

Lecture 7b

Queue ADT

A First-In-First-Out Data Structure

Lecture Overview

■ Queue

- Introduction
- Specification
- Implementations
 - Array Based
 - Linked List Based
- Application
 - Palindrome checking

What is a Queue

- Real life examples
 - A **queue** for movie tickets, Airline reservation **queue**, etc.
- First item added will be the first item to be removed
 - Has the **First In First Out (FIFO)** property
- Major Operations
 - **Enqueue:** Items are added to the **back of the queue**
 - **Dequeue:** Items are removed from the **front of the queue**
 - **Get Front:** Take a look at the first item

Queue: Illustration



A **queue** of
3 persons



Enqueue a new
person to the **back**
of the queue



Dequeue a person from
the **front of the queue**

Queue ADT: C++ Specification

```
template<typename T>
class Queue {
public:
    Queue();

    bool isEmpty() const;
    int size() const;

    void enqueue(const T& newItem);
    void dequeue();
    void getFront(T& queueTop) const;

private:
    // Implementation dependant
    // See subsequent implementation slides
};
```

Queue ADT: Design Considerations

- How about the common choices?
 - Efficiency of **array based implementation**
 - Removing item at the head is the worst case
 - Adding item at the back is the best case
 - Efficiency of **singly linked list implementation**
 - Removing item at the head is the best case
 - Adding item at the back is the worst case
- Is it possible to have both efficient ***enqueue()*** and ***dequeue()*** operations?

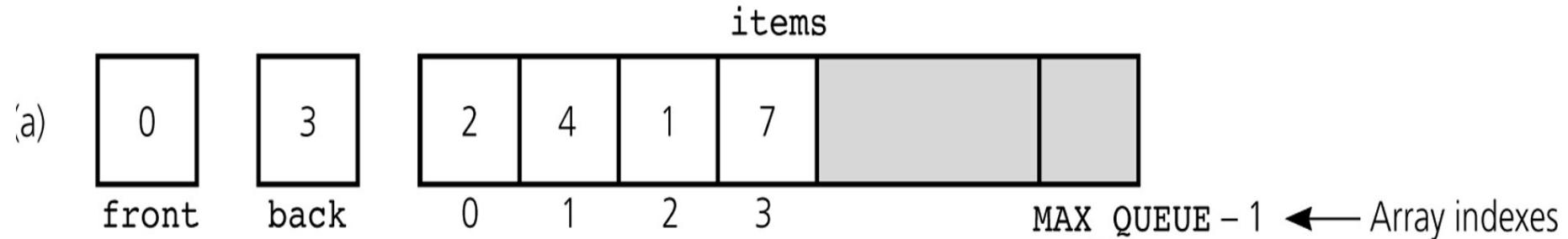
Queue ADT using Array

Array Implementation Issues

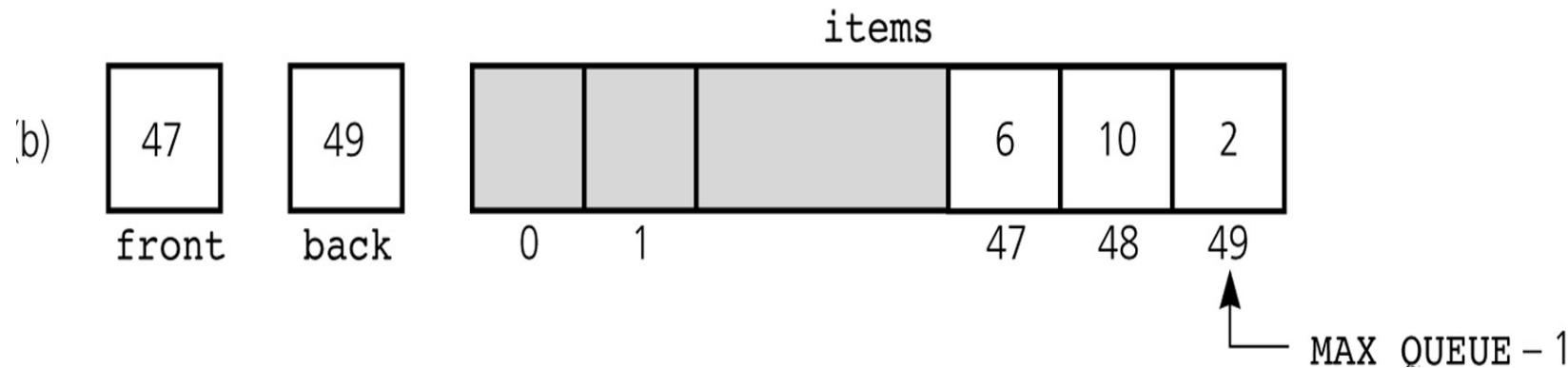
- Removing item from the front is inefficient
 - **Shifting items is too expensive**
- Basic Idea
 - The reason for shifting is
 - Front is assumed to be at index 0
 - Instead of shifting items
 - **Shift the front index**
- So, we have two indices
 - **Front:** index of the queue front
 - **Back:** index of the queue back

