Scalable View-Dependent Progressive Mesh Streaming



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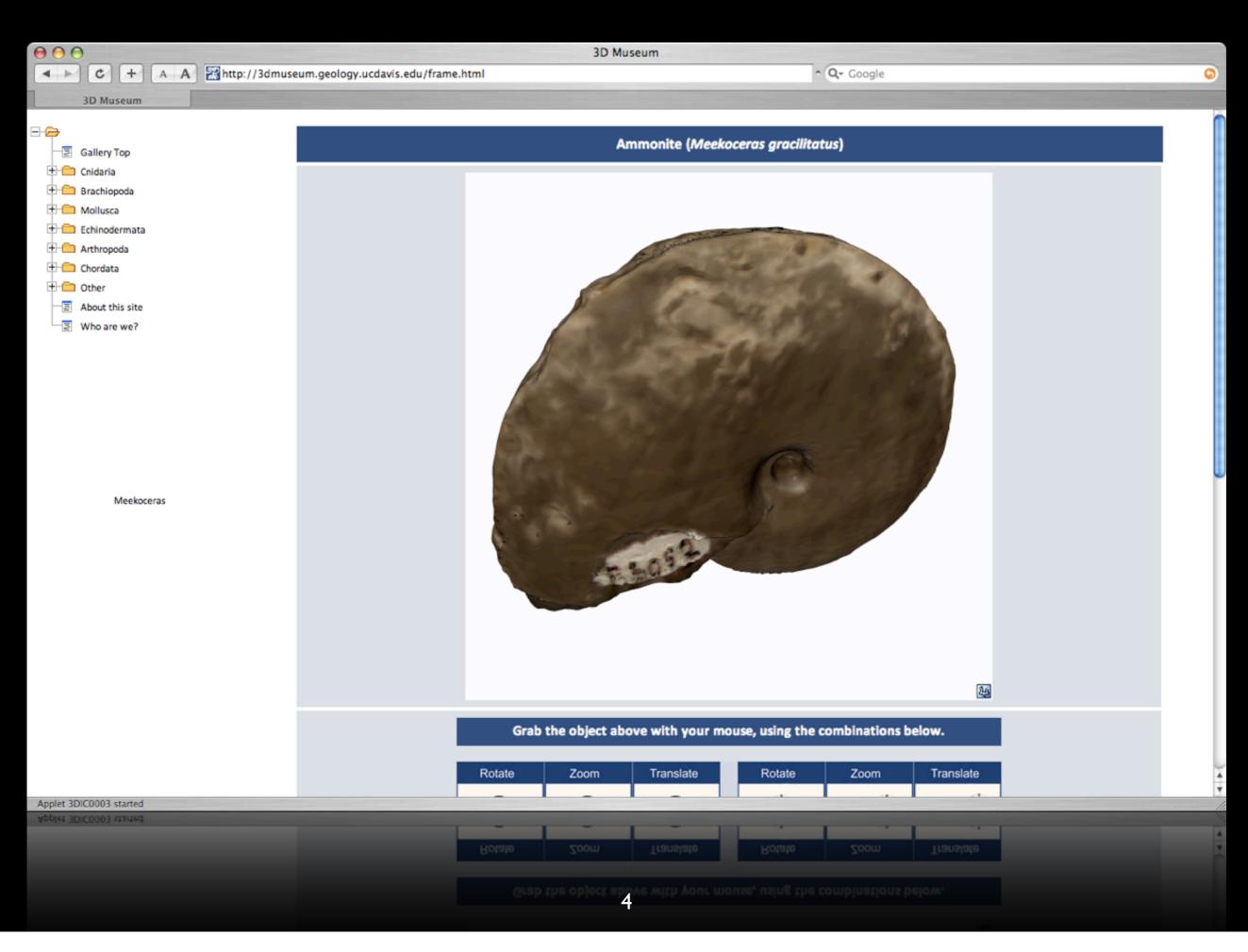
joint work with Cheng Wei National University of Singapore



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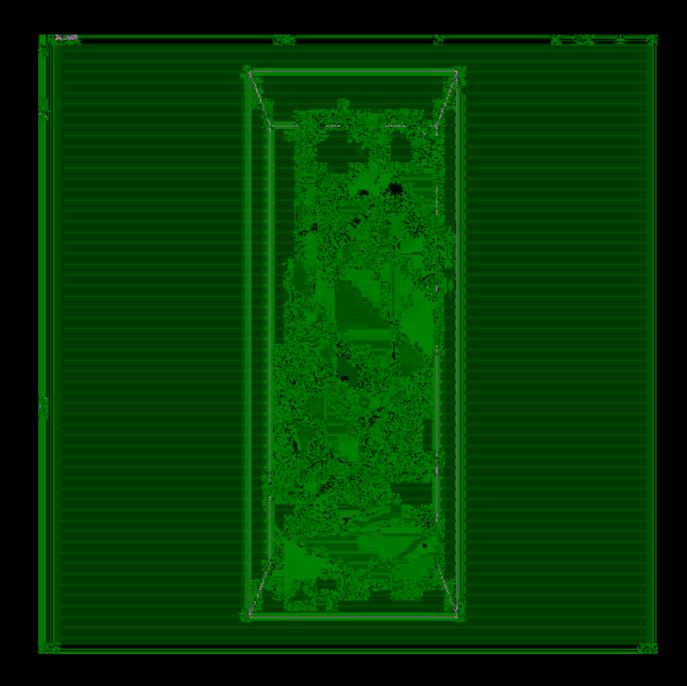




