplicate ACK

4. Retransmission Bundling



Retransmit all unacknowledged data in queue

5. Redundant Data Bundling

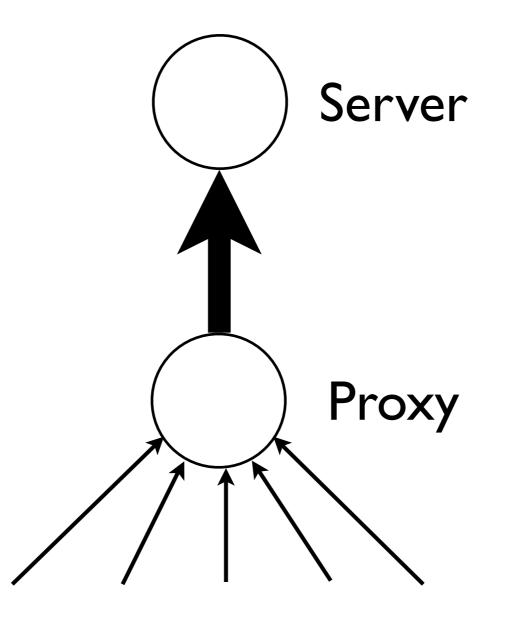


Send any unacknowledged segment in queue as long as there is space. Lost data gets recovered in the next transmission before retransmission.

6. Turn off or reduce Delayed ACKs

Packet interarrival time on average > 200ms (can't combine two ACKs into one anyway)

7. Combine Thin Streams into Thicker Stream



TCP for Games

- remove exponential backoff
- reduce RTO
- make fast retransmit faster
- retransmit aggressively
- don't delay ACK
- combine into thick streams

With Linux kernel, TCP_THIN_LINEAR_TIMEOUTS TCP_THIN_DUPACK