Incorrect Implementation

- At the beginning, with 4 items queued



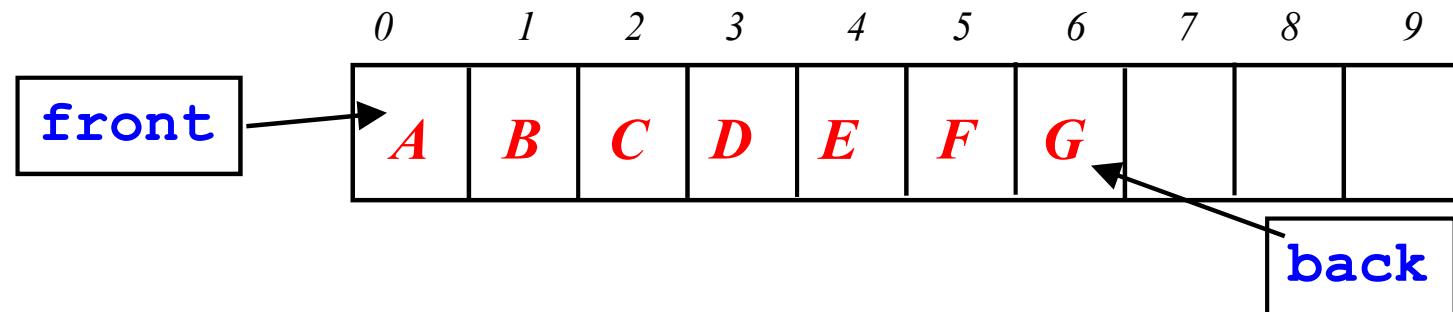
- After many queue operations



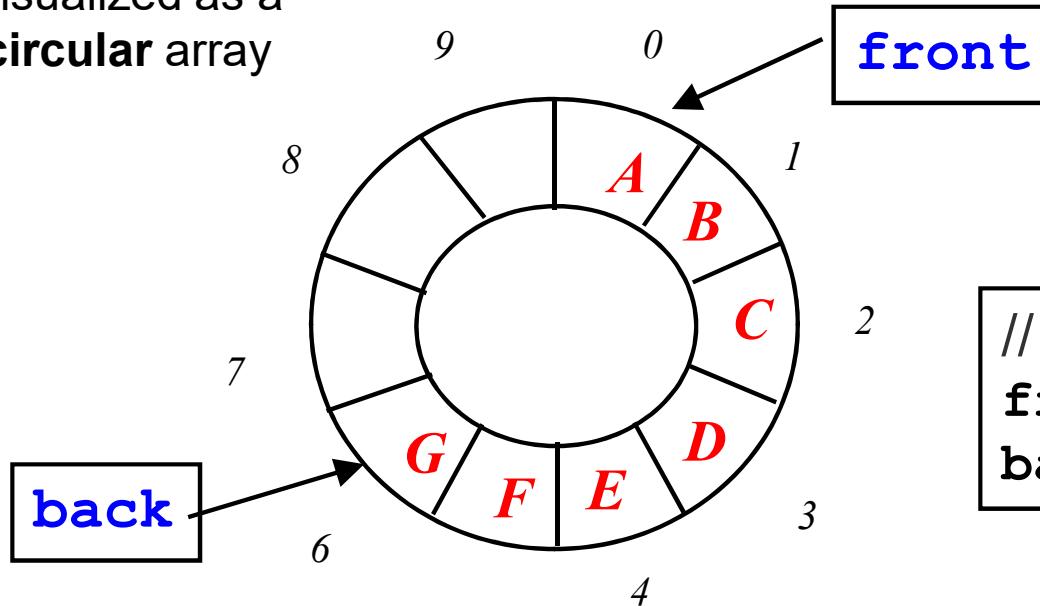
- The front index will drift to the right,
 - Most array locations empty and unusable

Circular Array

- Allow both indices to “wrap” back to index 0 when they reached the end of array
 - Effectively making the array “circular”



Visualized as a
circular array



```
// Advancing front and back index
front = (front+1) % maxsize;
back = (back+1) % maxsize;
```