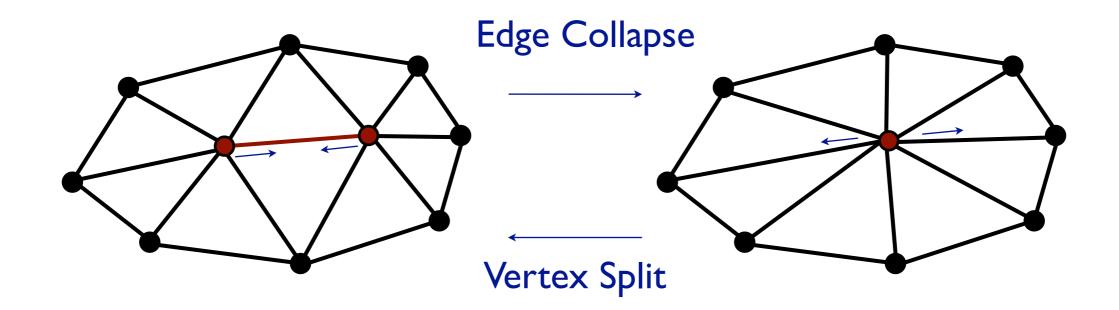
I0 MB



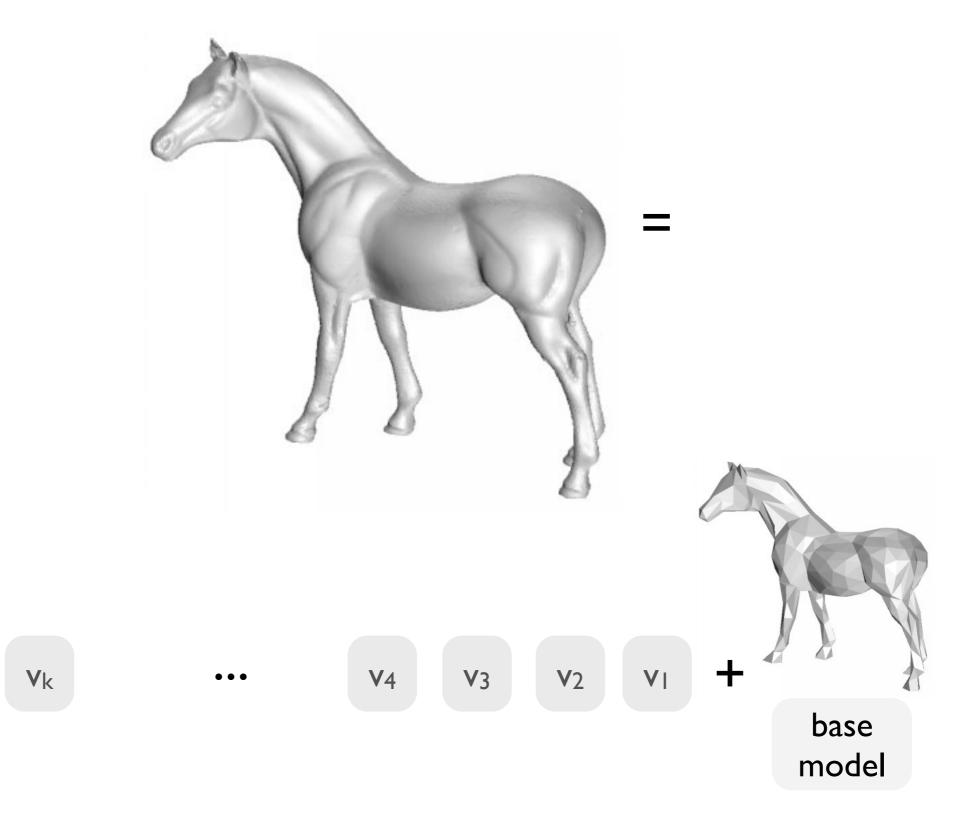
2 GB



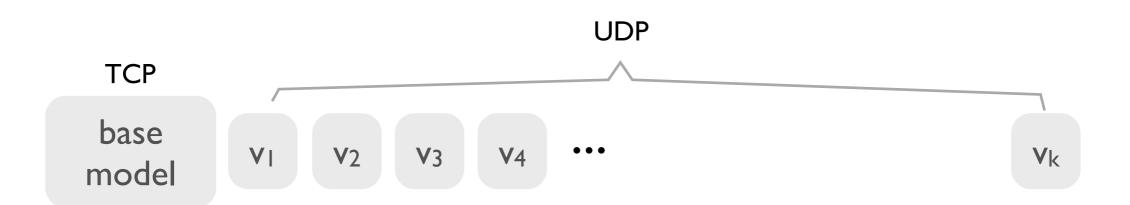
Hoppe's Progressive Mesh



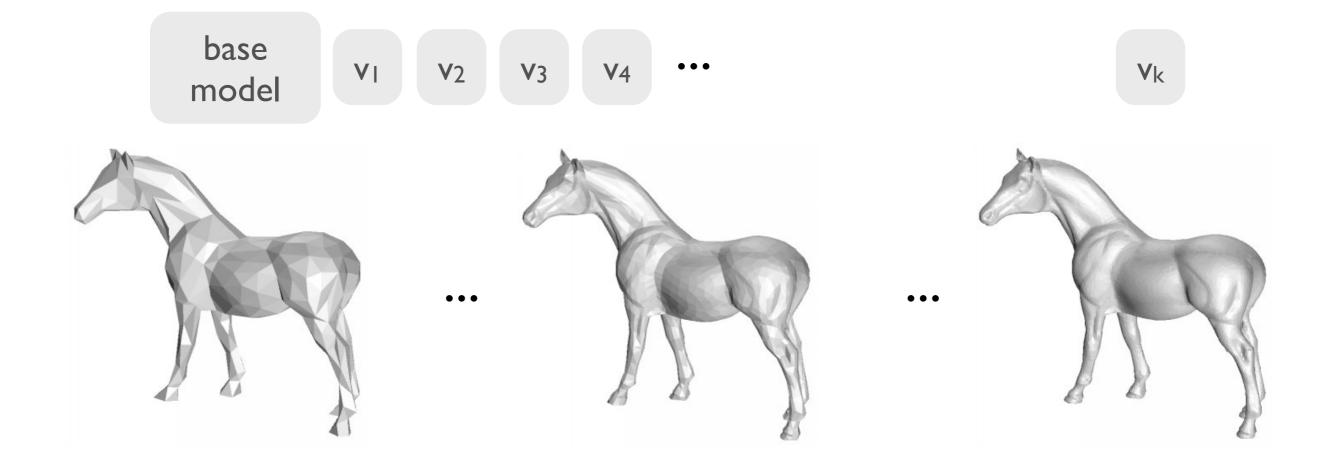
At the sender

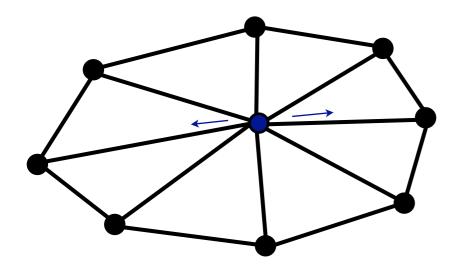


Transmission

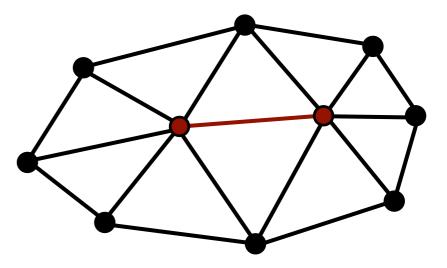


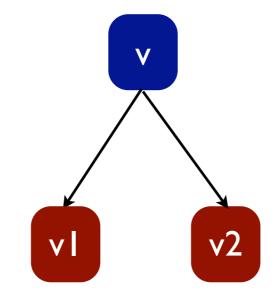
At the receiver

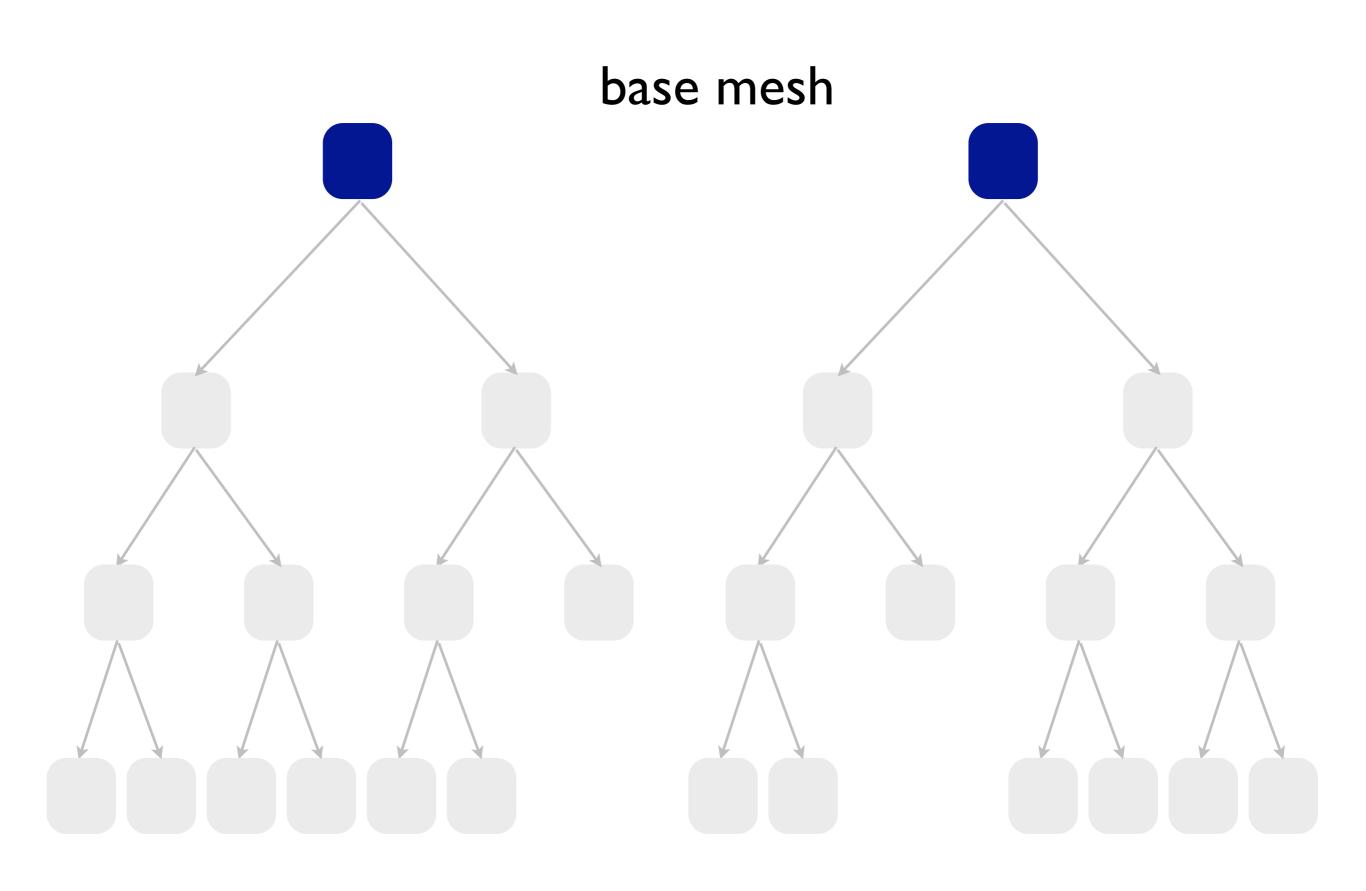


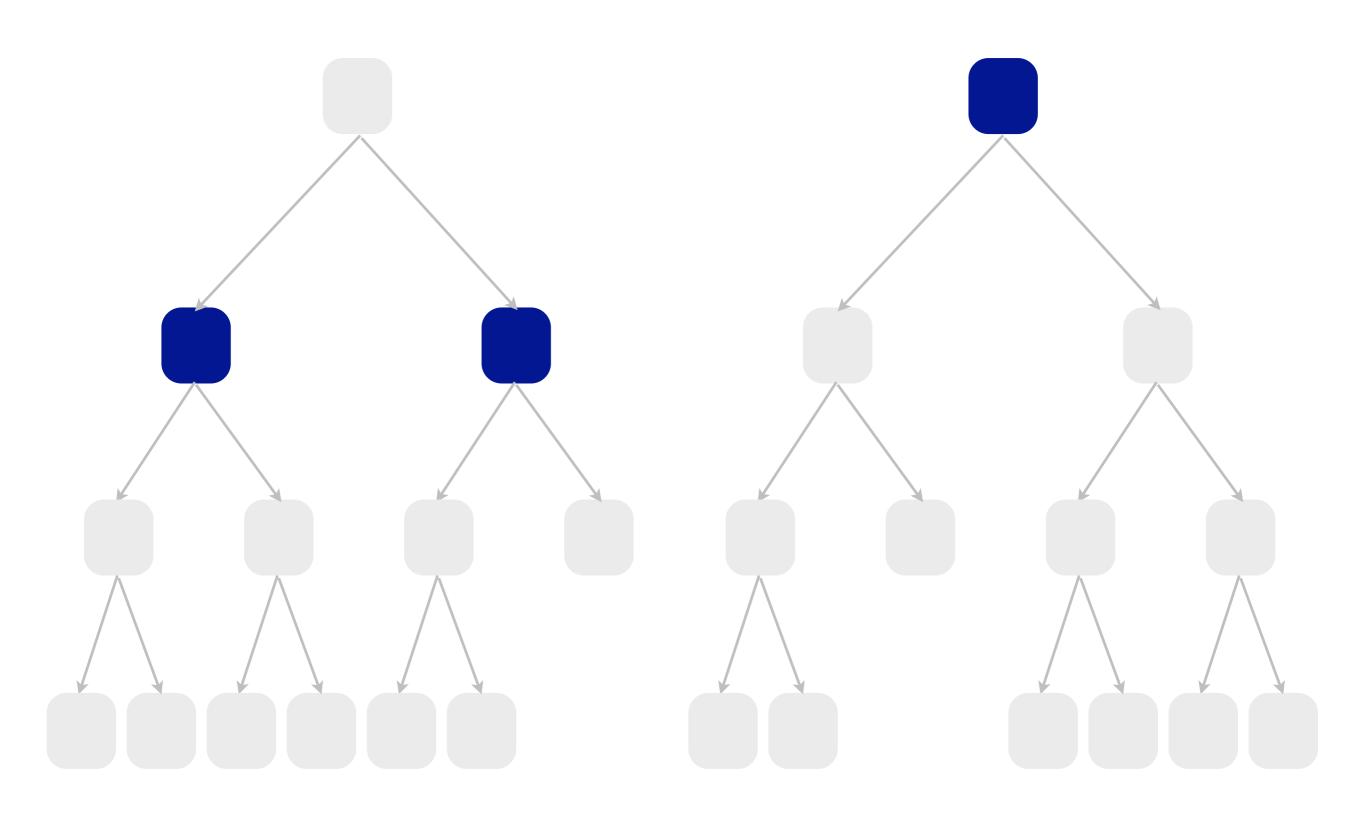


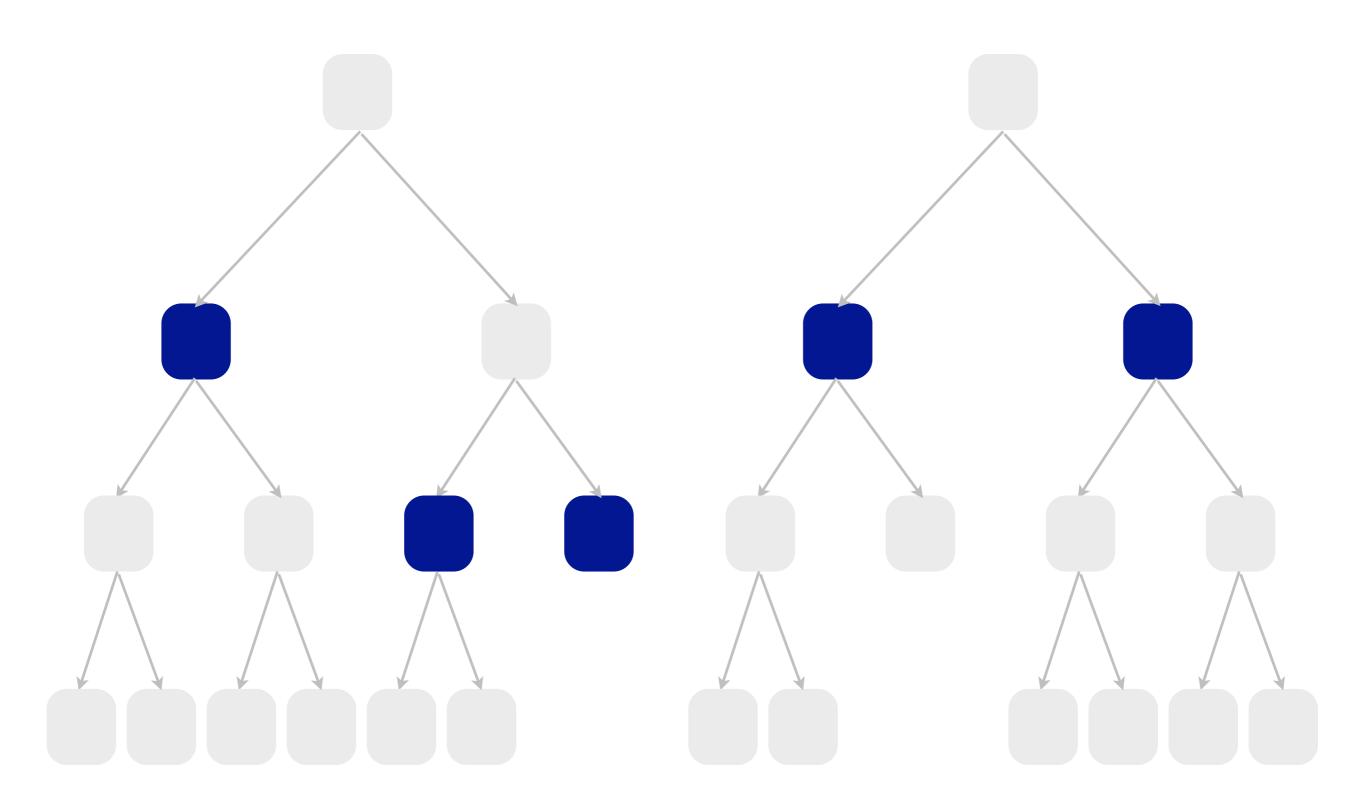
Vertex Split

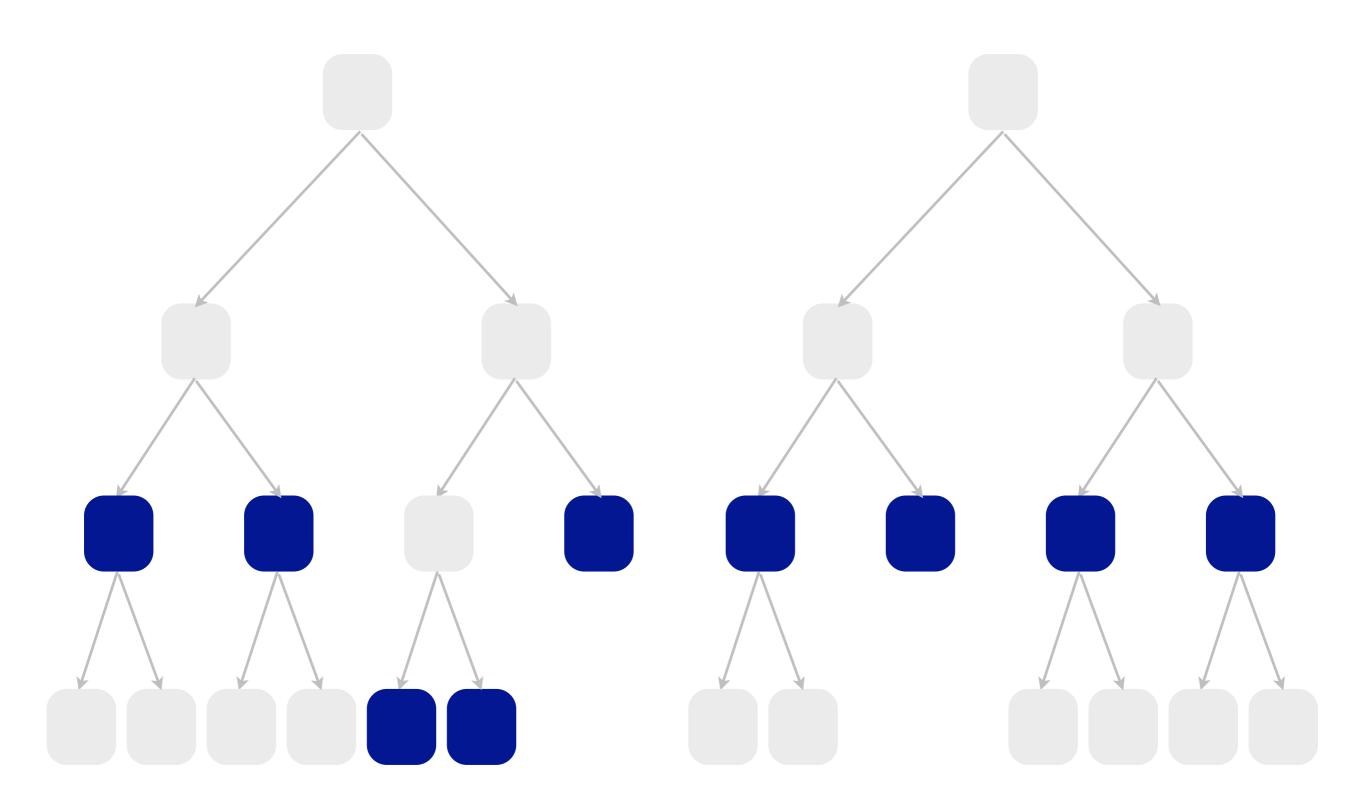