Queue ADT (Array): C++ Specification

```
const int MAX_QUEUE = 50; // here is the main problem of array
                         // our queue cannot be that large

template<typename T>
class Queue {
public:
    Queue();
    bool isEmpty() const;
    int size() const;

    void enqueue(const T& newItem);
    void dequeue();
    void getFront(T& queueFront) const;

private:
    T _items[MAX_QUEUE];
    int _front, _back, _count;
};
```

QueueA.h

Implement Queue ADT (Array): 1/2

```
#include <string>
using namespace std;
const int MAX_QUEUE = 50;
template<typename T>
class QueueA {
public:
    QueueA() {
        _front = 0;
        _back = MAX_QUEUE-1;
        _count = 0;
    }
    bool isEmpty() const { return _count == 0; }
    int size() const { return _count; }
    void enqueue(const T& newItem) {
        if (_count == MAX_QUEUE)
            throw string("Queue is full");
        else {
            _back = (_back+1) % MAX_QUEUE;
            _items[_back] = newItem;
            ++_count;
        }
    }
}
```

QueueA.h, expanded

Implement Queue ADT (Array): 2/2

```
void dequeue() {  
    if (isEmpty())  
        throw string("Queue is empty");  
    else {  
        _front = (_front+1) % MAX_QUEUE;  
        --_count;  
    }  
}  
  
void getFront(T& queueFront) const {  
    if (isEmpty())  
        throw string("Queue is empty");  
    else  
        queueFront = _items[_front];  
}
```