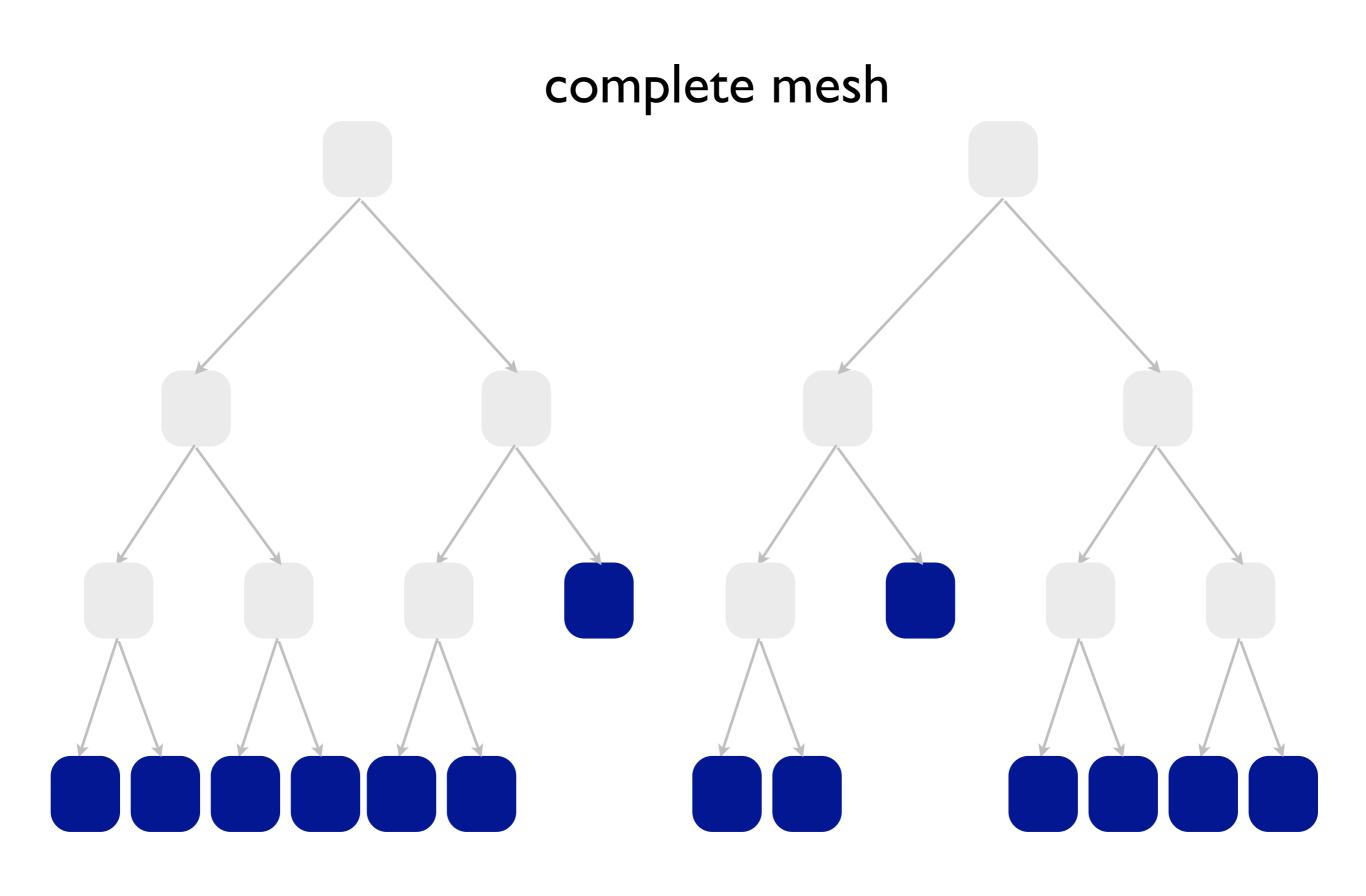








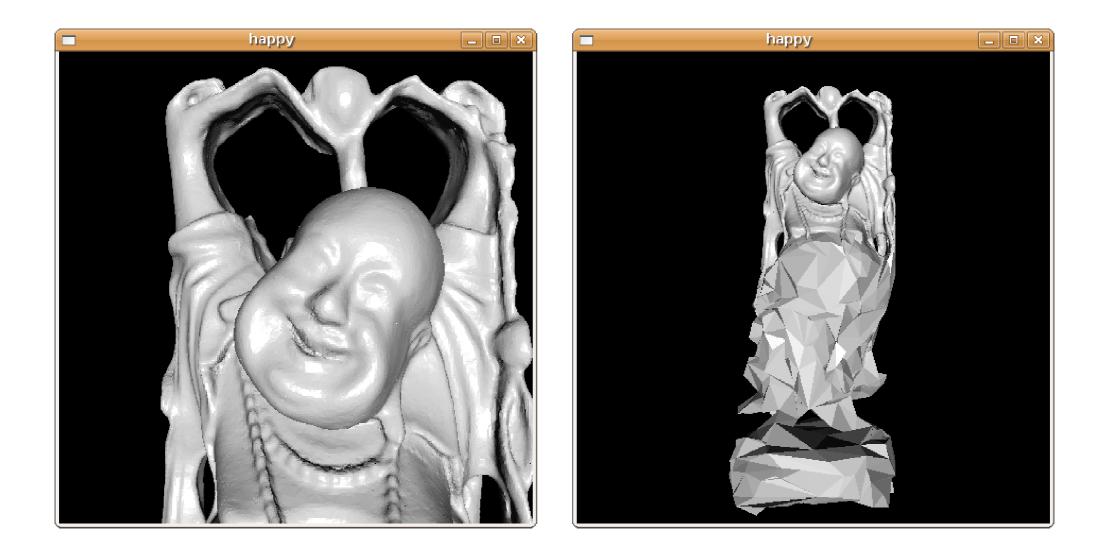


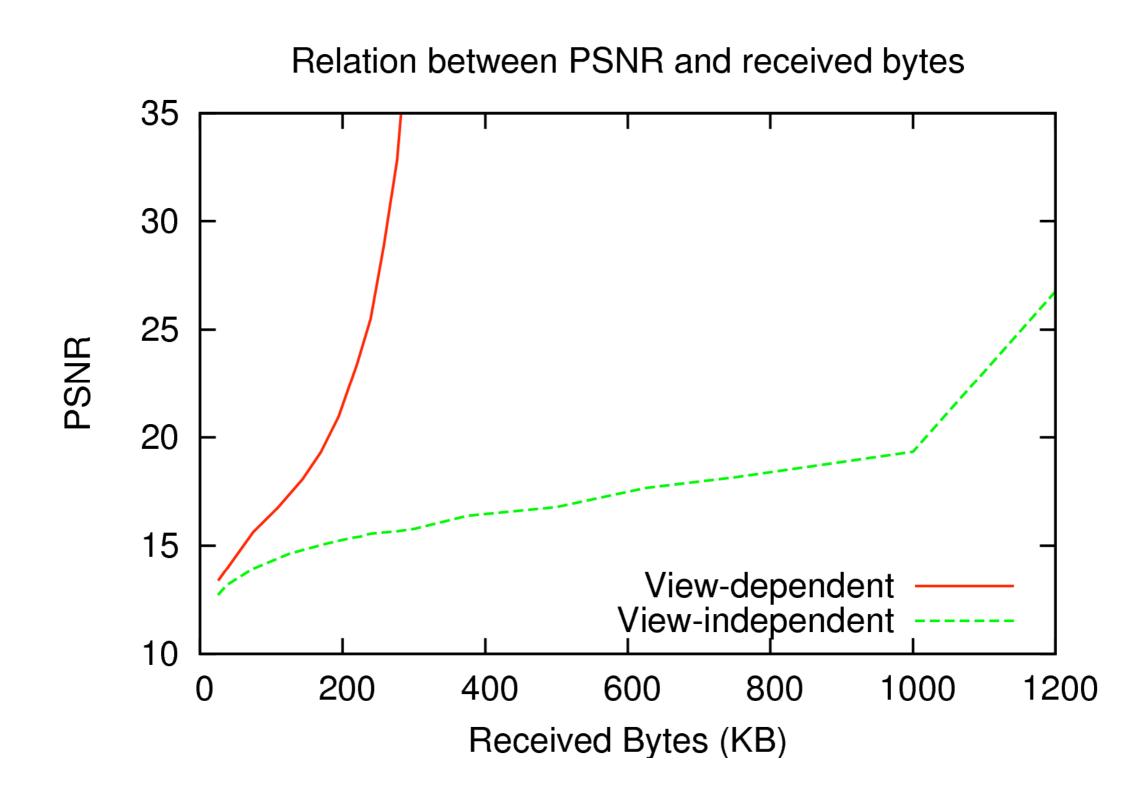


view-dependent streaming:

only send what the receiver can see







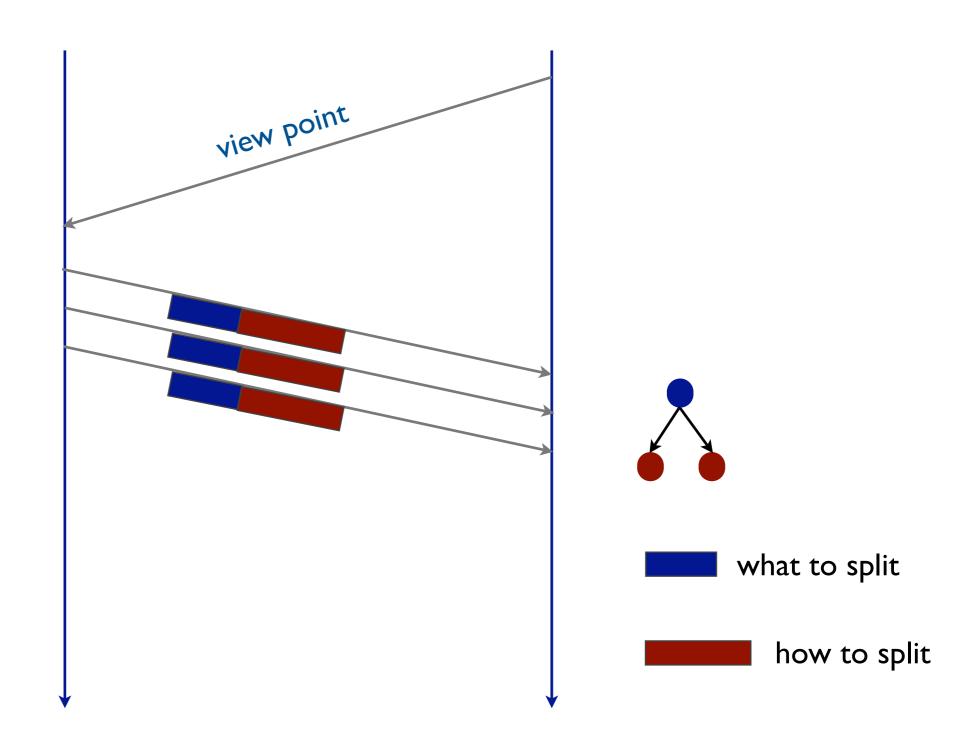
what to send?

in what order?