QueueA.h, expanded

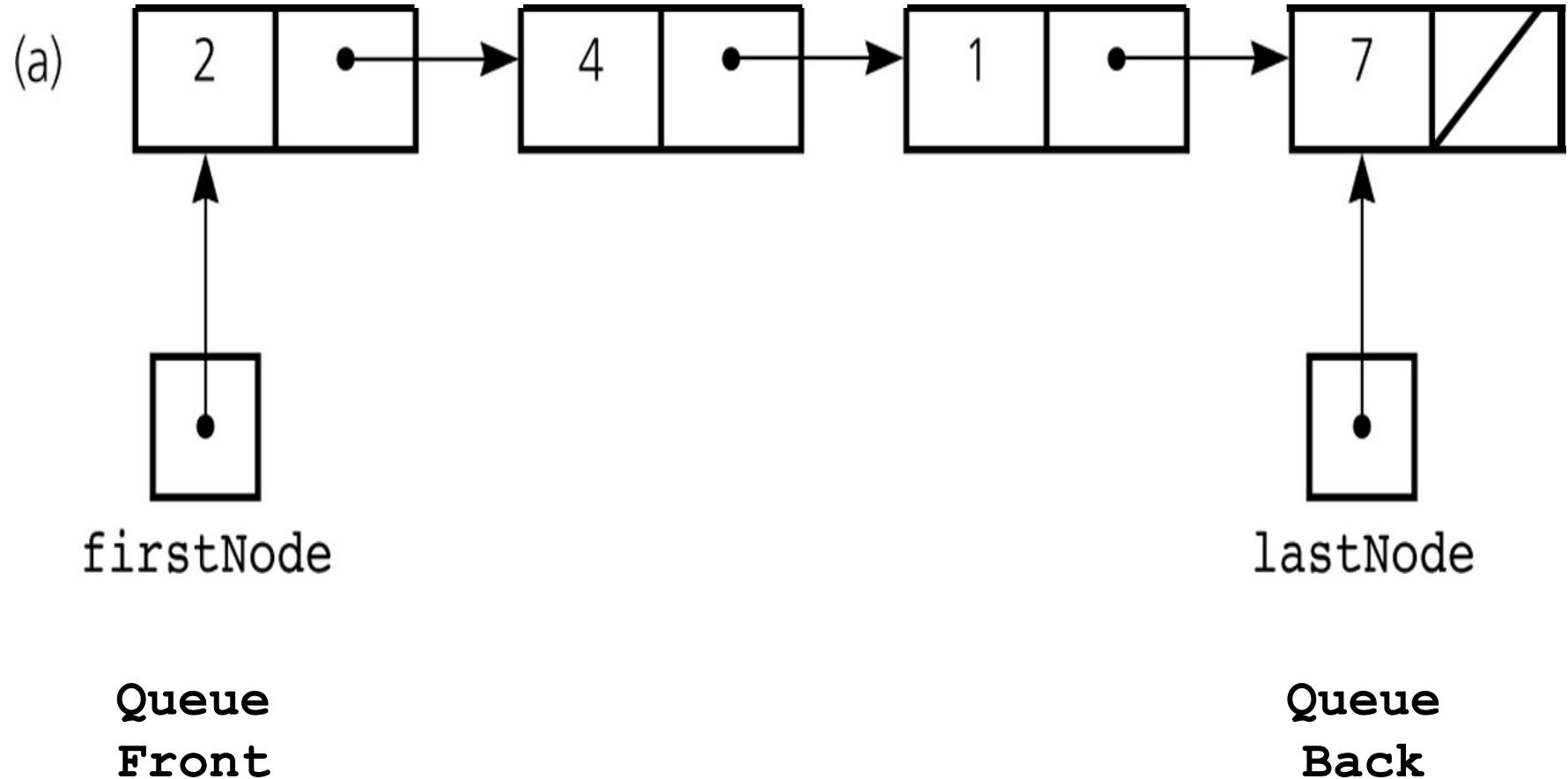
Queue ADT using Modified Linked List

Conceptual Discussion Only

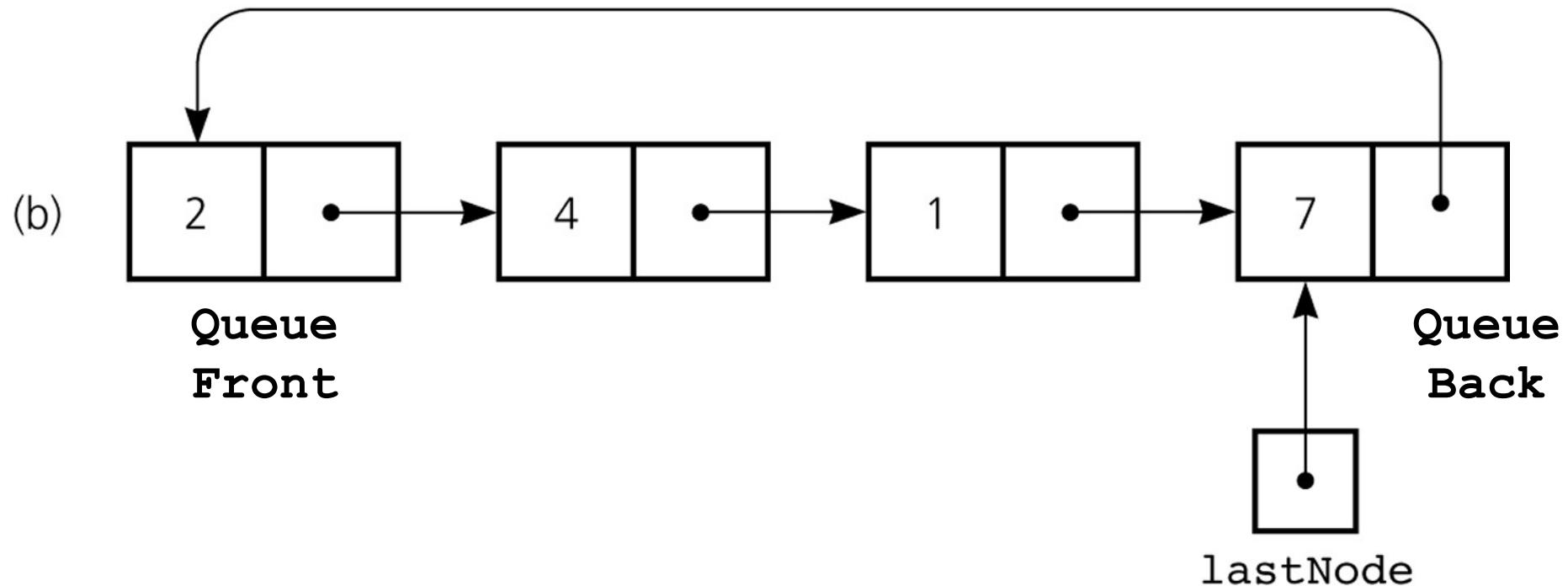
Improving the Singly Linked List

- Singly linked list performs badly for ***enqueue()***
 - Need to traverse all the way to the last node
 - Takes longer time as the queue grows
- How to avoid the traversal to the last node?
 - Easy: Just need to “know” where the last node is all the time
- **Solutions**
 - Keep an additional pointer to the last node, OR
 - Circular linked last with a tail pointer

Linked List: with “head” and “tail”



Circular Linked List



- Only keep track of `lastNode` (tail) pointer
 - `firstNode` pointer can be set when needed
 - `firstNode = lastNode->next;`
- Will use circular linked list for subsequent discussion