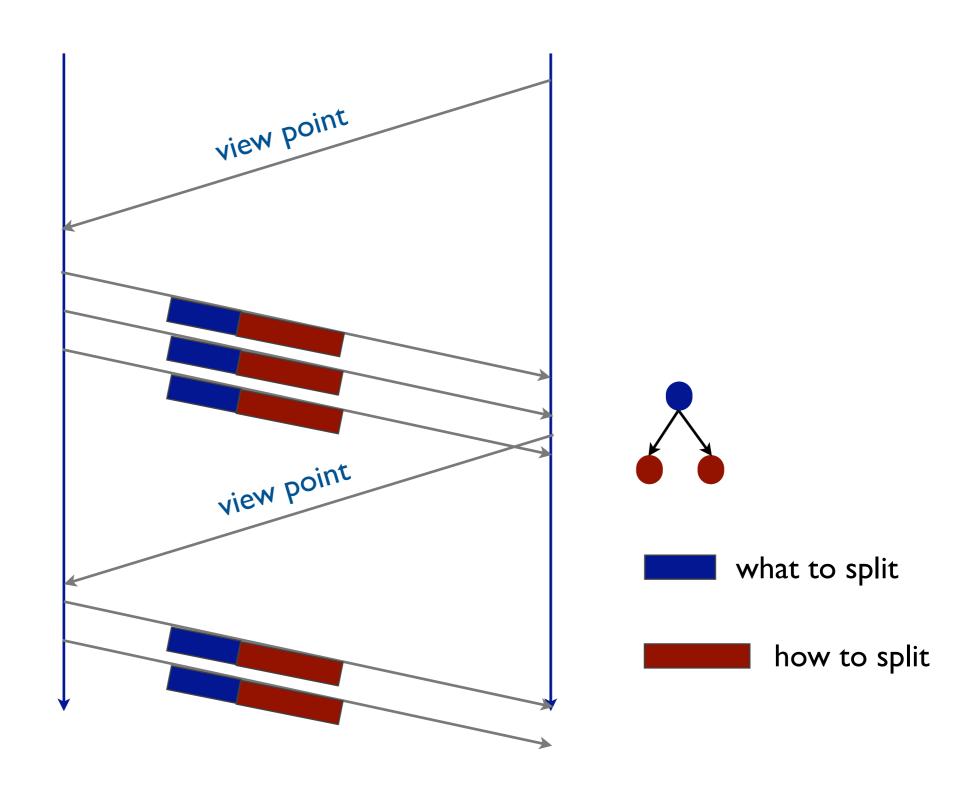
what to send? determined by view point

in what order? determined by visual contributions

Existing Approach



Existing Approach



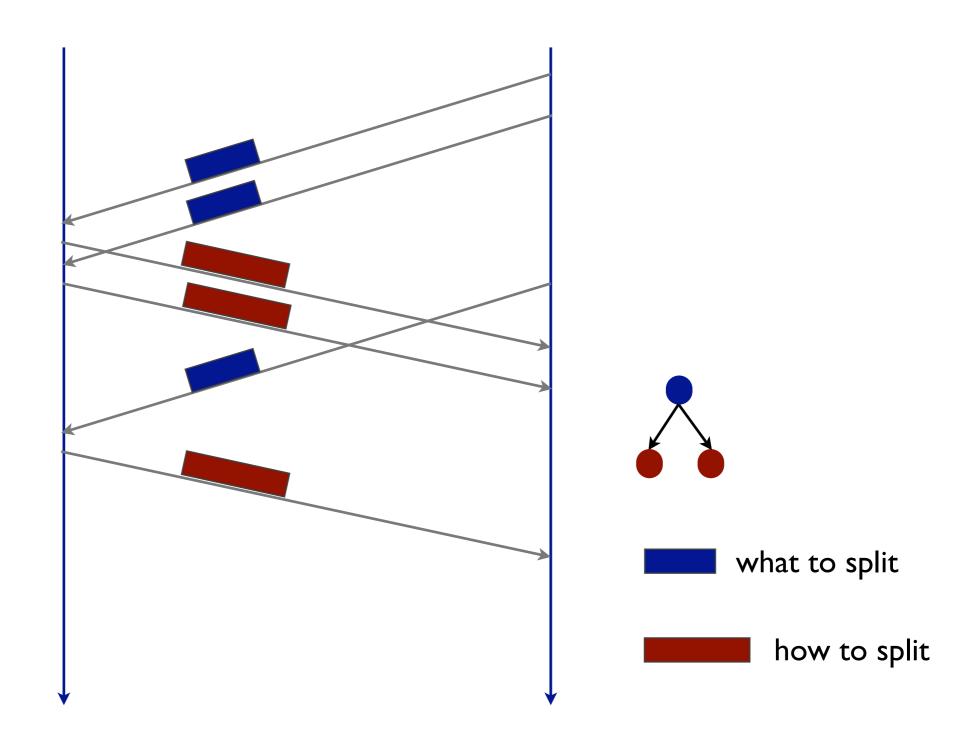
For each receiver, server needs to:

- compute visibility
- compute visual contribution of each vertex split
- sort vertex splits
- remember what has been sent

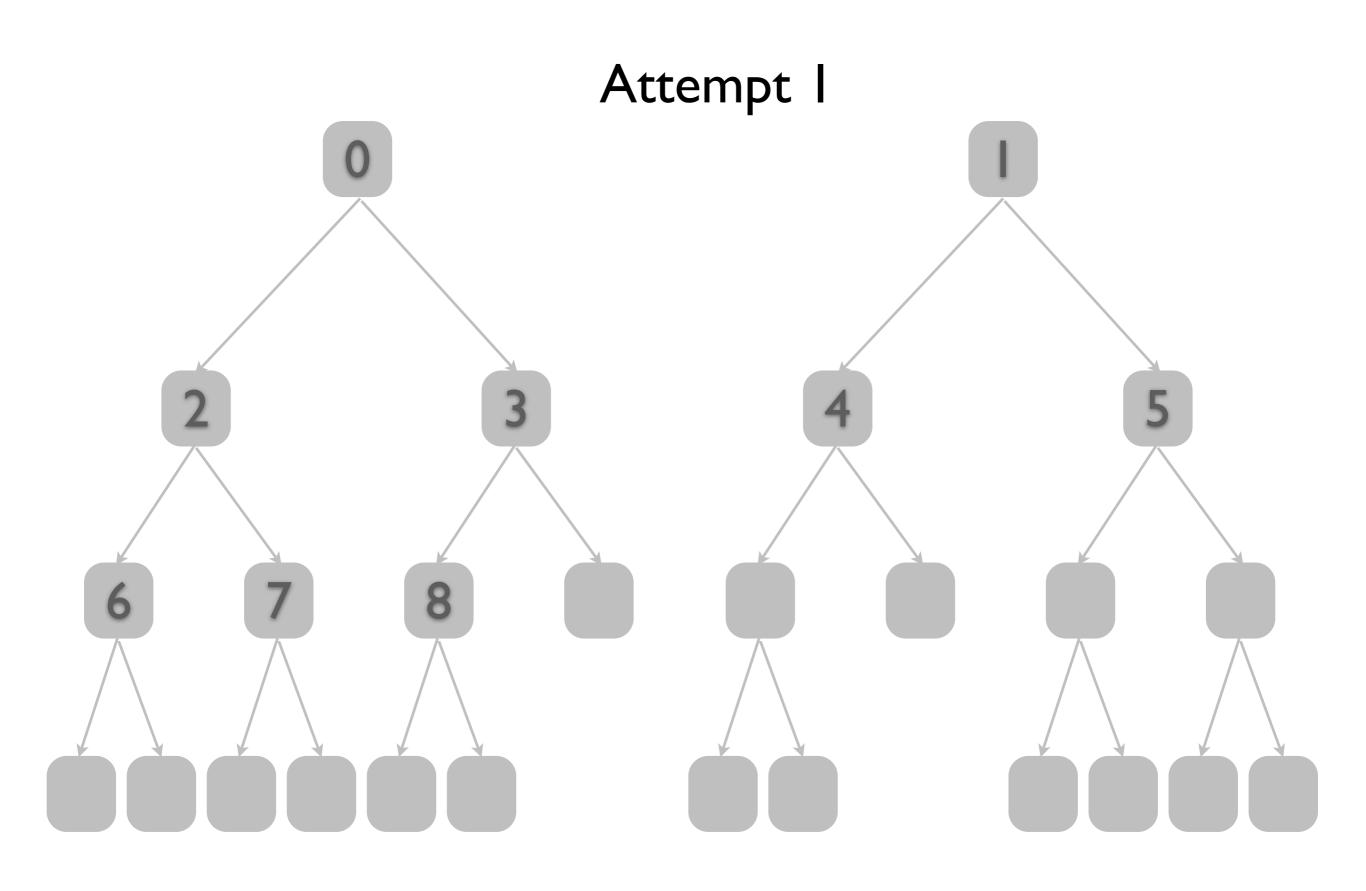
"dumb client, smart server"

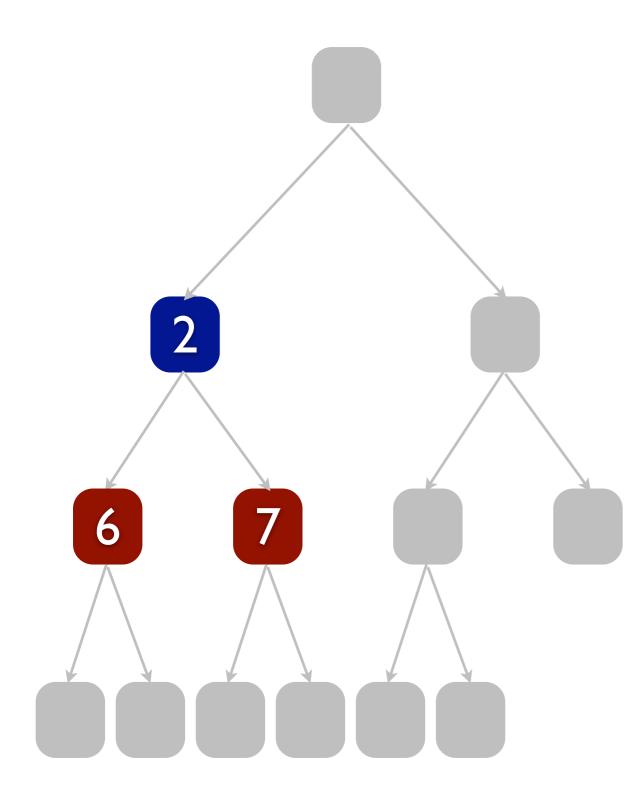
does not scale

Receiver-driven Approach



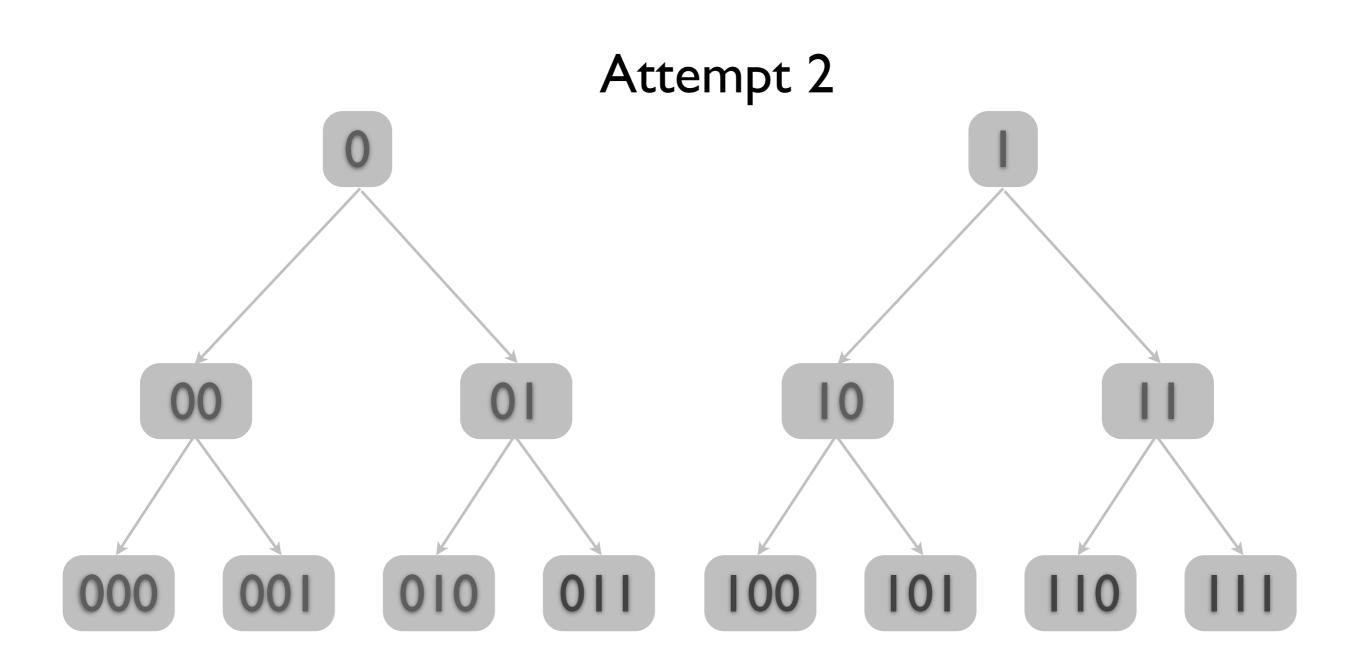
how to identify a vertex split?