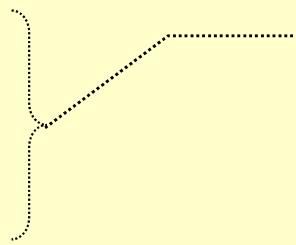
Queue ADT: C++ Specification

```
template<typename T>
class Queue {
public:
    Queue();
    ~Queue();

    bool isEmpty() const;
    int size() const;

    void enqueue(const T& newItem);
    void dequeue();
    void getFront(T& queueTop) const;

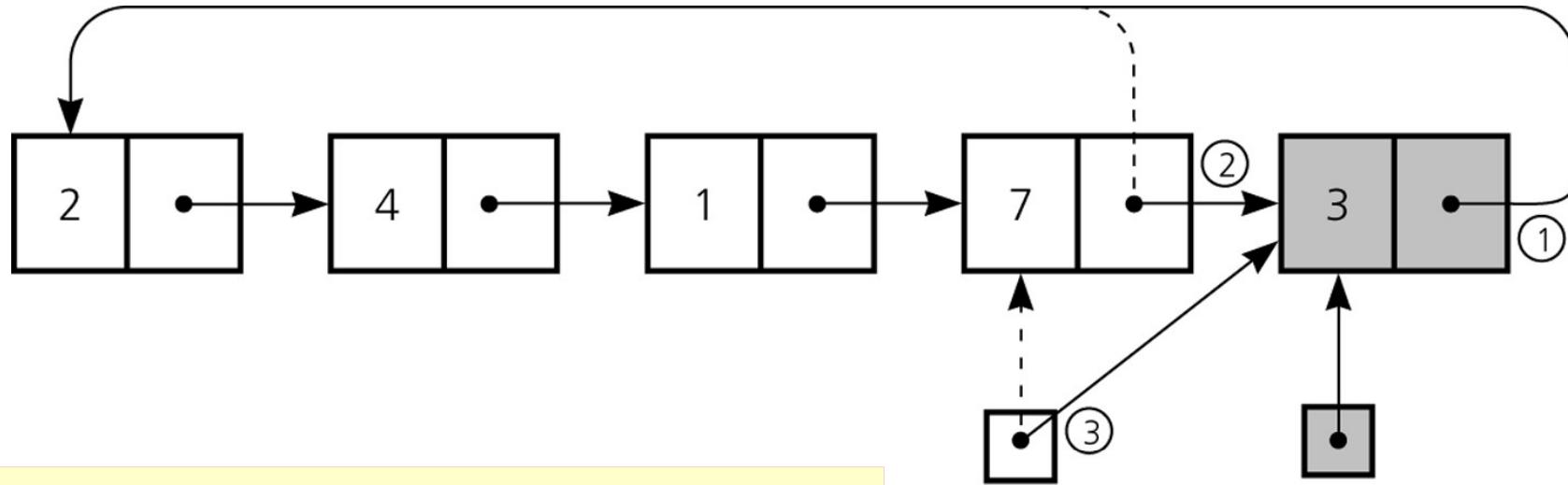
private:
    struct QueueNode {
        T item;
        QueueNode *next;
    };
    int _size;
    QueueNode *_lastNode;
};
```



Just like a **ListNode** structure, yes we can use inheritance but ListLL.h in Lecture6 is an SLL

QueueLL.h

Insertion: Non-Empty Queue



Step 1:

```
newNode = new QueueNode;  
newNode->next = lastNode->next;  
newNode->item = 3;
```

Step 2:

```
lastNode->next = newNode;
```

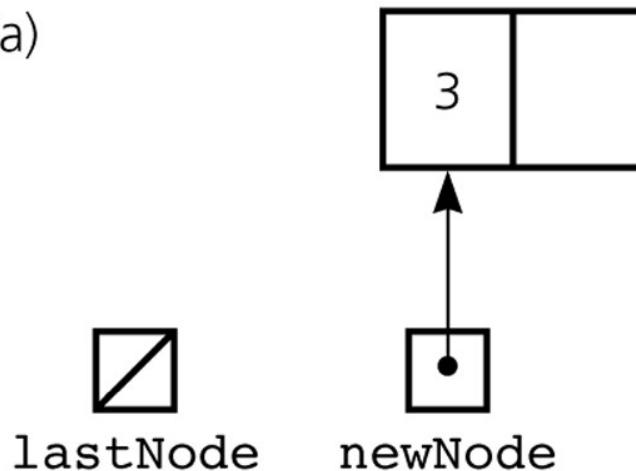
Step 3:

```
lastNode = newNode;
```

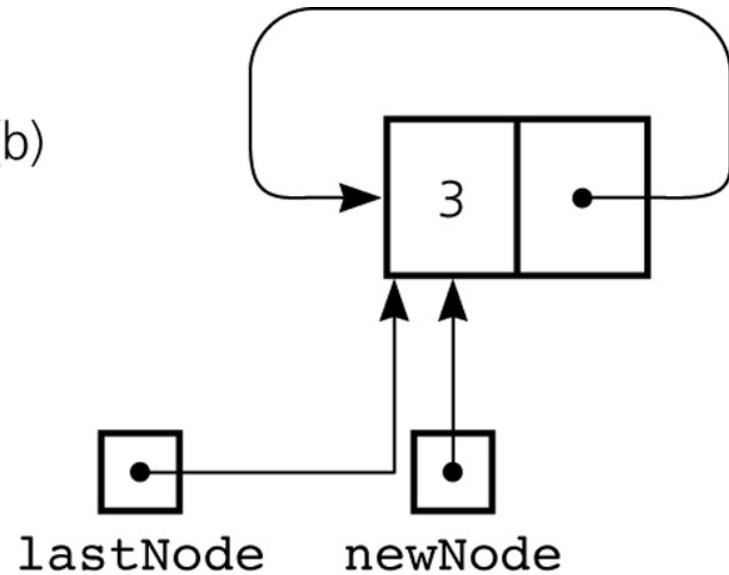
This value is just
an example only

Insertion: Empty Queue

(a)



(b)



Step (a):

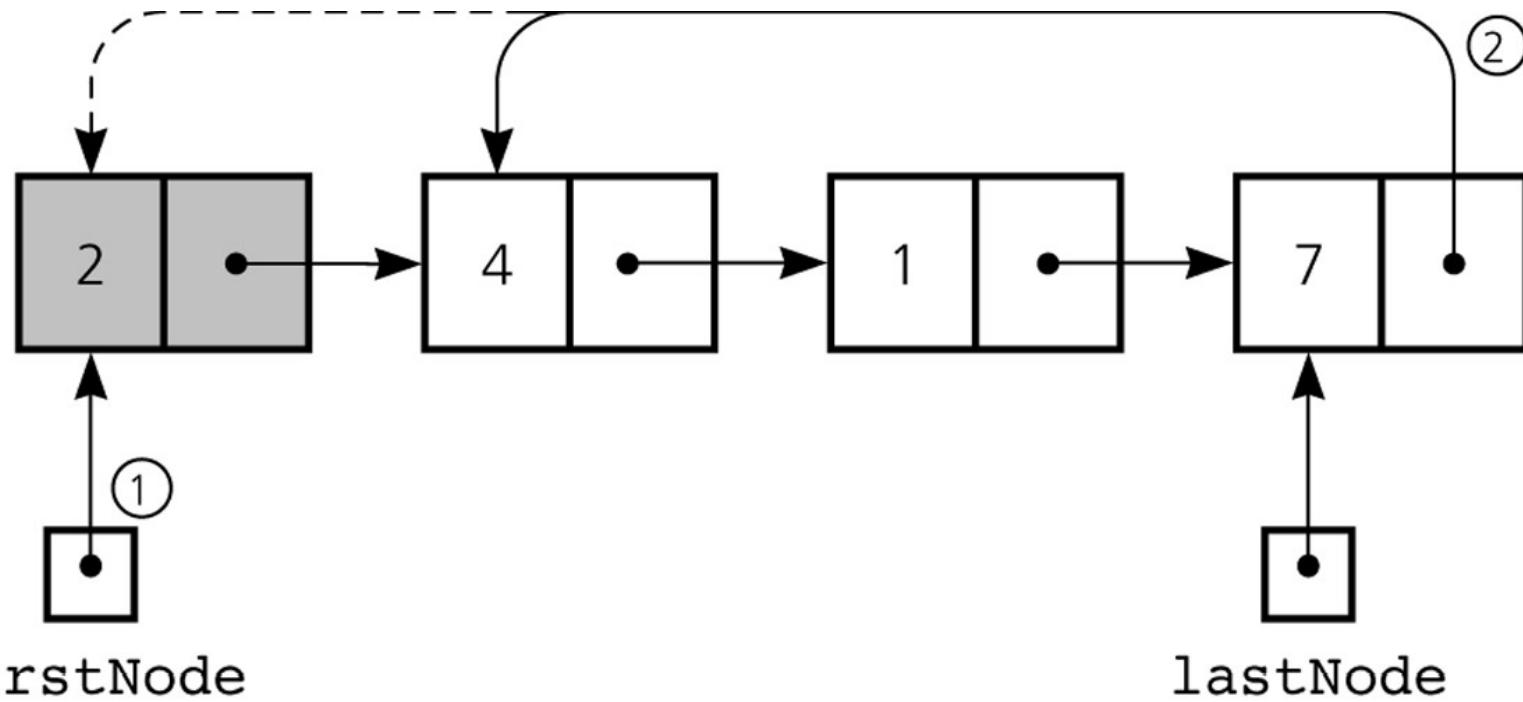
```
newNode = new QueueNode;  
newNode->item = 3;
```