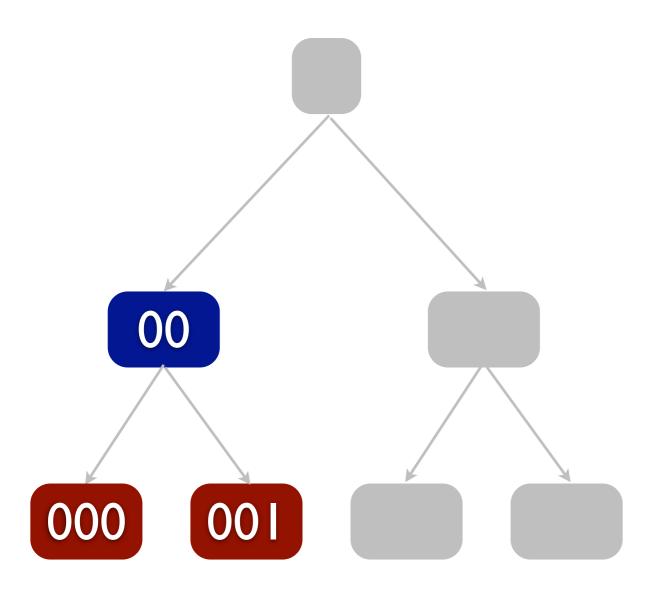


want to split vertex 2

here is how to split, and 2 splits into 6 and 7



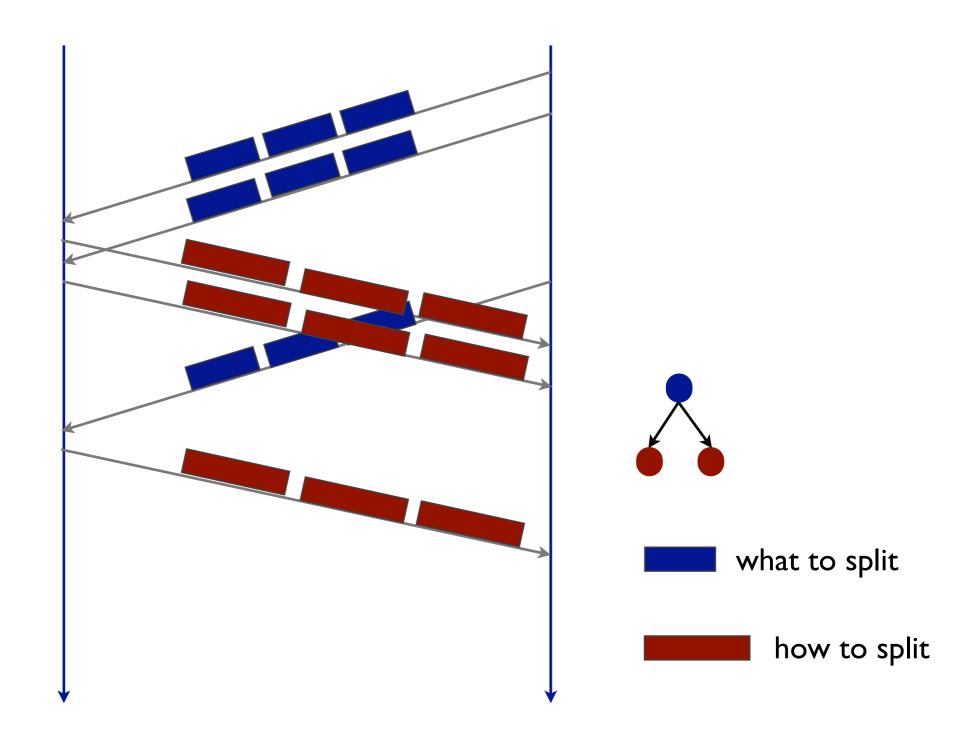
Kim, Lee, "Truly selective refinement of progressive meshes," In Proceedings of Graphics Interface, pages 101–110, June 2001

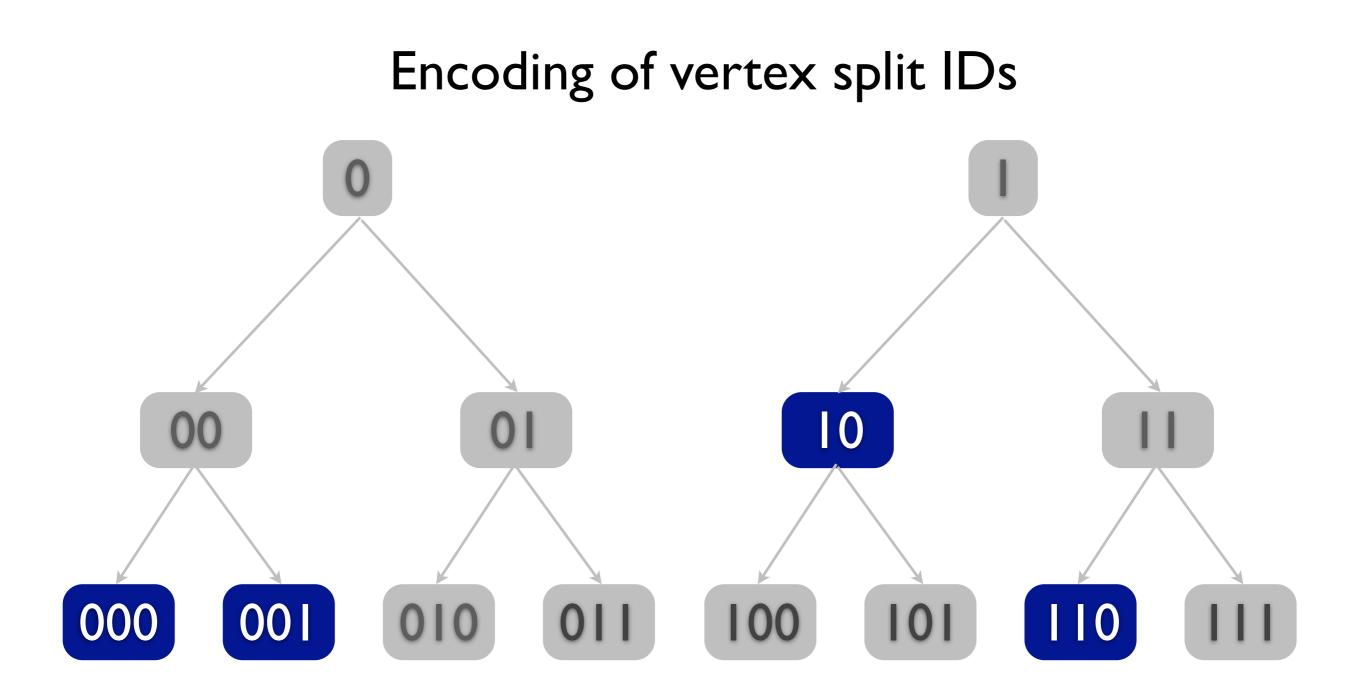


want to split vertex 00

here is how to split 00

Receiver-driven Approach





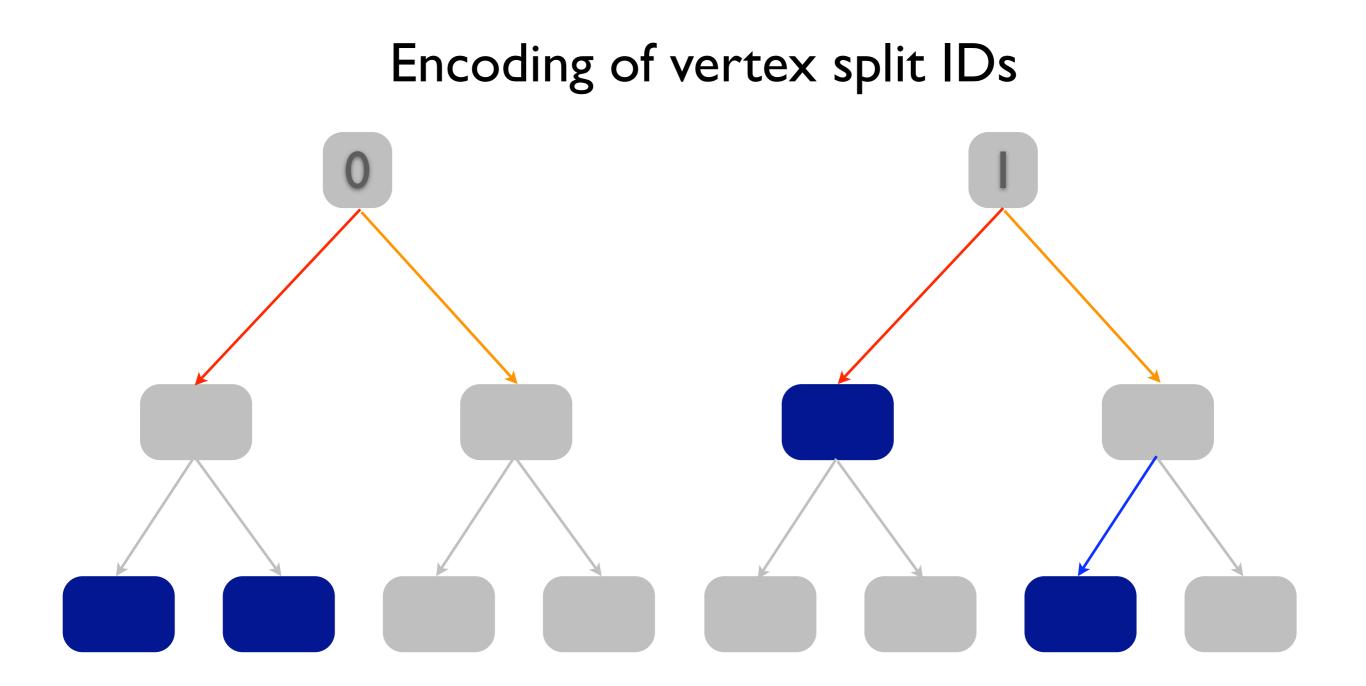


proc encode(T)

if no vertices to be split in T return 0

else

return I + encode(T.left) + encode(T.right)