Step (b):

```
newNode->next = newNode;  
lastNode = newNode;
```

Set up the “loop”

Deletion: Queue size larger than one



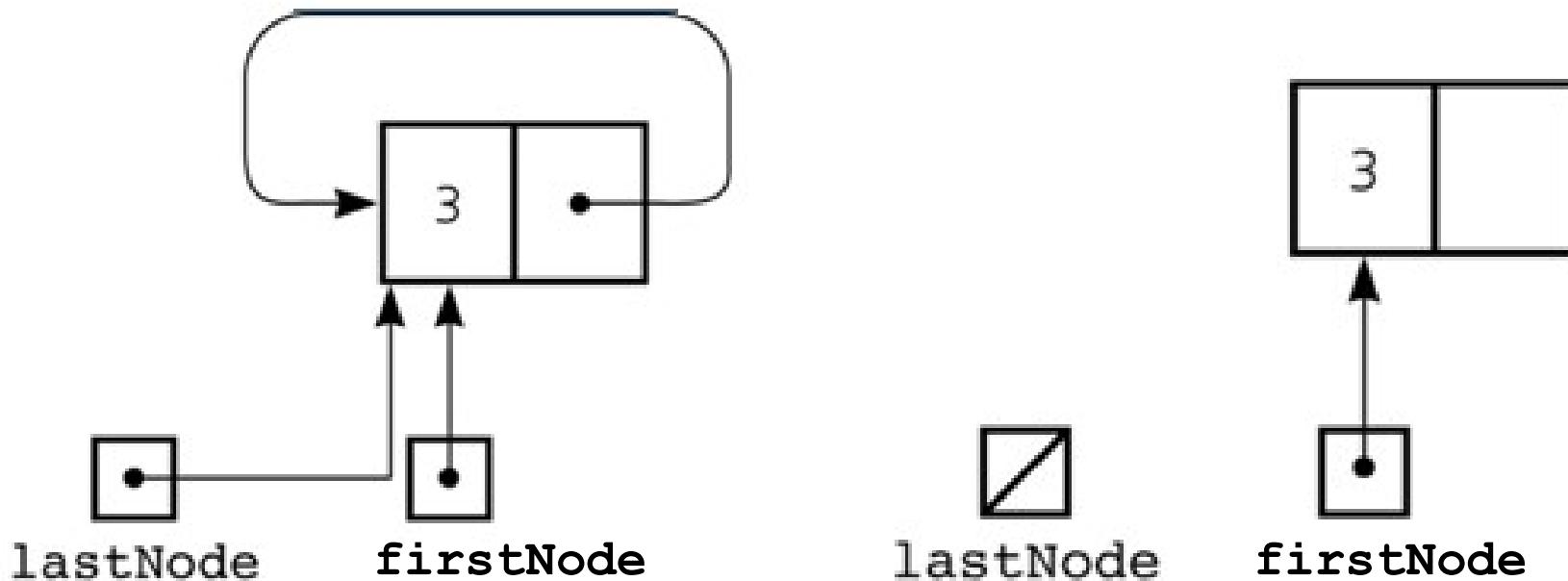
Step 1:

```
QueueNode* firstNode = lastNode->next;
```

Step 2:

```
lastNode->next = firstNode->next;  
delete firstNode;
```

Deletion: Queue size equal to one?



Step 1:

```
QueueNode* firstNode = lastNode->next;
```

Step 2:

```
lastNode = null;  
delete firstNode;
```

STL queue

You should have guessed it by now
STL has a built-in queue ADT

<http://en.cppreference.com/w/cpp/container/queue>

STL queue: Specification

```
template <class T>
class queue {
public:
    bool empty() const;
    size_type size() const;

    T& front();
    T& back();

    void push(const T& t);
    void pop();
};
```

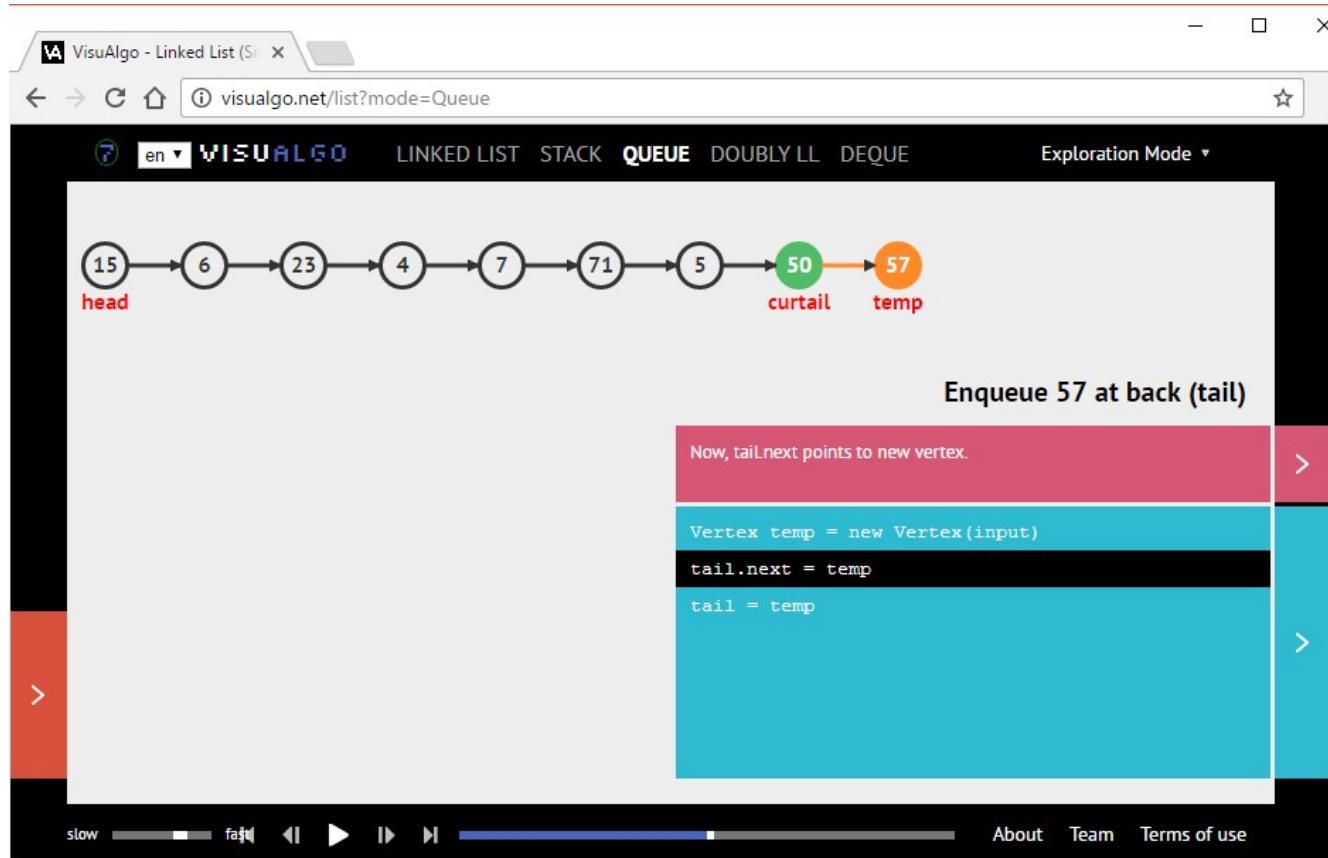
We can see both
front and back

enqueue() is known as
push() in STL Queue

This is the dequeue()
equivalence

VisuAlgo

- <http://visualgo.net/list?mode=Queue>
- I use Tailed Linked List approach instead of Circular Linked List but the idea is the same



Queue Application

Checking for Palindrome

Palindrome: Problem Description

- **Palindrome** is a string which reads the same either *left to right*, or *right to left*
 - **Palindromes:** “r a d a r” and “d e e d”
 - **Counter Examples (most random strings):** “d a t a”
- Many solutions
 - But for the sake of discussion, let's use the two newly learned ADTs
 - Highlight the difference of **LIFO** and **FIFO** property
- Main Idea
 - Use **stack** to reverse the input
 - Use **queue** to preserve the input
 - The two sequence should be the same for palindrome

Palindrome: Implementation

```
#include <queue>
#include <stack>
using namespace std;

bool palindrome(string input) {
    stack<char> s ;
    queue<char> q ;

    for (int j = 0; j < input.size(); j++) {
        s.push(input[j]);
        q.push(input[j]);
    }

    while (!s.empty()) {
        if (s.top() != q.front())
            return false;
        s.pop();
        q.pop();
    }

    return true;
}
```

Push the same character into both queue and stack

Queue has the original sequence, Stack has the reversed. Compare to make sure they are the same

Summary