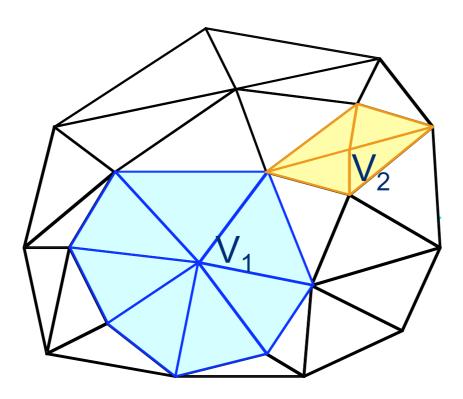


0 | 100 | 000 | 100 | 100

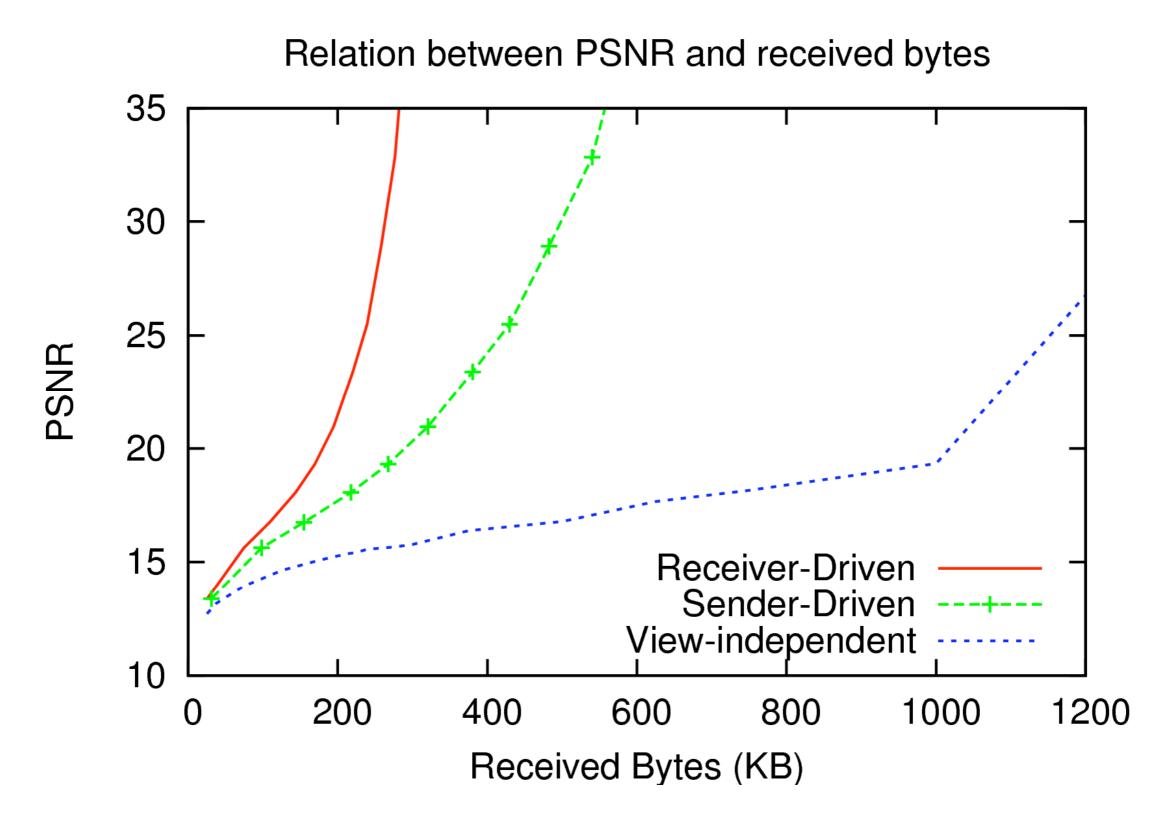
how to compute visibility + visual contributions?

(without possessing the complete mesh?)

Estimate with screen space area of vertices



	Sender-driven	Receiver-driven
send base mesh	1.4	1.13
decode IDs	-	1.55
search vertex split	I.85	1.85
determine visibility	0.41	-
update state	1.41	-
encode IDs	0.94	-
others	0.16	0.16
total	6.17	4.69



receiver-driven protocol alleviates the **computational** bottleneck at the sender.

the other bottleneck is **bandwidth**.

goal: reduce server overhead by retrieving vertex splits from other clients if possible **difficulty**: need to quickly and efficiently determine who to retrieve the vertex splits from low server overhead

low response time

low message overhead

common P2P techniques:

- I. build an overlay and push
- 2. use DHT to search for chunks
- 3. pull based on chunk availability

common P2P techniques:

1. build an overlay and push

2. use DHT to search for chunks

3. pull based on chunk availability

peer-to-peer file transfer:

a needed chunk is likely to be available in any peer

peer-to-peer video streaming:

a needed chunk is likely available from a peer that has watched the same segment earlier

(temporal locality)

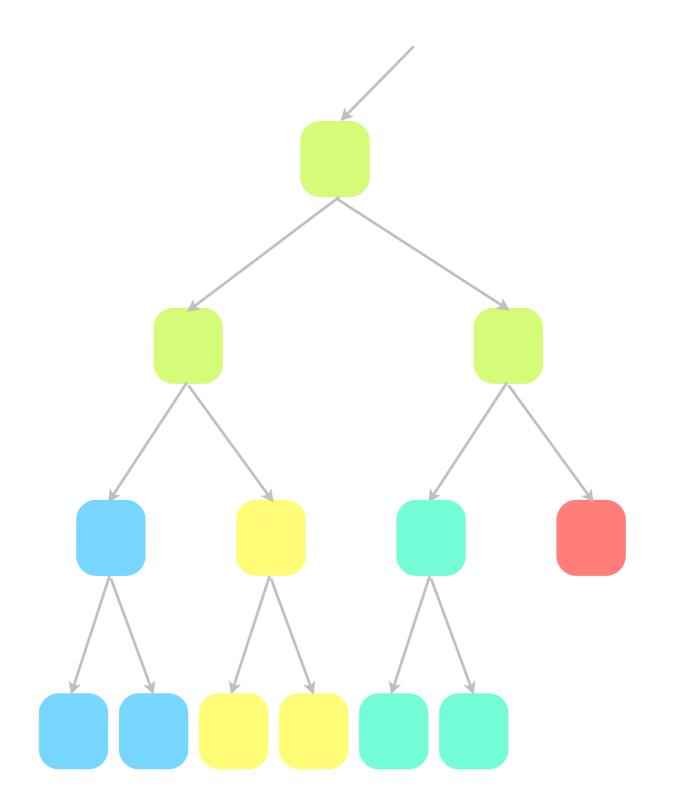
peer-to-peer mesh streaming

a needed chunk is likely available from a peer that is viewing the same region

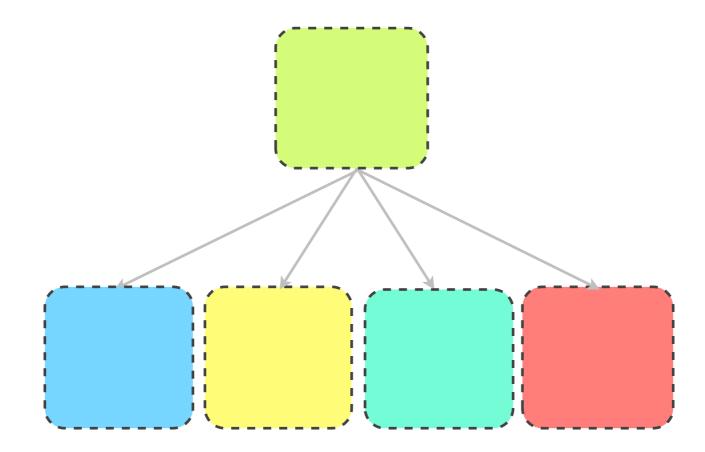
(spatial locality)

idea: exploit spatial locality to reduce message overhead.

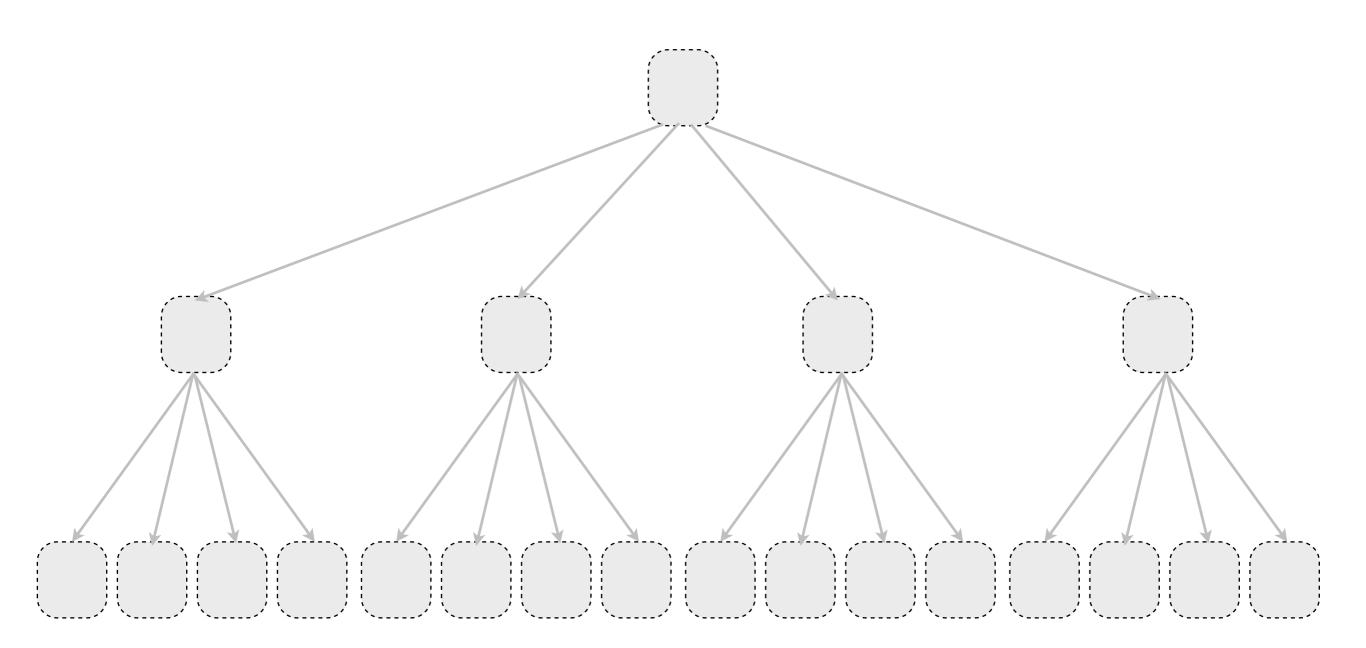
chunks

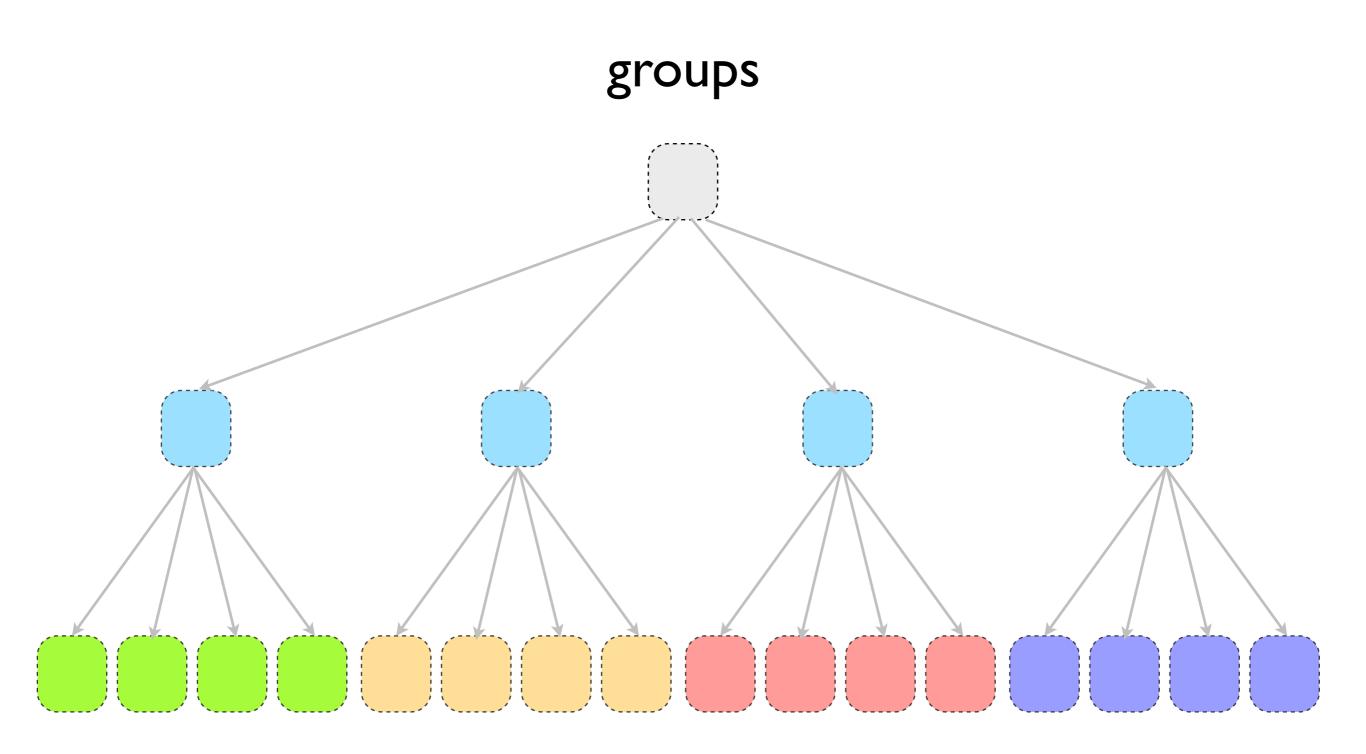


chunks



(I chunk = 240 vertex splits)



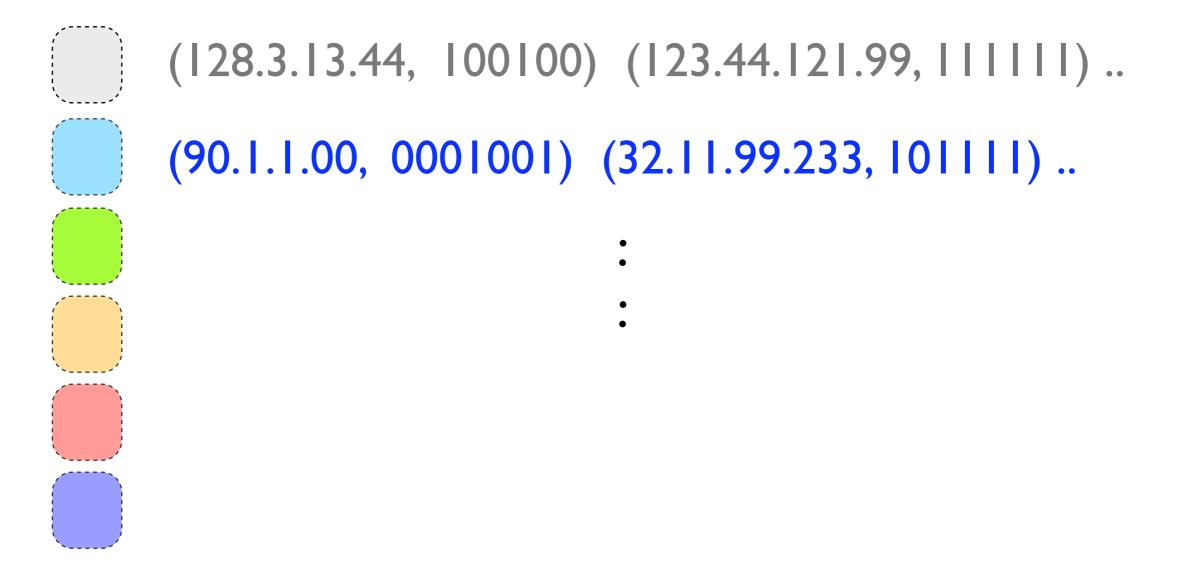


(| group = |6 chunks)

Only exchange messages between peers that need chunks from the same group.

how the protocol works

server maintains a list of group members for each group, and who possesses which chunk.



client: "I want to view mesh M"

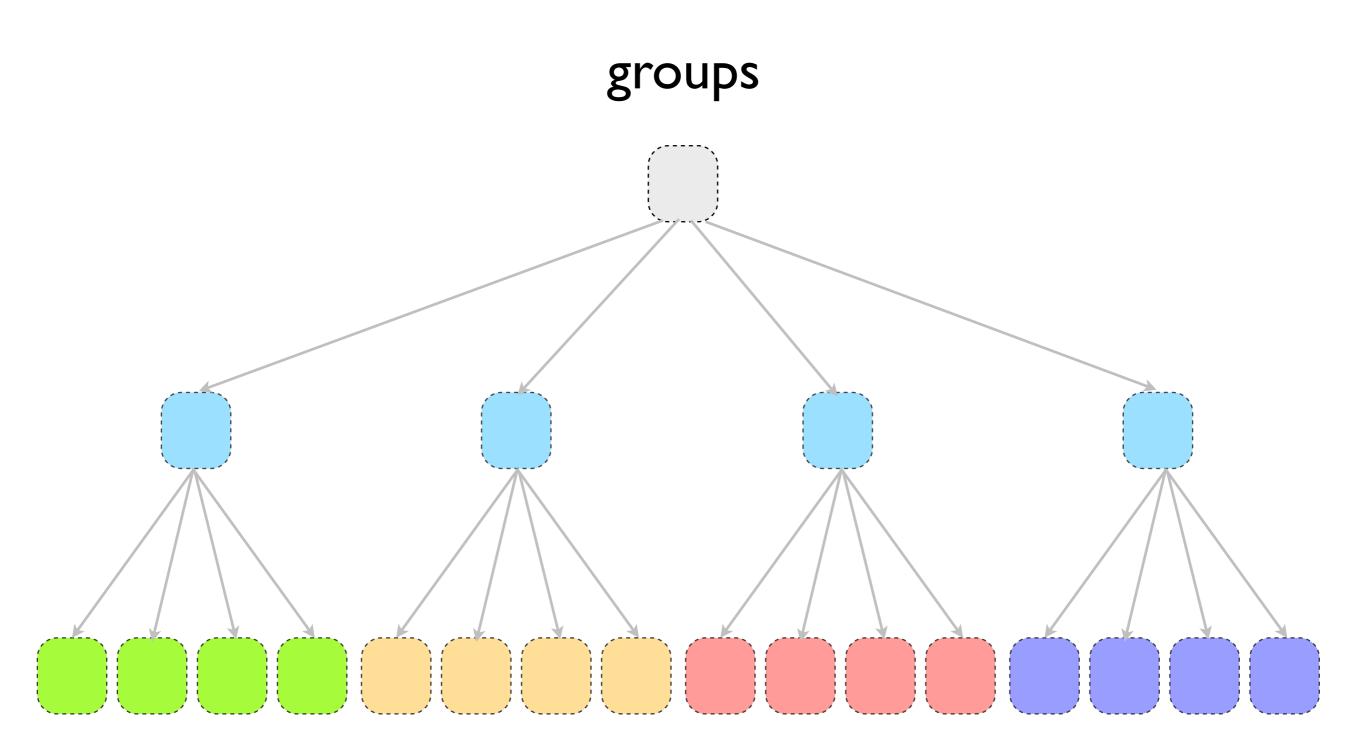
server sends :

- (i) base mesh
- (ii) group members of the highest group.
- (iii) what each member possess

client decides which vertex splits (chunk) to refine

if some peer has that chunk, request from peer else request chunk from server peers inform server when they received a chunk

if a chunk in the next group can be decoded, server sends group members of the next group



if too many group members, server sends only most recent subsets + some seeds

on-going work:

- I. evaluation using user traces and simulator
- 2. other design parameters
- 3. further reduce the role of server

summary

receiver-driven design to reduce CPU cost

peer-to-peer design to reduce bandwidth cost

